Vacuums in the Cloud: Analyzing Security in a Hardened IoT Ecosystem

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Abstract

With the advent of robot vacuum cleaners, mobile sensing platforms entered millions of homes. These gadgets not only put “eyes and ears” into formerly private spaces, but also communicate gathered information into the cloud. Furthermore, they reside inside the customer’s local network. Hence, they are a prime target for attacks and if compromised become a privacy and security nightmare. Vendors are aware of robots being a target of interest; they employ various security mechanisms against tampering with devices and recorded data in the cloud.

In this paper, the Neato BotVac Connected and Vorwerk Kobold VR300 ecosystems are analyzed and the robot firmware is reverse engineered. To achieve the latter, a technique to bypass the devices’ secure boot process is presented revealing the firmware, which is then dissected to evaluate device-specific secret key generation and to trace vulnerabilities. We present flaws in the secret key generation and provide insight on the occurrence and exploitation of a buffer overflow, which give an attacker complete control not only in the local network but also via the robots’ cloud interface. Eventually, multiple attacks based on the findings are described and security implications are discussed. We shared our findings with the vendors, who further increased their otherwise commendable security mechanisms, and hope more vendors can take away valuable lessons from this highly complex Internet of Things (IoT) ecosystem.

1 Introduction

Smart vacuum cleaning robots contain many interesting sensors, such as laser scanners, ultrasonic distance sensors, cameras, and bump detectors. Combined with cleaning schedules and result maps, this exposes a substantial amount of private information about their owner—accordingly, iRobot’s announcement to sell Roomba data for marketing purposes has been discussed controversially [2]. With the customer’s shipping address, WiFi configuration, or simply position data collected by the robot or the app, apartment size and furniture can be associated with locations. All market leading vacuum cleaning robots are embedded into complex IoT ecosystems. Xiaomi robots were subject of significant security research including an open source replacement for their proprietary cloud components [7]. However, Xiaomi uses a different infrastructure, hardware architecture, and operating system—meaning that while Neato and Xiaomi are functionally similar, the security measures have to be different.

With a market share of 62.7 million € in 2017, Neato is one of the bigger players on the vacuum robot market [20]. Vorwerk acquired Neato in 2017 and sells a robot with almost identical firmware, which is part of the vacuum cleaning product line Kobold listed with 800 million € revenue in total.

Even though vacuum cleaning robots are a target of interest, not much security research has been published on Neato, except for an exploit chain for attackers in the same local network [10, 11]. For Vorwerk robots, which are re-branded Neato robots with minor hardware and settings upgrades, not a single Common Vulnerabilities and Exposures (CVE) number is known. Despite robot vacuums not being broadly considered to be security critical, Neato put plenty of advanced security techniques in place, which are interesting to analyze and learn from for all kinds of IoT devices. We show that the combination of multiple vulnerabilities in different components leads to a powerful attack compromising the entire ecosystem. Our main contributions are:

- Secure boot bypass and extraction of Random Access Memory (RAM) with previous contents,
- execution of commands on robots over the cloud with superuser privileges without authorization,
- massive entropy reduction in robot secret key generation,
- decryption of RSA private keys on robots and enabling arbitrary command and control channels over their cloud,
- information leakage of end-user smartphone’s public Internet Protocol (IP) addresses.
All security issues were communicated to the vendor. Remote exploitation over the cloud was filtered immediately, and firmware updates fixing bugs closer to hardware were rolled out until June 2019. Adequate patches on some encryption flaws required changes in the device manufacturing process.

Our research highlights that obscurity and the usage of strong encryption methods without proper knowledge of the underlying schemes does not provide sufficient security. Our analysis and recommendations are practice-oriented and of interest to other vendors and security analysts. This is especially relevant as some technologies utilized by Neato are also present in the automotive and other interconnected sectors. We encourage vendors to work on open solutions for security and establish better standards. In the long run this will solve security issues similar to the ones found in our analysis.

This paper is structured as follows: Section 2 gives a brief overview on the robot ecosystem components and the messages exchanged with special emphasis on security. While Section 3 focuses on exploitation with local hardware or network access, Section 4 shows cloud exploitation techniques. Security implications of our findings are discussed in Section 5. Section 6 concludes our work.

## 2 The Neato and Vorwerk Ecosystems

Advanced smart robot vacuum cleaners are able to clean the house almost autonomously. The only parts requiring interaction are emptying the dirt bin and starting the robot. Neato advertises the connected series with the ability to control the robot and get notifications via smartphone as well as obtaining a cleaning summary [16]. To achieve this, the ecosystem features multiple components displayed in Figure 1. The smartphone app communicates with two different application programming interfaces (APIs), namely Beehive and Nucleo. Beehive (beehive.neatocloud.com) manages all account-related tasks, such as account information and linked robots. The robot itself communicates with the Beehive cloud during initial pairing to exchange authentication information. In contrast, Nucleo (nucleo.neatocloud.com) provides robot functionality. It receives commands from the smartphone app and forwards them to the robot. In reverse direction, Nucleo provides the app with robot status messages. Moreover, a subset of commands can be sent to the robot directly without using the cloud in manual mode. The robot platform itself is based on a QNX system [3].

Neato, like most vendors, does not provide public details about the inner life of their ecosystem. All information in this section was gathered either by sniffing network traffic or by reverse engineering the Android smartphone app. Additionally, Neato provides a developer API documentation [17].

We dissected multiple products, which we compare in Section 2.1. Section 2.2 explains further details on the Beehive cloud while Section 2.3 focuses on the Nucleo cloud.

### 2.1 Product Comparison

<table>
<thead>
<tr>
<th>Product Name</th>
<th>Firmware Version</th>
<th>Build Date</th>
</tr>
</thead>
<tbody>
<tr>
<td>Neato BotVac Connected</td>
<td>2.0.296</td>
<td>Dec 12 2016</td>
</tr>
<tr>
<td>Neato BotVac Connected D7</td>
<td>4.2.0-102</td>
<td>Jul 12 2018</td>
</tr>
<tr>
<td></td>
<td>4.4.0-72</td>
<td>Dec 22 2018</td>
</tr>
<tr>
<td>Vorwerk Kobold VR300</td>
<td>4.2.2-137</td>
<td>Aug 30 2018</td>
</tr>
<tr>
<td></td>
<td>4.2.5-166</td>
<td>Jan 11 2019</td>
</tr>
</tbody>
</table>

Table 1: Firmware version comparison. Neato is used to test new features, Vorwerk is getting stable releases only.

To confirm and compare our findings, we tested three vacuum cleaning robot models. Vorwerk robots are very similar to Neato; in fact they use the same firmware just with different settings applied. Firmware versions are listed in Table 1. The Vorwerk Beehive and Nucleo cloud are hosted on beehive.ksecosys.com and nucleo.ksecosys.com. Vorwerk hardware has some upgrades compared to Neato hardware, such as more ultrasonic sensors, a larger battery, different power and speed settings for normal and eco vacuum cleaning mode, a green brush, and rainbow colored LEDs.
2.2 Beehive Cloud

Beehive is responsible for all account-related information. It holds the secret_keys of all robots that are connected to a user account. The secret_key is required for communicating with the Nucleo cloud and to send manual commands to the robot. Beehive is hosted at Amazon [1] via Heroku [18].

2.3 Nucleo Cloud

The Nucleo API represents the interface for smartphone↔robot communication. It provides robot status information and forwards commands to the Neato BotVac Connected. The Nucleo infrastructure is hosted with Amazon via Heroku as well. All messages to the API are payload-based with the constant path /vendors/neato/robots/[robot_serial]/messages.

Authentication In the Nucleo API, Authorization and Date header fields are used for authentication. The authorization field is generated by concatenating the message body to the date and afterwards concatenating the result to the lower case robot_serial. Then the SHA-256 Keyed-Hash Message Authentication Code (HMAC) [6, 9] of the resulting string is computed, using the robot secret_key. The computation instructions in pseudo-code are given in Listing 1.

On sending a message to a distinct robot, the server checks for Date header validity. If a date is provided and lies within a 20 min time window of the current time, the server checks for presence of the Authorization header. When providing the “NEATOAPP” string in the Authorization header but not giving a HMAC or an invalid HMAC, the server’s response is 

"Could not find robot_serial for specified vendor_name". This behavior also applies when trying to access random robot_serials, which means it is not possible to confirm the existence of robot_serials in the Nucleo API. Note that the server does not check the Authorization header validity itself, instead the request is forwarded to the robot which then validates the HMAC. To summarize, anyone who is in possession of a robot_serial and secret_key is able to send commands to the corresponding robot.

```python
authorization = "NEATOAPP " + signature
signature = HMAC_SHA256(robot_secret_key,
UTF8_encoded(string_to_sign))
string_to_sign = lower(robot_serial) + "\n" +
date_header + "\n" + body
```

Listing 1: Neato Nucleo authentication between smartphone and backend [17].

3 Local Vulnerabilities

We use a new secure boot bypass technique in Section 3.1 to extract system image and RAM contents. Static system image analysis in Section 3.2 is the base for findings across the remaining part of the paper, which includes extensive secret_key generation entropy reduction in Section 3.3.

3.1 Bypassing Secure Boot

Access to the firmware running on the robot is required for more advanced security research. This includes but is not limited to analysis of mutually authenticated Transport Layer Security (TLS) connections and debugging of running applications. Neato is using a custom AM335x chip with secure boot enabled [14]. The flash chip only contains an encrypted and signed system image; a physical attacker cannot simply unsolder the chip to obtain it. Firmware updates use the same encrypted and signed format. Changing and debugging firmware on a robot therefore requires bypassing secure boot.

We start with a hardware analysis in Section 3.1.1, which leads to a hidden boot menu in Section 3.1.2. We exploit this to extract Neato’s system image (see Section 3.1.3). The described security issue is filed as CVE-2018-20785. It has been fixed in the Neato BotVac Connected D7 firmware version 4.4.0-72, all other versions listed in Table 1 are vulnerable. As of June 2019, also the Vorwerk firmware has been fixed.

3.1.1 Hardware Analysis

After teardown of the Neato BotVac Connected, further analysis of all components is required. The motherboard has many testing points and undocumented chips, with an important part of it depicted in Figure 2. We connected all testing points to a logic analyzer, and rebooted the robot multiple times. Most of the testing points do not show any interesting output, some might be for sensors and other features not relevant for system access. In addition to test points, the right front of the motherboard contains three extra large unlabeled pins, which are a serial interface with clock, read, and write. The ground needs to be taken from another source, such as the internal Universal Serial Bus (USB) port. It prints the following text during startup and remains silent afterward:

```
BotVac Connected "<Fast Memory>"(84038)
```

The Neato BotVac Connected D7 features a similar serial port, it is located on the motherboard’s left and already has four connection pins. Access to the Vorwerk Kobold VR300 serial port does not require disassembling it, it is located next to the USB port.

The main chip on the Neato Botvac Connected is the SN120952608ZCE in the center of Figure 2. Since it is a custom chip there is no information or documentation available. The form factor 962B leads to the AM335x chip family manufactured by Texas Instruments.
3.1.2 Hidden Boot Menu

During startup, QNX first loads an Initial Program Load (IPL), which then loads an Image File System (IFS) containing the operating system. Instead of a QNX IPL, Linux typically uses the more common setup of MLO and U-Boot on AM335x. Either way, these setups enable to boot an operating system. Secure boot uses encryption keys configured on the AM335x decrypt and verify both, IPL and IFS, prior to execution.

Neato’s standard setup is to simply boot the IFS containing the firmware from the external flash chip without showing any options. While testing the partially documented Neato USB command line [15], we found that TestMode On followed by SetSystemMode PowerCycle enables a hidden IPL boot menu on the serial interface:

```
BotVac Connected "<Fast Memory>"(84038)
Press enter twice within the next 2 seconds for boot menu
**Commands:**
Press 'M' to load IFS from main image flash
Press 'X' for serial download, using XModem-1k
Press 'S' to scan existing memory without download
```

Note that XModem-1k download makes the robot download a system image from the serial port and boot it. This does not enable system image extraction. On uploading a valid QNX IFS no secure boot checks are performed, representing a bypass. Loading a BeagleBone QNX IFS from Foundry27 for AM335x works in principle, it prints some information on the serial interface but crashes when initializing processors.

The original AM335x Foundry27 board support package includes a very similar IPL procedure in src/hardware/ipl/boards/am335x/main.c. It enables sending an IFS over the QNX serial protocol sendnto. Neato’s implementation represents a variation and more AM335x-based products might be affected by the secure boot bypass. The more common setup of MLO and U-Boot also allows to load arbitrary images by default. In general, vendors relying on secure boot should carefully check for such contradicting setups.

3.1.3 Memory Extraction

We exploit the secure boot bypass to launch a cold boot attack, which is possible as neither the AM335x chip nor the IPL reset memory contents during startup. When recompiling and booting the Foundry27 image, the main issue is that hardware cannot be initialized. This is since memory is mapped to different addresses on the Neato chip variant and any invalid memory access freezes the chip.

At this stage an attacker is able to print memory contents until an invalid memory access is made. In our experiments, increasing the serial interface speed up to 1 Mbit/s worked. Despite slow speed and partially unknown memory layout, this poses a huge attack surface for cold boot attacks.

By default, QNX is located at address 0x81000000. First memory dumps show that despite booting our own QNX IFS this memory region still contains some of the original Neato IFS contents. We relocate our custom QNX IFS to a free memory region and perform a full memory dump of the upper 2 GB of the address space. This area contains the same 256 MB, which are mapped 8 times to the memory region. Each mapping contains the full Neato IFS and any RAM contents present before the reboot, including WiFi credentials.

3.2 Static Firmware Analysis

Neato’s IFS is a rich source for information and enables static firmware analysis. It contains all important binaries and shell scripts coordinating them. For example, /bin/robot is called by a script /etc/brain.sh and the watchdog /bin/pinky is started via /etc/pinky.sh. An overview of the main tasks is shown in Figure 3. robot is the largest binary and responsible for the main features, starting at the interactive command line and ending at all cleaning logic. It receives Nucleo cloud commands from astro and generates replies. The manual drive mode which is issued locally via the smartphone app is first redirected from webservice to astro and then processed similarly. conman manages WiFi connections. Messages are exchanged via proprietary libNeatoIPC calls as well as standard QNX message queues. In the following we focus on those components relevant for further findings.
Testing the robot’s wireless interfaces produced coredumps
that are separate files in the so-called “Blackbox Logs”, which

are located in

final

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This function calls multiple other functions to finally
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type declarations, we were able to reconstruct the function
as shown in Listing 2. The first important observation is
the usage of the OpenSSL library for computing the final
secret_key, which appears to be a SHA-1 [5] hash. Type
declarations were derived from the OpenSSL documentation.
Contents passed to SHA-1 are the following blocks of bytes:

t_shift[0-3] = time_now

t_shift[4-6] = 0

t_shift[7] = 16

t_shift[8] = rnd + rnd / 0xFFFF

t_shift[9] = ~(~(((unsigned int) unsigned short)

(rnd + rnd / 0xFFFF) >> 8) & 0x3F) << 25) >> 25

t_shift[10-15] = robot_MAC

v_sha1 = EVP_sha1()

Listing 2: Secret Key Generation.

and is used for authenticat-
ing the robot against the cloud components. By examining the
key decryption procedure in the robot binaries, we were able
to decrypt the keys and launched several tests as described in
Section 4.1. Due to the high impact of these keys becoming
public, we agreed to not disclose details on the decryption
process.

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Contents passed to SHA-1 are the following blocks of bytes:

t_shift[0-3] = time_now in seconds.

t_shift[4-6] is set to constant 0.

t_shift[7] is set to constant 16.

t_shift[8] is set to a random value with some mathemati-
cal operation.

t_shift[9] is set to a random value with multiple mathemati-
cal operations. Since one of them is an AND with the
value 0x3F, the first two bit of this byte are set to
constant 0.

t_shift[10-15] is the robot’s Media Access Control
(MAC) address, which is a 6 B hex-string.

In the next step, a hash for the 16 B t_shift variable
is computed. If computing the digest was successful, the hash is then stored at the address contained in
the pointer emmc_keys_robotkey which leads to the file
/emmc/keys/RobotKey. The first 16 B of the robot_key
variable are then written to that file.

In Section 5.1.3 we show that the robot_serial (and
thereby the robot_MAC) has to be considered as known,
which eradicates the entropy of t_shift[10:15]. The same
is true for the constant values, eliminating the entropy for
t_shift[4:7].

The bytes t_shift[8:9] are random, but one of them
has two bits constantly set to 0. Given the random number
generator (RNG) was properly seeded, this would add 14 bit
of entropy to the secret_key. Since the RNG is not seeded,
these two bytes do not add entropy in case the attacker knows
the state of the RNG. Our experiments show that the RNG
state is always the same after rebooting or resetting a robot.
We were able to confirm, that newly unboxed robots have the
same RNG state as well. Note that the RNG state is the same
amongst all robots and not limited to single devices.

The first 4 B of the t_shift structure are the Unix time
in seconds, which is close to the time where the robot
connects to the Beehive cloud for the first time. Since all other
values for the secret_key generation are known, the entropy
consists of timestamp entropy plus 14 bit from the random
values in case the RNG state is not known. If the RNG state
at the time of secret_key computation is known, the en-
tropy exclusively relies on the timestamp. Further security
implications and attacks are discussed in Section 5.1.3.

4 Remote Vulnerabilities

In Section 4.1 several attacks on the robot interface of the
Neato and Vorwerk cloud are presented. Section 4.2 analyzes
a buffer overflow vulnerability in the robot, which is triggered
by sending forged requests to the cloud.
4.1 Robot Backend Vulnerabilities

Obtaining the private RSA keys for the Neato and the Vorwerk cloud (compare Section 3.2) offers a new attack vector against the ecosystems. We give an overview of the message types that are exchanged between robot and cloud and discuss resulting security problems.

4.1.1 Authentication

Robots authenticate against the backend by using RSA signatures in the Authorization header, as detailed in Listing 3. The header consists of two parts, separated by the : character. The first part is the robot_serial and the second part the base64 encoded RSA signature. The signature signs similar content as the Authorization header presented in Section 2.3. It signs the robot_serial, the Hypertext Transfer Protocol (HTTP) request method, the request Uniform Resource Identifier (URI), and the request body. Since some of the backend messages exchange the secret_key discussed in Section 3.3, backend requests do not utilize it and instead entirely rely on the RSA private key.

Robot Impersonation Attack The RSA private key is identical for all robots and the robots do not submit any other authentication factor. This enables an attacker to send messages to the backend on behalf of any robot with a known robot_serial. Problems caused by this issue are explained in the following sections.

4.1.2 Beehive

The Beehive backend is exclusively used for linking a robot to a new account and exchanging the new secret key. Under normal circumstances a robot only has to send a request to the backend once in a lifetime. A linking request contains several information about the robot, among others the robot’s secret_key and the user_id of the account which the robot will be linked to. A linking request as sent by a robot is displayed in Listing 4.

Robot Stealing Attack The Robot Impersonation Attack enables an attacker to send such linking requests on behalf of arbitrary robots. This allows to change the user account to which a robot is linked to. With the means presented in Section 3.3 it is possible to compute the current valid secret_key and fully take over control of a robot.

Changing Keys Possession of the private RSA key also allows for changing the secret_key stored in the cloud for a given robot_serial, which results in the Robot Blinding Attack described in Section 4.1.3.

4.1.3 Nucleo

The Nucleo cloud backend forwards app commands to the robots and receives their status updates. A robot->Nucleo connection is a bidirectional stream, one connection is used for multiple messages. This portion of Nucleo is not HTTP conform. When a robot logs into the cloud with a request to /vendors/neato/login (compare Listing 5), the cloud answers with transfer-encoding: chunked and does not give a Content-Length. It also sends a Connection: close header, which normally ends the connection. If the connection is not closed from client side, the server starts to send
heartbeat messages and control commands. The connection is used as a socket from this point on.

The data format of one message in this socket connection is a length field followed by the Authorization and Date header data explained in Section 2.3. Afterwards the payload is sent. An example exchange is displayed in Listing 6.

Input Validation  The backend is not filtering the payload sent by the app or robot, since the payloads of frontend messages are identical. It is even possible to change data types in the JavaScript Object Notation (JSON), which leads to abnormal behavior in the app. This might represent an attack vector against the app.

Smartphone IP Disclosure  The most interesting functionality provided by the app is manual control, where a user directly controls a robot via websocket. To start manual mode, the smartphone sends a message to the Nucleo cloud, requesting the IP and port where it should connect to. The robot then sends a response with this information and additionally communicates which Service Set Identifier (SSID) it is connected to. The manual driving mode should only work in the same WiFi and local network, as stated by the manufacturer. By impersonating a robot and sending a forged external IP and port, we were able to make the victim’s smartphone connect to our given IP address. Neither the IP range nor the SSID are verified to be local before connecting to the given IP. This poses a severe threat since the victims public IP address is leaked in this attack. With this information, an attacker is able to derive geolocation information or target the victim’s smartphone directly.

Forwarding the unchecked Authorization headers and payloads to the robot has several side-effects:

Robot Blinding Attack  First, changing the secret_key of a robot as described in Section 4.1.2 results in a Denial of Service (DoS) attack. The robot will still receive heartbeat packets and is convinced it is still connected, but it is not able to validate any packets with payload. Therefore the robot cannot be controlled via the cloud anymore.

Secret Key Extraction  Second, in Section 5.1.3 an offline attack against a robot’s secret_key with manual driving packets is explained. Since the messages sent from Nucleo backend to the robots are not additionally encrypted and can be leaked by impersonating the robot, this enables an offline attack as well. The severity of this attack is further increased due to eliminating the constraint of having access to the victim’s local network.

Buffer Overflow  Third, the buffer overflow explained in 4.2 could as well be prevented by not forwarding the header information to the robots unchecked.

4.2 Buffer Overflow

We discovered a buffer overflow vulnerability in the Neato BotVac Connected firmware version 2.2.0 and developed a proof of concept (PoC) exploit, which enables arbitrary unauthenticated remote code execution over the Nucleo cloud with root permissions.

4.2.1 Vulnerability

Commands can be sent to robots over the Nucleo cloud as described in Section 2.3. If the forwarded request body contents are too long, astro crashes locally on the robot and the QNX utility dumper automatically creates a core dump. Listing 7 shows a request that triggers the buffer overflow. It is triggered for arbitrary HMACs, meaning it does not require valid authentication. The only required information is the robot’s serial number.

We analyze the binary files webserver, astro, and the core dump of the latter to locate the root cause. The Link Register (LR) is used to save the return address if a function is called in an ARM program [12]. The return address points to the next instruction to execute after the function call returns.
The LR of the astro core dump belongs to the function HMAC_CTX_cleanup. It is part of the library OpenSSL, but the problem occurs earlier in program execution. We examine the call graph of HMAC_CTX_cleanup shown in Figure 4 to locate the issue, which indicates that astro verifies the HMAC. Note that the function name VerifySignature is misleading here; examination of the binary file reveals that this function calculates a HMAC. Since the crash occurs during HMAC verification, the exploit can be exploited unauthenticated.

Listing 8 shows the part of VerifySignature that contains the buffer overflow vulnerability. The NeatoCrypto.UTF8_Encode function encodes an input string in UTF-8 and stores it in the buffer utf8_str. This is done without verification if the input string fits into the target buffer. The input string str can contain 7265 B, while the target buffer utf8_str is only 2048 B, thus causing the overflow.

The input string is generated by concatenating the robot’s serial number, the date header, and the JSON-formatted commands. The Program Counter (PC) register, which holds the currently executed instruction on ARM [13], retrieves its content from the stack in function epilogues. The corresponding stack value is overwritten by the payload at the position 2032. Moreover, a malicious payload controls registers r4–r11.

```c
Nucleo_message_handler
    VerifySignature
        libNeatoCrypto_VerifySignUsingRobotPassword
            HMAC_CTX_cleanup

Figure 4: Simplified call graph to the function HMAC_CTX_cleanup.
```

VerifySignature(*this, const char *cmd, const char *date, const char *robot_serial, const char *signature) {
    char *str; //str_to_sign
    char *utf8_str;
    // [...] Memset(&utf8_str, 0, 0x800);
    memset(&str, 0, 0x1C61u);
    robot_serial_lower = str_lower_case(str);
    snprintf(&str, 0x1C61u, "%.s%.s%.s\"s", 32, robot_serial_lower, 64, cmd_date, 7168, cmd_str);
    len = strlen(&str);
    NeatoCrypto_UTF8_Encode(&str, len, &utf8_str);
    // [...] }
```

Listing 8: Extract of the source code of the reverse engineered function VerifySignature.

### 4.2.2 Proof of Concept Development

The QNX 6.5 operating system running on Neato BotVac Connected robots supports various exploit mitigation techniques. However, neither Address Space Layout Randomization (ASLR) nor Data Execution Prevention (DEP) are enabled by default [21]. Thus, the buffer overflow exploit development is simplified. Nonetheless, request body payloads must satisfy two conditions. First, a string containing null bytes (0x00) is rejected by the Nucleo cloud. Second, only American Standard Code for Information Interchange (ASCII) characters are accepted (byte values 0x00–0x7f).

To execute arbitrary commands, our proof of concept aims at executing system [4]. Since the astro binary runs as root, commands are executed with the same privileges. astro is using the shared library libc.so.3, which provides the system function at address 0x0101b0a8. However, this address contains non-ASCII bytes. In addition, system requires the command to be passed in r0, which is not directly controlled by the buffer overflows. Therefore, we construct a Return Oriented Programming (ROP) chain containing five gadgets, which are located at valid addresses. An ADD gadget solves the non-ASCII issue by adding up ASCII values to non-ASCII ones, which ultimately leads to root Remote Code Execution (RCE).

### 5 Discussion

In this section, the most important findings and their security implications and impact are discussed. We continue using the same structure as in the actual findings and first discuss local vulnerabilities in Section 5.1, remote vulnerabilities in Section 5.2, and then discuss remaining issues in Section 5.3.

#### 5.1 Local Vulnerabilities

In this section, security implications related to findings in Section 3 are discussed.

##### 5.1.1 Secure Boot Bypass

Fixing the bootloader has a high potential in brickering robots. If anything fails during the bootloader update, there is probably no recovery for customers. However, a hard-coded password and a memory reset would be sufficient to stop an attack at this stage in the first place. Neato’s solution to this issue rolled out in version 4.4.0-72. It fixes this issue by detecting if the robot is in factory mode and skips the boot menu if not.

##### 5.1.2 Static Firmware Analysis

On evaluating the binaries, we found that they are not obfuscated at all—they even contain symbols which made analysis fast and efficient. The efforts required to fully understand and exploit an ecosystem could be significantly increased by obfuscation, while the costs for doing so are fairly low.
5.1.3 Secret Key Generation

Secret keys are computed based on robot MAC addresses, a timestamp, and an unseeded random element. In this section we discuss the security of serial numbers and the resulting implications on the secret key.

Serial Number Security A lot of the security and authentication depends on a robot’s serial number—which turns out to be fairly easy to find out. First, serial numbers are printed on the robot’s packaging and the robot itself. The former enables attackers to go into shops and note their numbers or checking someone’s trash. The latter is often photographed when selling used robots. Second, serial numbers can be enumerated as shown in Section 4.1.2. The only information they contain is the MAC address and manufacturing date, with the former being observable in—for instance—the network and the latter being within a limited time range. Taking this into consideration, a robot’s serial number is not a secret and security solely relies on secret_key security.

Secret Key Security Considering the findings in Section 3.3, where we presented the NeatoCrypto_GenerateRobotPassword, several derivations are made. If an attacker knows in which year a robot was first connected to the cloud, the timestamp contributes \( \log_2(365 \times 24 \times 60 \times 60) \approx 25 \) bit of entropy. In contrast, if he is able to tell the timestamp with an accuracy of \( \pm \frac{1}{2} \mathrm{h} \), the number of unknown bits is \( \log_2(60 \times 60) = 12 \). Table 2 gives an overview of different timespans and the resulting entropy and average number of tries needed for bruteforcing a secret_key, depending on a known respectively unknown RNG state.

As visible in Table 2, an attacker is able to test all possible secret_keys for a given robot_serial against the cloud, if the RNG state is known and the time of first connection can be determined with a 30 min accuracy. If the time is not known to a sufficient accuracy, it is still possible to execute an online attack distributed over a longer period of time.

For an unknown RNG state, testing against the Nucleo API, since knowledge of the secret_key and the robot_serial is sufficient to forge a valid authentication header. This is an upgrade to an attack presented in [19], since for the new attack the robot is not paired with a new account. Instead the old account connection stays intact and the attack is thereby very hard to detect.

Since this is found in the NeatoCrypto library, this attack is valid for the most recent Neato and Vorwerk vacuum robots as well. With the secret_key being the root of trust in their cloud environment, this has further security implications.

5.2 Remote Vulnerabilities

In this section the security implications of the findings presented in Section 4 are discussed.

5.2.1 Robot Backend Vulnerabilities

We revealed several problems in the backend and the common root of problem is the usage of an identical RSA private key on all devices. Having different devices using the same private key defeats the purpose of public key cryptography—a symmetric cryptography scheme would be sufficient. Since RSA consumes a lot of computation time, this also allows to send lots of random Authorization headers to the backend, which results in an amplified Denial of Service (DoS) attack.

With an attacker being able to send and receive arbitrary data from both the app and the robot cloud interface, the server could be exploited as a botnet command and control server. An attacker could register multiple fake robot vacuums with arbitrary robot_serials and use them as an interfacing point for control data. Since knowledge of a robot_serial is enough to communicate with a fake robot, the app interface would be used by the botnet’s bots in this scenario. An attacker could also register high amounts of fake robots sending regular traffic to execute a stealthy Distributed Denial of Service (DDoS) attack. This would cause a high message load on the backend that is almost impossible to be filtered.

<table>
<thead>
<tr>
<th>RNG state</th>
<th>Timespan</th>
<th>Entropy</th>
<th>Avg. # tries</th>
<th>Bruteforce</th>
</tr>
</thead>
<tbody>
<tr>
<td>known</td>
<td>±30 min</td>
<td>12 bit</td>
<td>2048</td>
<td>0.00015 s</td>
</tr>
<tr>
<td></td>
<td>±1/2 year</td>
<td>25 bit</td>
<td>16,777,216</td>
<td>1.1 s</td>
</tr>
<tr>
<td>unknown</td>
<td>±30 min</td>
<td>26 bit</td>
<td>33,534,432</td>
<td>2.2 s</td>
</tr>
<tr>
<td></td>
<td>±1/2 year</td>
<td>39 bit</td>
<td>2.75 \times 10^{11}</td>
<td>5 h</td>
</tr>
</tbody>
</table>

Table 2: Complexity computations for secret_key.
5.2.2 Remote Command Execution

An attacker could use the command execution on the robot to extract sensitive information from the robot like maps and other sensor readings. This violates the privacy of the victim. An attacker gains access to the victim’s local network via the vulnerability. The access can be used to attack further devices inside the local network which would be protected by the firewall of the network otherwise. The attacker can also use the exploit to execute a DoS attack against the robot if he executes a command which bricks the robot. Last but not least the exploit could be used to build a botnet consisting of Neato robots, which can be utilized for DDoS attacks.

5.3 General

We finalize the discussion with general security-related observations.

Update Process The update process for the older connected series is not user-friendly. Updating the Neato BotVac Connected requires users to download firmware to a correctly formatted USB stick which has to be then plugged in with a special cable. Finally, the update process has to be invoked manually. We assume most customers will not update their robots, unless they are provided with a over-the-air (OTA) update functionality, as it is featured by the most recent models. Customers could be forced to do updates by locking their robots, however this might cause angry customer feedback.

Responsible Disclosure Establishing contact with Neato turned out to be more of an obstacle than expected. After first contacting them about our findings on March 12, 2018, customer support repeatedly claimed to have forwarded our messages to development, which was later disputed. Even sending a PoC Python script for the cloud-based unauthenticated command execution was ignored. In fact, it was not until November 27, 2018 when we were finally approached by Neato core development. Afterwards, the issues were fixed in a timely fashion and Neato even provided us with the latest Neato BotVac Connected D7 for further testing.

6 Conclusion

In this paper we presented multiple attacks on the Neato and Vorwerk ecosystems, which ultimately lead to various possible security breaches. Amongst them are those to start and pause the robot, extract a map of the victim’s apartment, read arbitrary sensor data, leak customer public IP addresses, and even gain access to the local network of a customer.

On a first glance, Neato did a very good job in securing their connected vacuum cleaners: secure boot and mutual authentication between robots and the cloud exceeds what is done for most IoT products. Their robots are in the upper price range and the customer receives premium security features in return. Dissecting their robots was not only challenging, but also now provides interesting insights in building secure IoT applications and other large-scale connected scenarios.

However, Neato followed a hard shell, soft interior approach. Once an attacker is inside the system with access to all binaries, many security best practices were missing: no privilege separation, no operating system security features, no binary obfuscation, and hand-written cryptography. Especially the cryptography part looks well-done from the outside and represents solid security without knowledge of the exact parameters used in the computation. Their implementation illustrates how scientifically proven strong theoretical security primitives still fail in practice. The mistakes were present over multiple years without being noticed by the developers, thus indicating a need for better books and standards to be guided.

Our analysis shows that cryptography on current devices is flawed and does not provide sufficient security. Neato will need to strictly filter invalid requests in the cloud to prevent attacks, which is not possible for targeted attacks on individual robots. To maintain security in a targeted attack scenario, users need to protect their robot’s serial_id by any means: they should buy robots exclusively from the vendor, not throw away the original packaging (or shred it), use a separate WiFi, and prevent physical access to the robot by any third parties.

From a customer perspective, more research on which data is shared depending on the vacuum cleaner model would be interesting. Collected and accessible data contains a lot of information and meta data about the customer’s everyday life. This includes the location of residence, apartment size and furniture, number of inhabitants including pets, unencrypted traffic exchange on the robot’s WiFi, daily scheduling, and more. This information is sensible enough to be worth all the protection mechanisms discussed in this paper.

Finally, we want to encourage the use of the knowledge gathered in this paper by the existing Neato and Vorwerk communities to build smartphone apps and gateways with additional features and capabilities.

Acknowledgments

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References


