

Demystifying Pointer Authentication on Apple M1

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Outline

- Pointer Authentication
 - What is Pointer Authentication (PAC)
 - Current Research State of Apple PAC
- Our Motivation
 - Lack of systematic analysis of Apple PAC Protection
- > Our Contribution
 - m1n1-based reverse engineering framework
 - Disclosure of Apple's hardware implementation
 - Comprehensive analysis of PA-based XNU kernel protection
- > Our Findings
 - Apple's PA Hardware
 - PA-based XNU Kernel Protection
 - Security analysis

(PAC) le PAC



1 control register

- SCTLR_EL1.EnIA/IB/DA/DB)



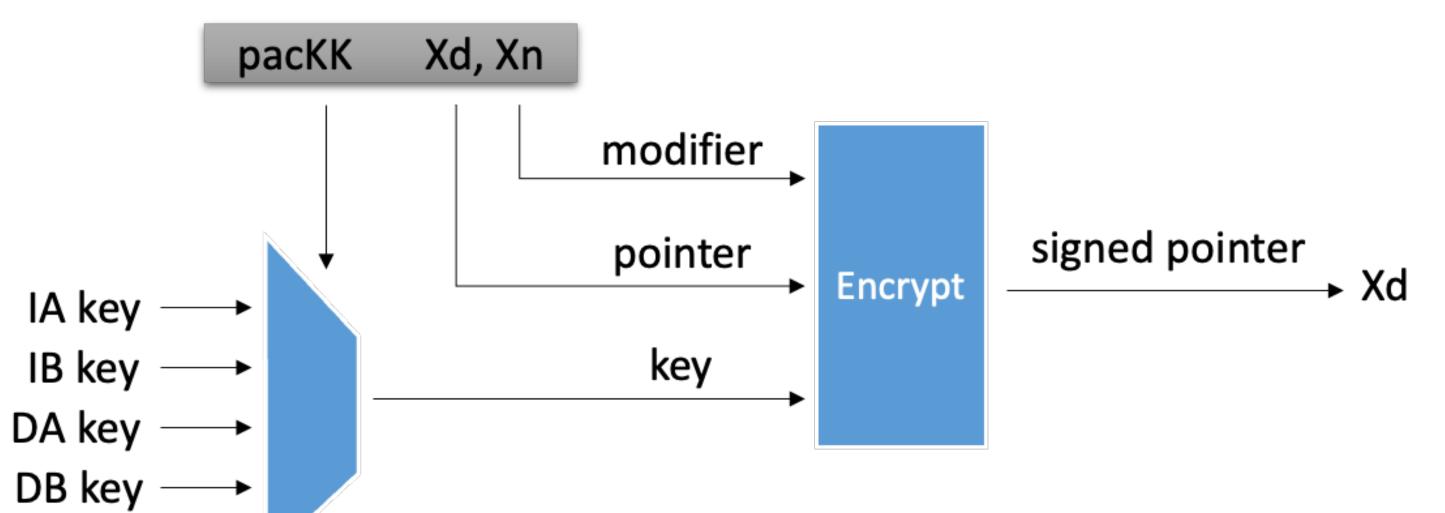
- APIA/APIB/APDA/APDB/APGA
- 2 kinds of instructions
 - pac*/aut*

0xfffff00010c00ac4 Pointer

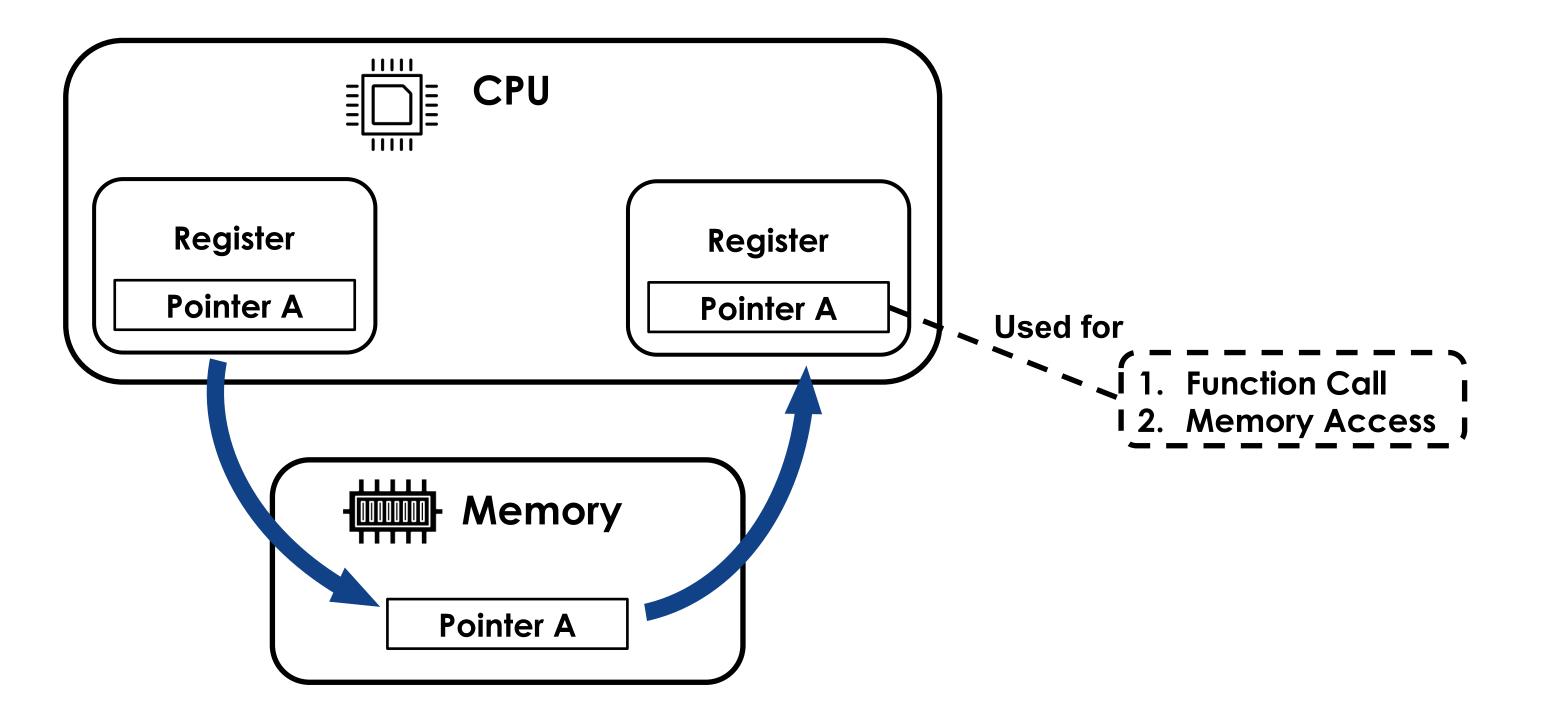
IA key

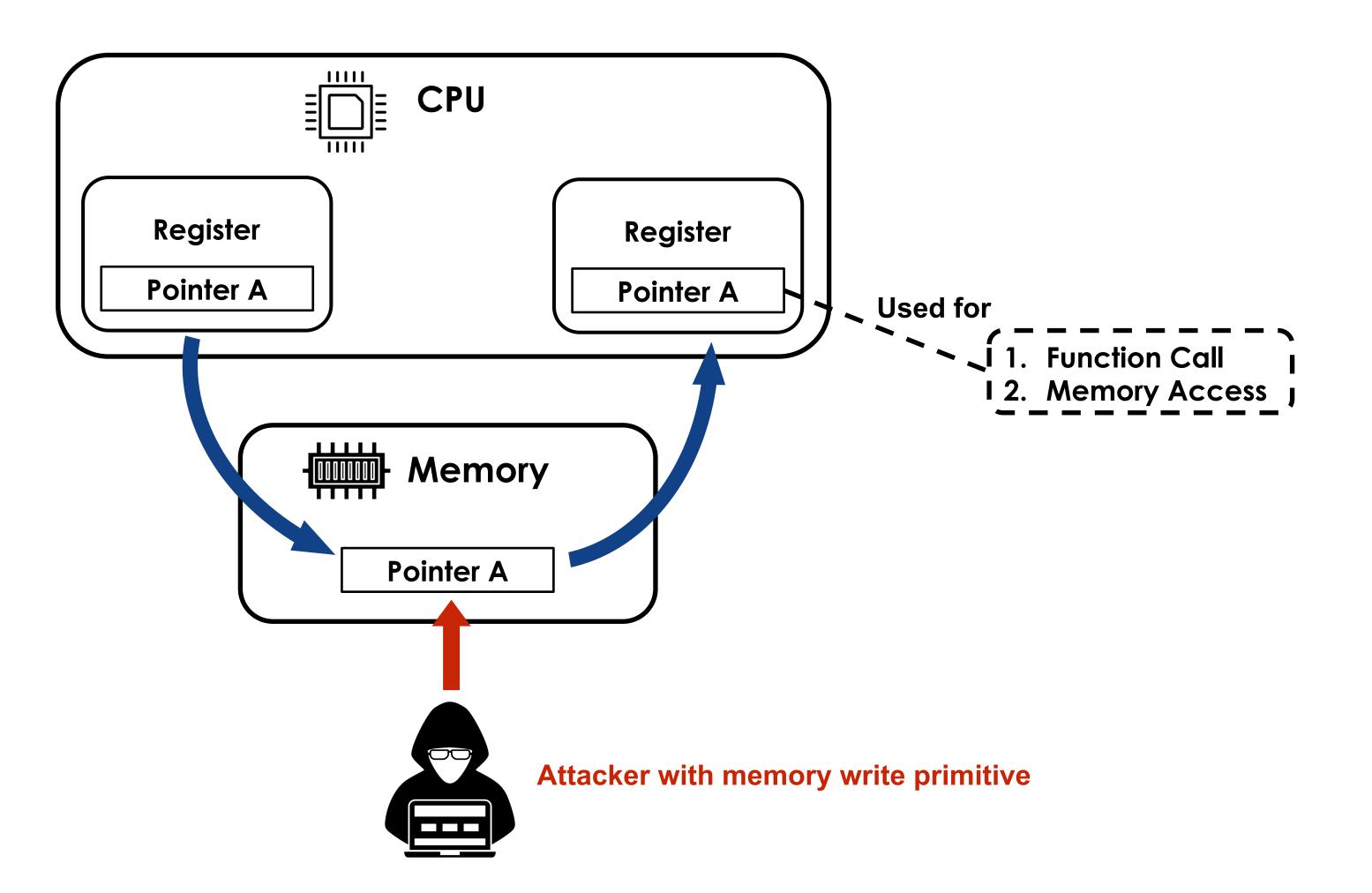
IB key -

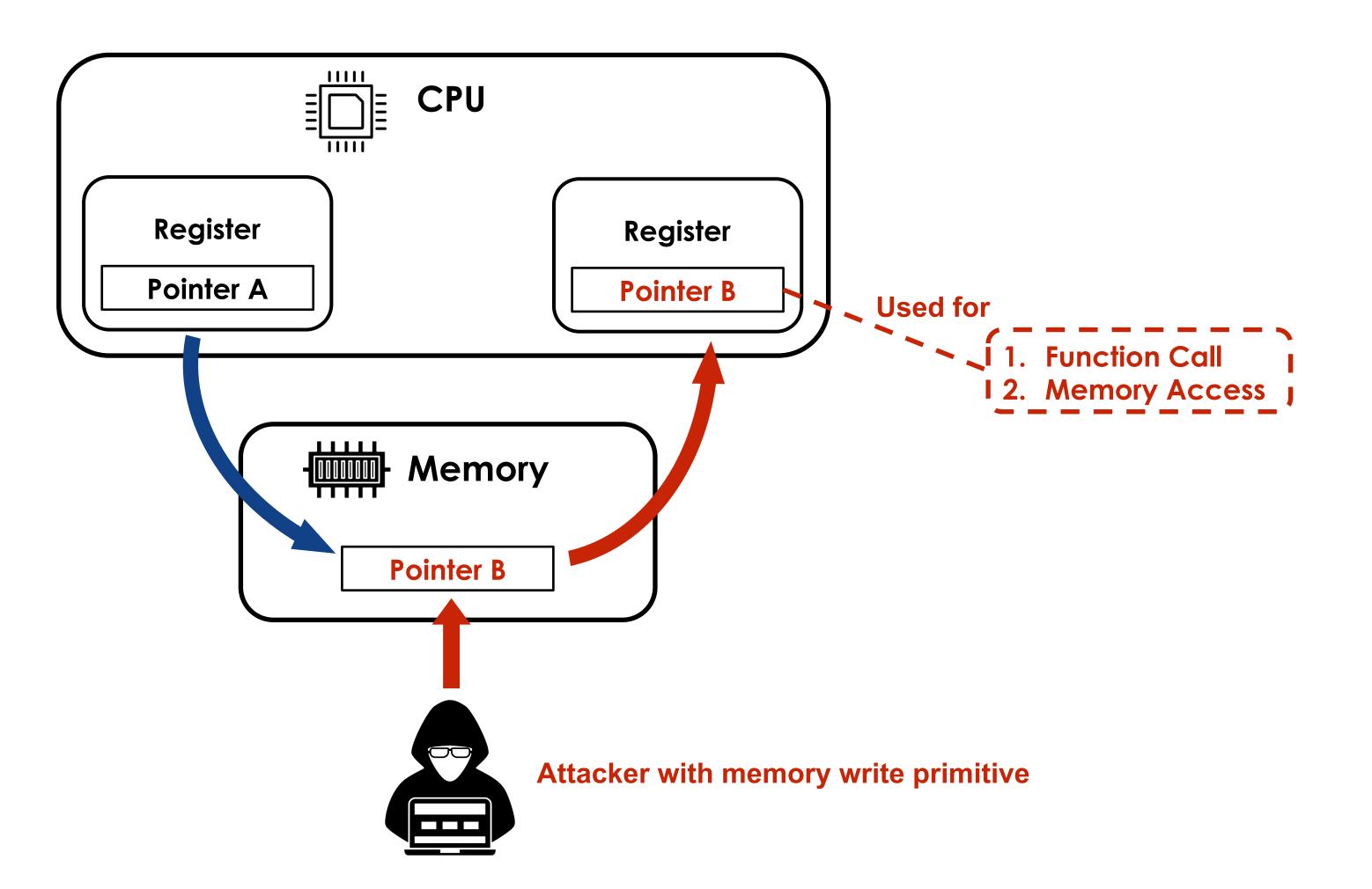
DB key

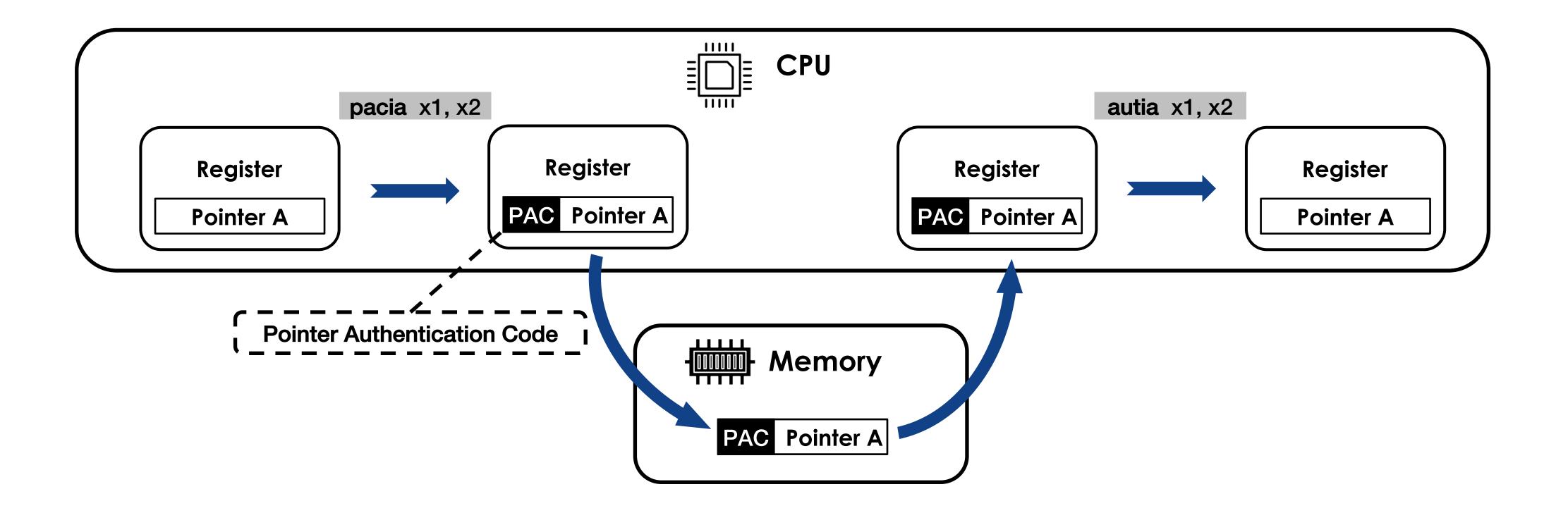


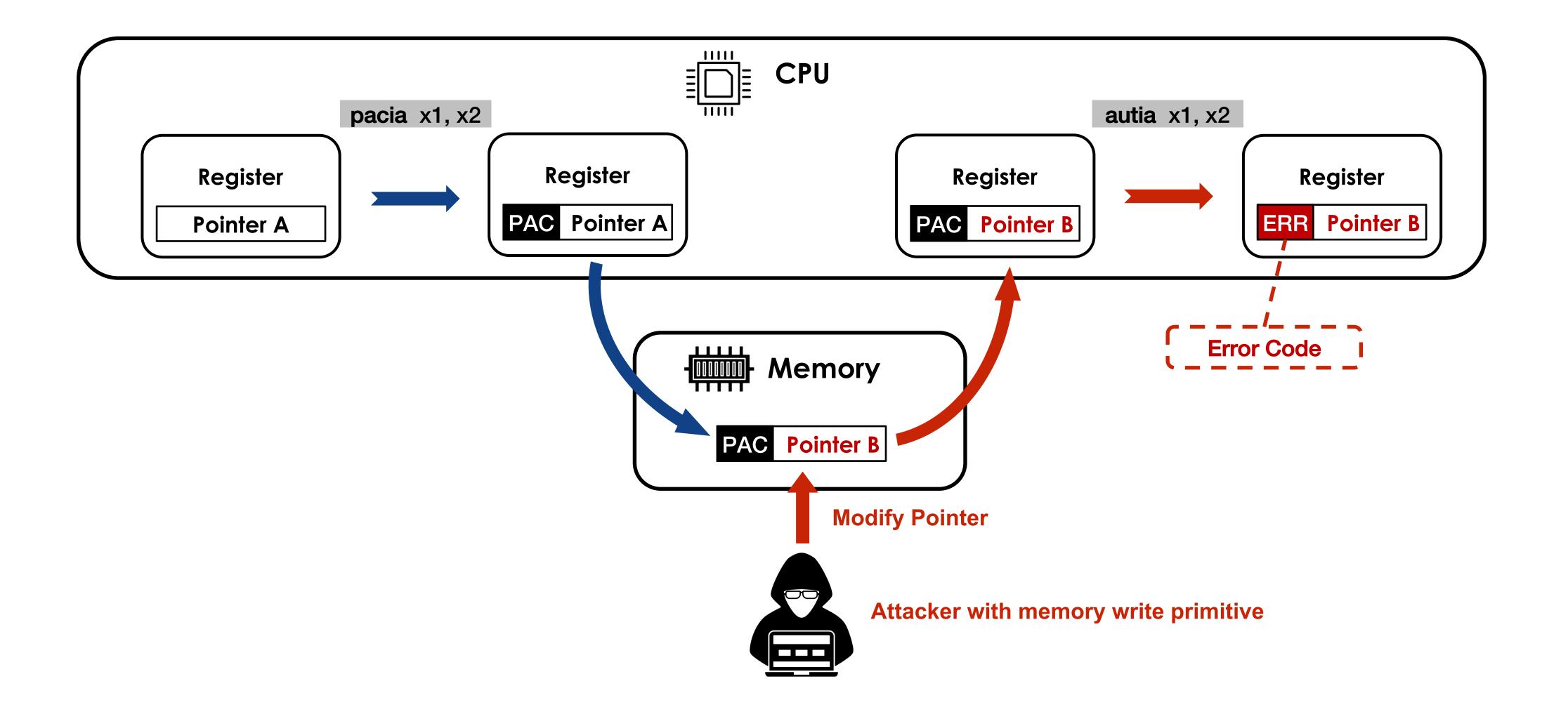


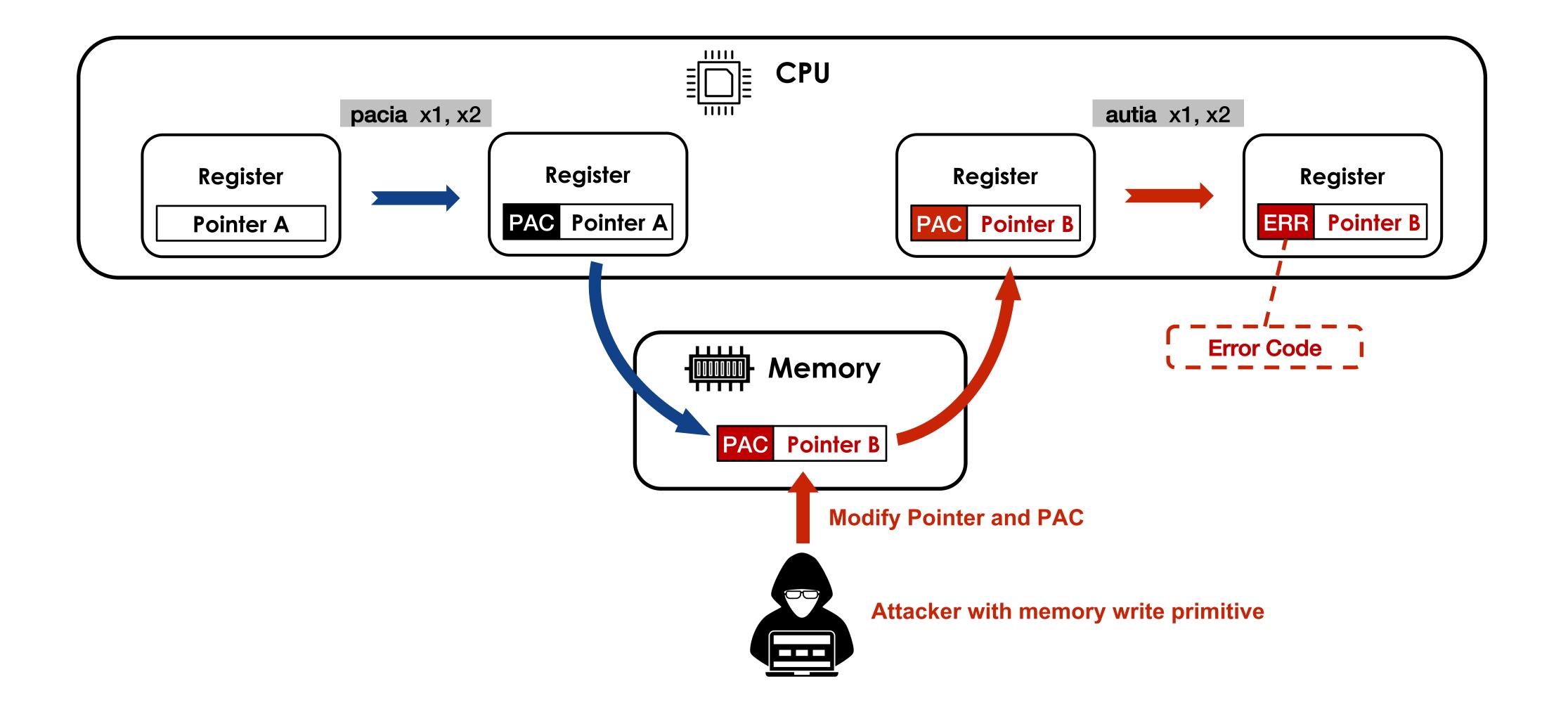








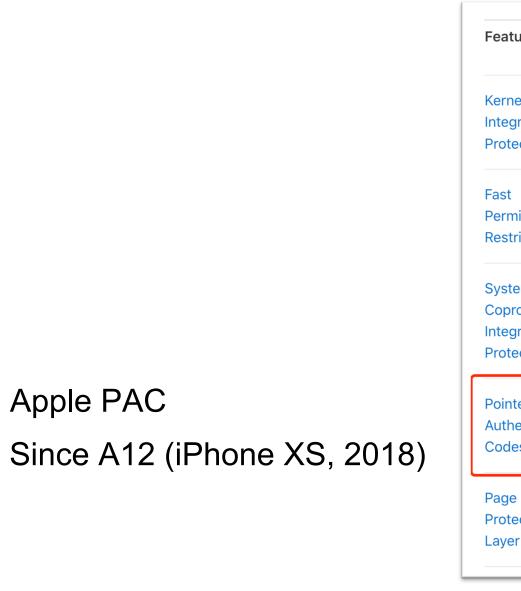




Apple's PA Protection

- Apple is the first one to implement and deploy PA hardware
 - State-of-the-art mitigation against pointer corruption attack
 - Since its debut, the number of public kernel exploits has decreased

Apple PAC



[1] Apple SoC Security. https://support.apple.com/guide/security/apple-soc-security-sec87716a080/web

ature	A10	A11, S3	A12, S4	A13, S5	A14, A15, S6, S7	M1 Family
rnel egrity otection	0	•	0	Ø	0	
st rmission strictions		Ø	0	Ø	Ø	Ø
stem processor egrity otection			0	Ø	Ø	Ø
inter thentication odes			٢	Ø	 Image: A start of the start of	 Image: A start of the start of
ge otection yer		Ø	0	Ø	Ø	See Note below

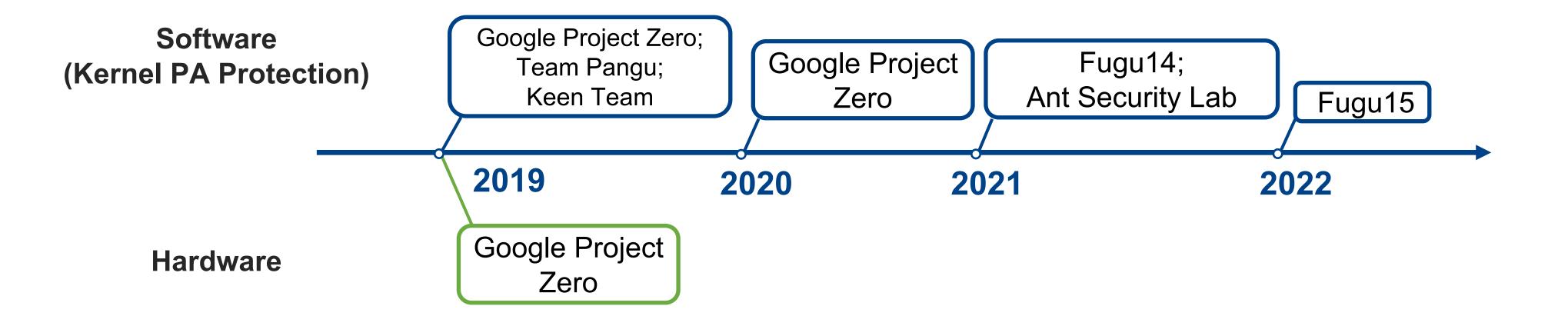
Apple SoC Security^[1]



Our Motivation

Although there are a lot of research works on of Apple's PA \bullet

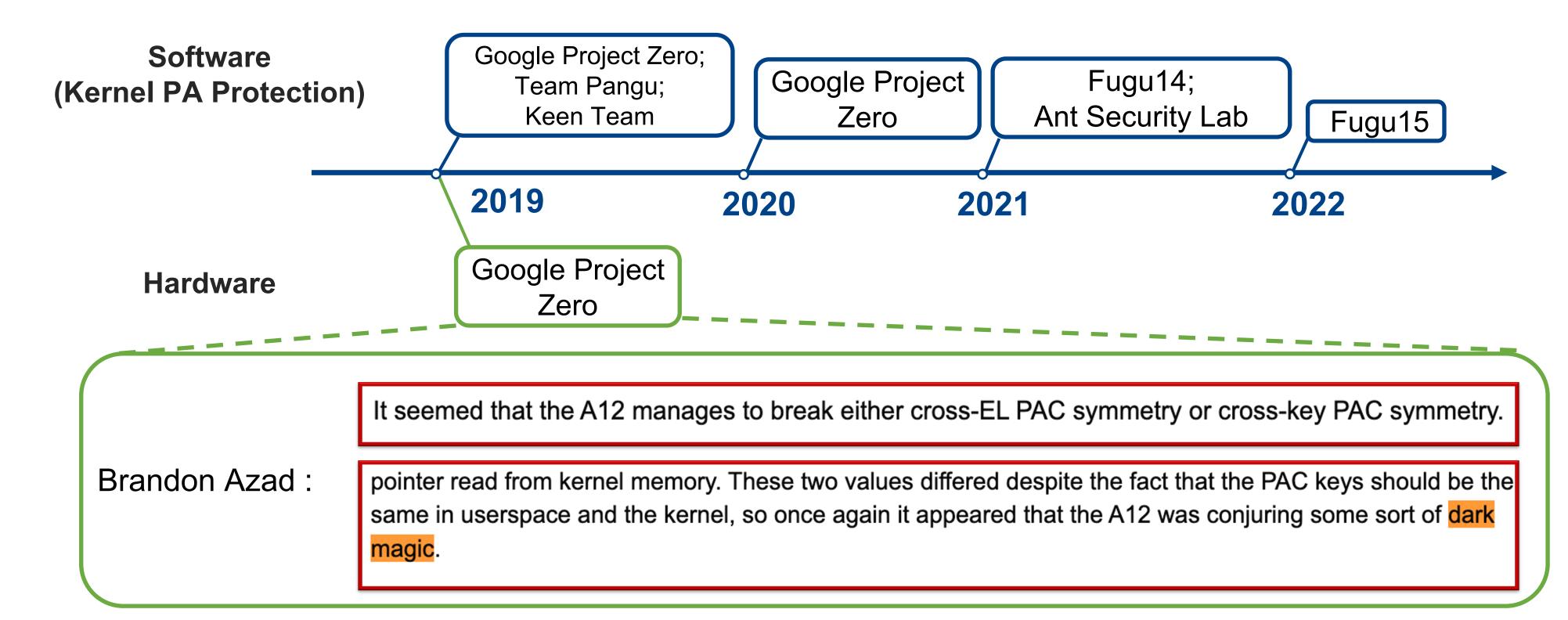
- Most of them focus on PA-based software protection
- How does Apple implements Apple PA hardware remain unknown





Our Motivation

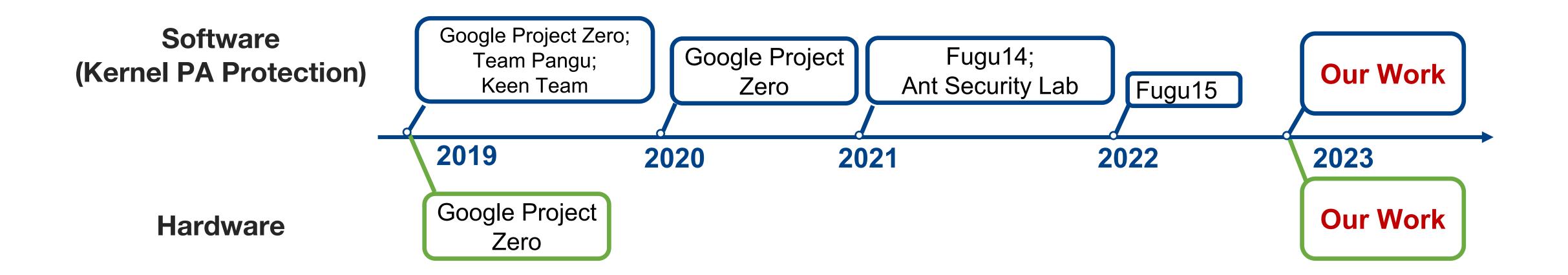
- Although there are a lot of research works on of Apple's PA
 - Most of them focus on PA-based **<u>software</u>** protection -
 - How does Apple implements Apple PA hardware remain unknown -



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Our Motivation

- The hardware implementation remains unknown
- As a result, the PA software can not be analyzed systematically
- There is still no systematic analysis of Apple PA hardware and software





Our Contribution

- to analyze PA hardware and software on Apple M1
- properties Cross-domain (Cross-EL/Key/VM/Boot) attack mitigation
- update and assigned us a CVE.

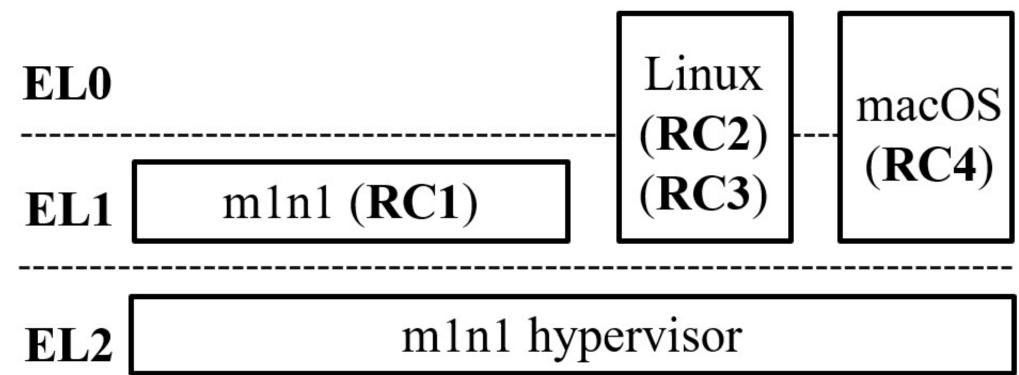
 \succ We build a reverse engineering framework based on m1n1 (an open-sourced hypervisor)

We reveal how Apple customizes the PA hardware to introduce undisclosed security

> We analyze the implementation of PA-based XNU kernel protection and identify four attack surfaces. Apple acknowledged our findings publicly, fixed these issues in a security



- Required Capabilities (RC)
 - **RC1: Identify** undisclosed PA-related Apple-specific system registers
 - **RC2: Read/Write** actual PA key values
 - **RC3: Profile** the undisclosed PA instruction behavior
 - **RC4: Debug** the XNU kernel dynamically



m1n1-based RE Framework



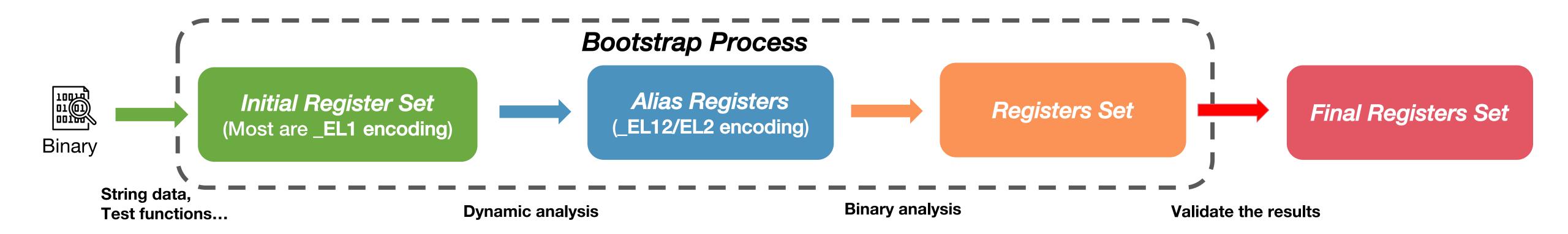
m1n1-based RE Framework

- RC1: Identify undisclosed PA-related Apple-specific system registers
 - Challenge: Apple introduced a lot of undisclosed system registers

Our solution: We identify registers based on

- **Binary** information

- System Register Redirection Hardware Feature





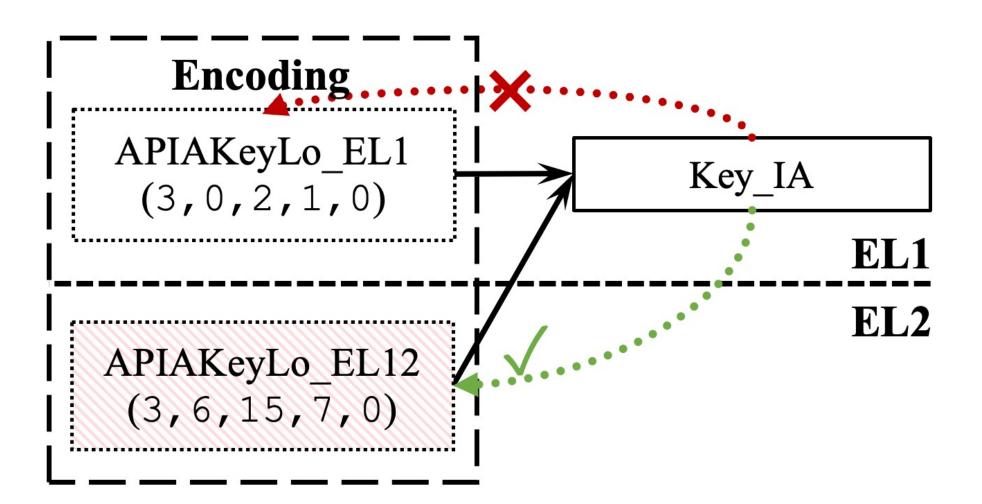
RC2: Read/Write the actual key value

Challenge: Apple implemented a hardware-based PAC key protection

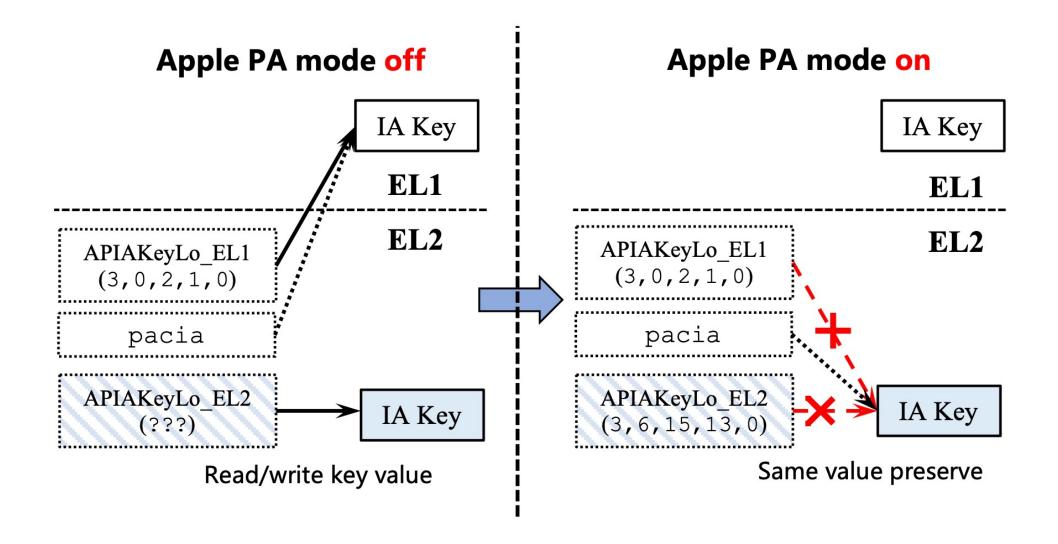
Our techniques:

- For EL1 Key, Read/Write the key from EL2 exception level

- For EL2 Key, Read/Write the key before Apple PA is enabled



m1n1-based RE Framework



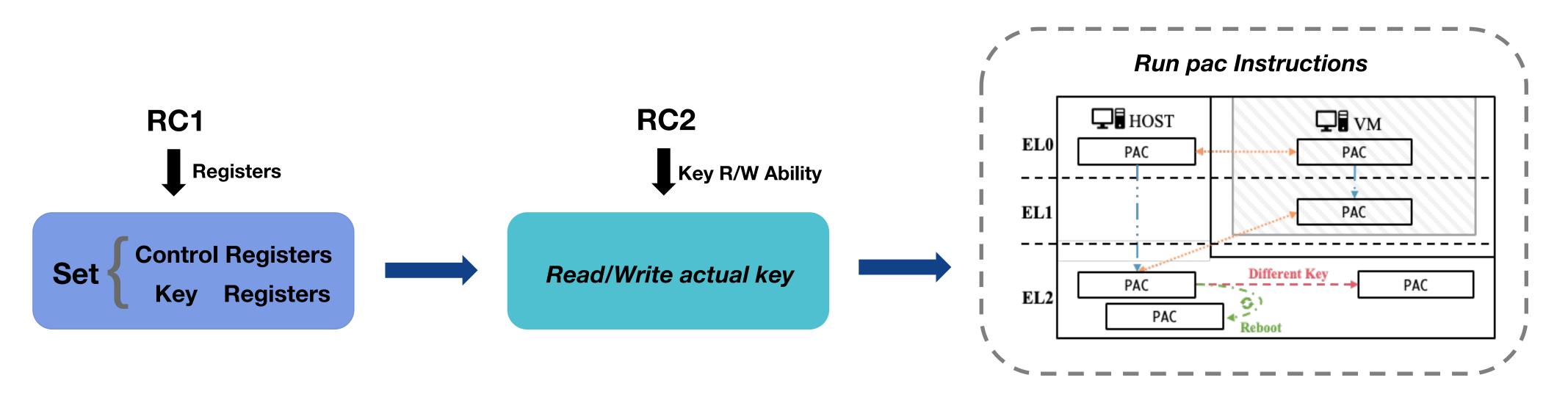


m1n1-based RE Framework

RC3: Profile the undisclosed PA instruction behavior

Challenge: We need to analyze the complex interplays between registers and intructions

Our solution:



Step 1

Step 2

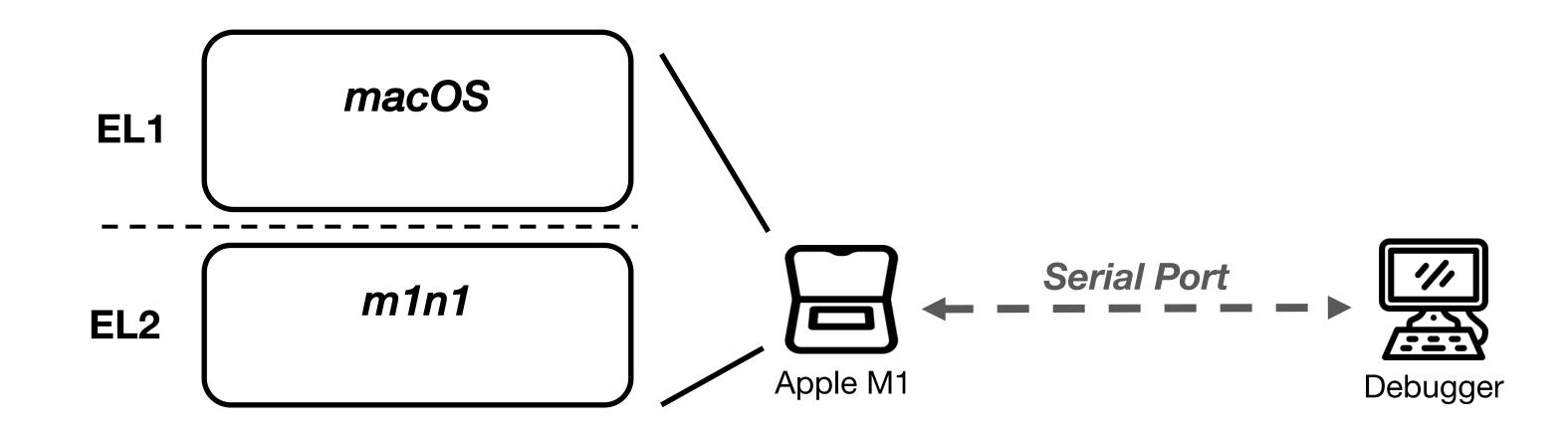




RC4: Debug the XNU kernel dynamically

Challenge: LLDB (provided by Apple) does not support active kernel debugging on Apple M1

Our solution: We implement active kernel debugging based m1n1 hypervisor



m1n1-based RE Framework

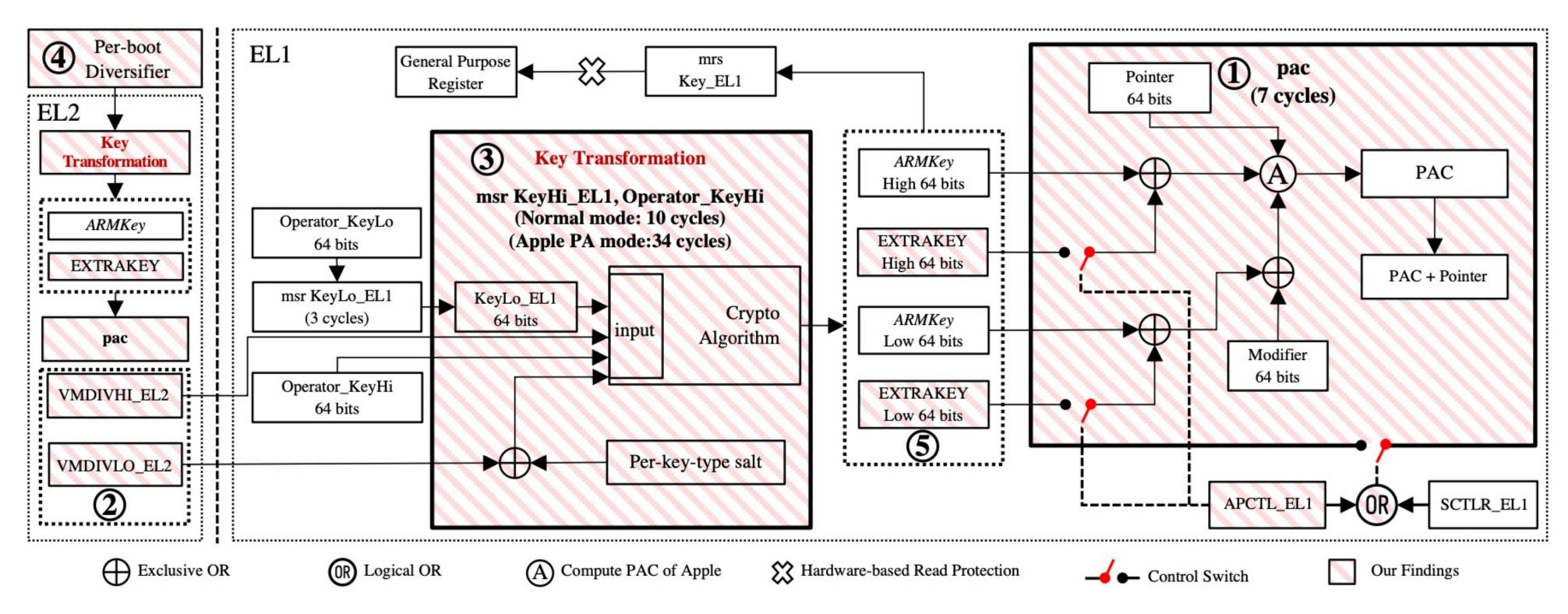
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Our Findings

• Finding Overview

1 Controllability, PAC algorithm

2-5 Cross-domain Attack Mitigation



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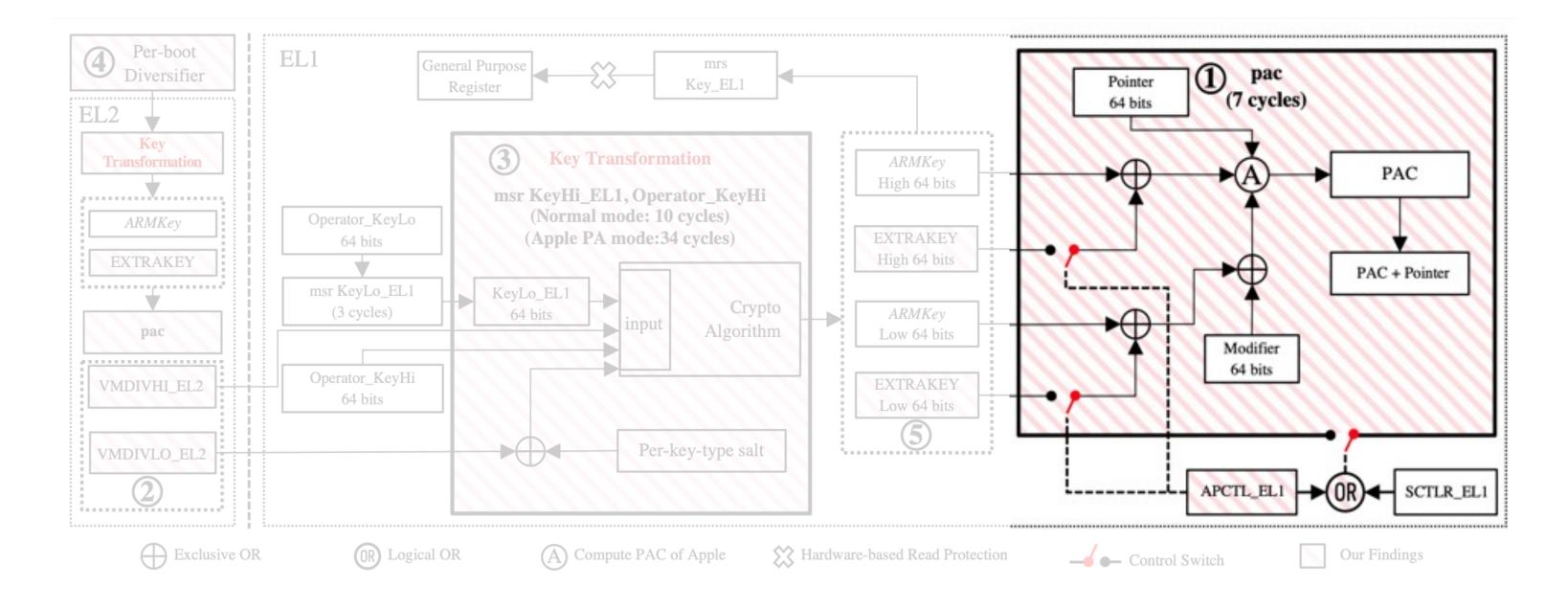
• Controllability & PAC Algorithm (1)

APCTL_EL1

- bit[0]: Enable Apple PA

-bit[2], bit[3]: Enable PA on user (bit[2]) or kernel space (bit[3])

- bit[1], bit[4]: Enable EXTRAKEY on user (bit[4]) or kernel space (bit[1])



PAC Algorithm is not QARMA

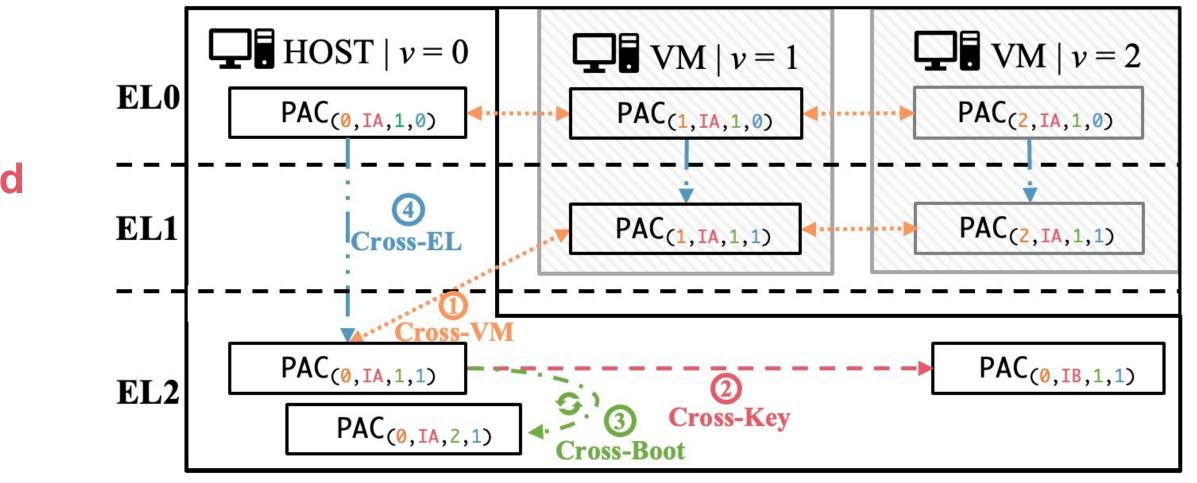
- (Modifier XOR KeyValue) is one of the inputs



- **Cross-domain Attack**
 - Pointer substitution attack across different domains

Cross-VM: From VM to Host and Other VMs Cross-Key: E.g., From APIA-signed to APIB-signed Cross-Boot: From Boot Round 1 to Round 2 Cross-EL: From User space to Kernel space

Formalization of Cross-domain Attack in the paper 🤐



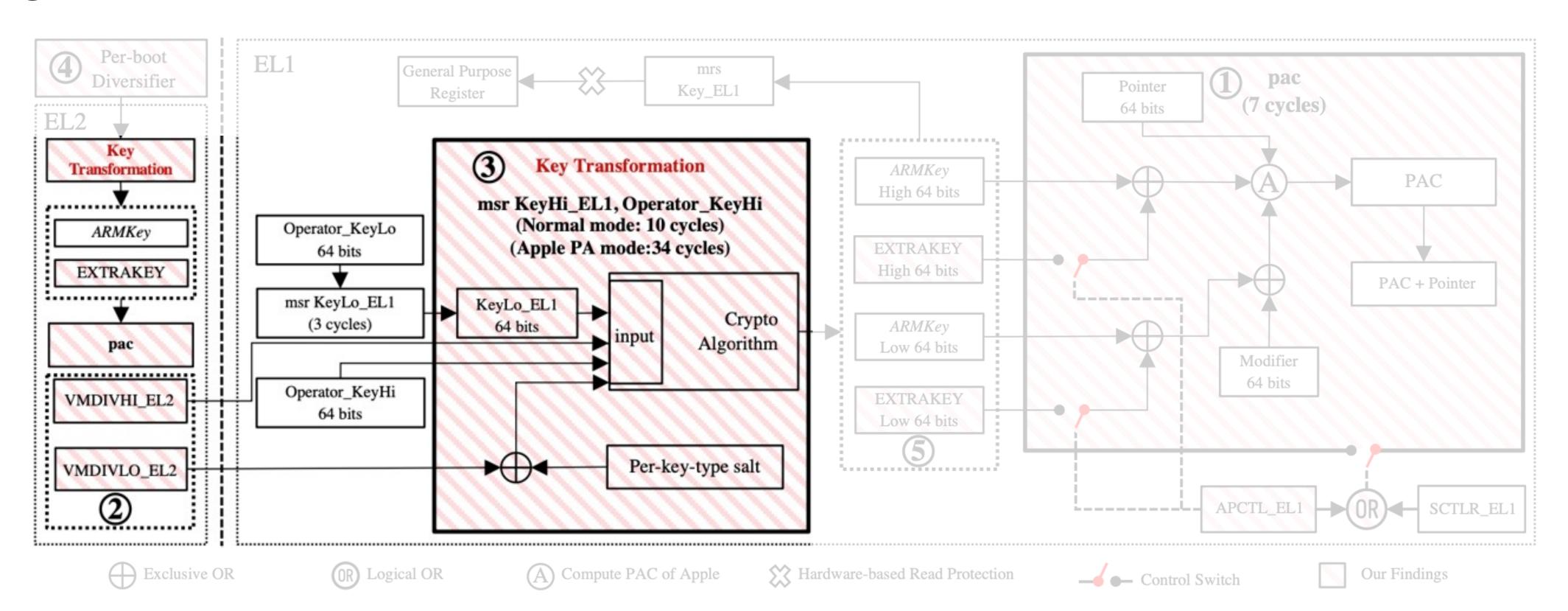




Cross-VM attack mitigation

Setting the higher 64-bit PAC Key will trigger a Key Transformation

(2) VMDIV_EL2 is used for differentiate the Key Transformation between VM and Host



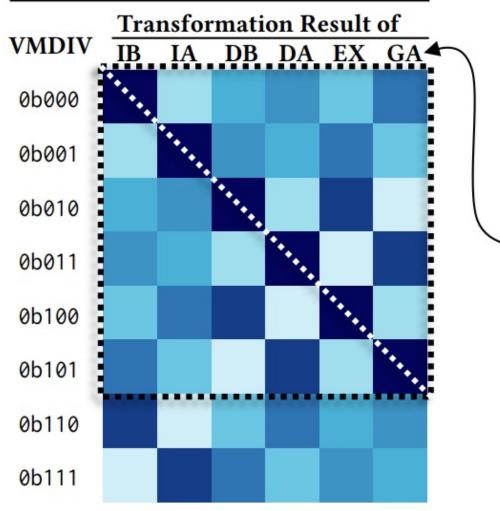


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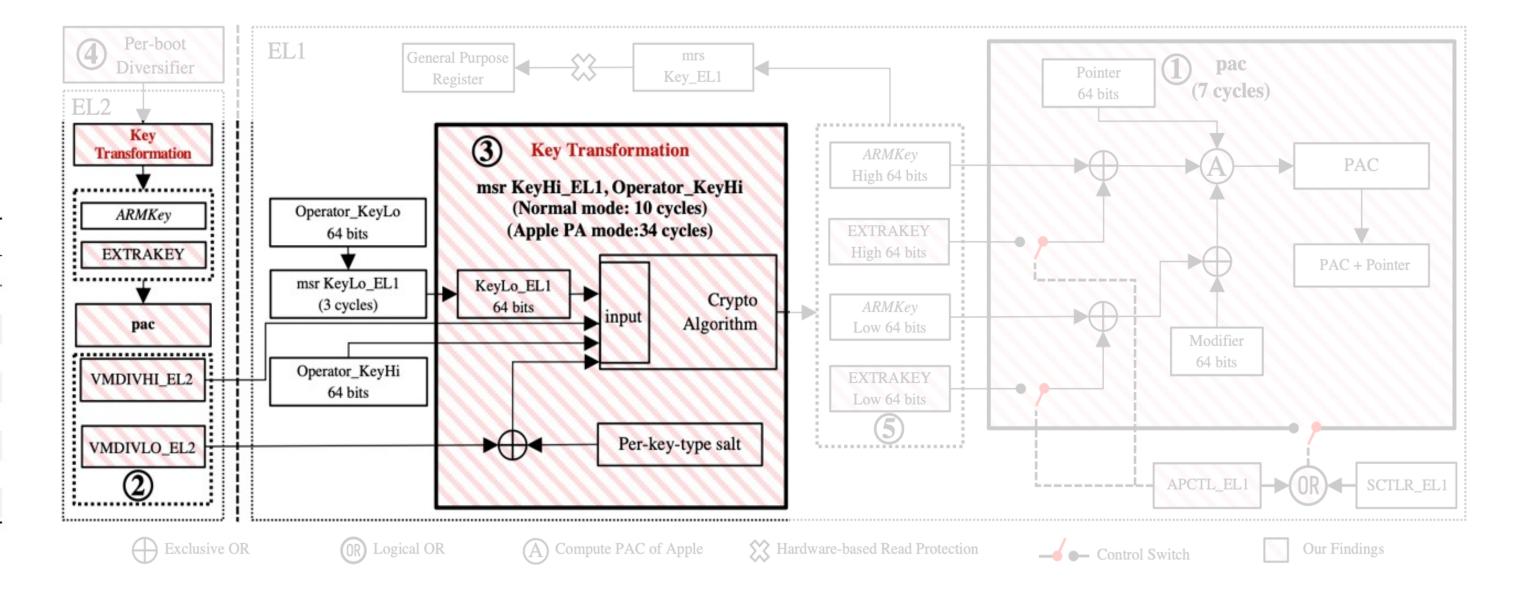
Cross-Key attack mitigation

Key Transformation introduces per-key-type salts to differentiate the results for different key types

3 VMDIV_EL2 • per-key-type salts is one of the inputs for Key Transformation

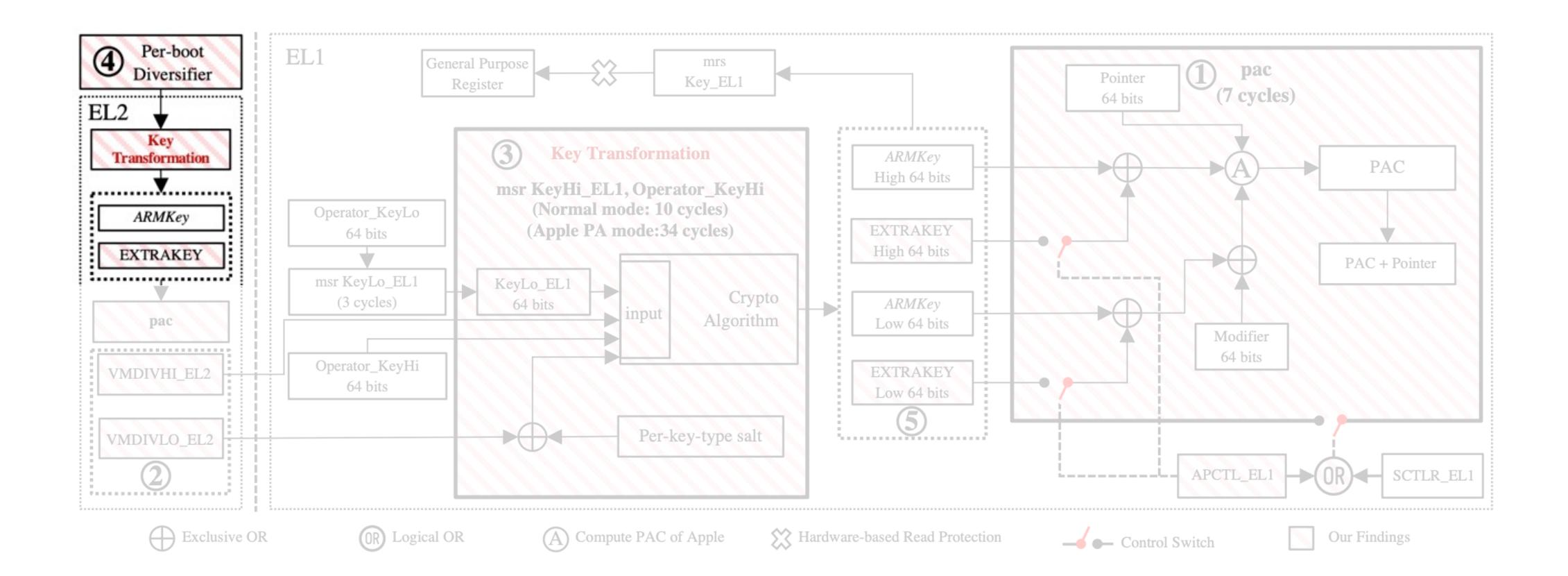


-	Per-key-type Salt of						
-	IB	IA	DB	DA	EX	GA	
	0	1	2	3	4	5	
	1	0	2	3	4	5	
	2	3	0	1	6	7	
	3	2	1	0	7	6	
	4	5	6	7	0	1	
	5	4	7	6	1	0	
	6	7	4	5	2	3	
	7	6	5	4	3	2	





Cross-Boot attack mitigation



EL2 Key Transformation introduces (4) per-boot diversifier to differentiate the results for different CPU Boots

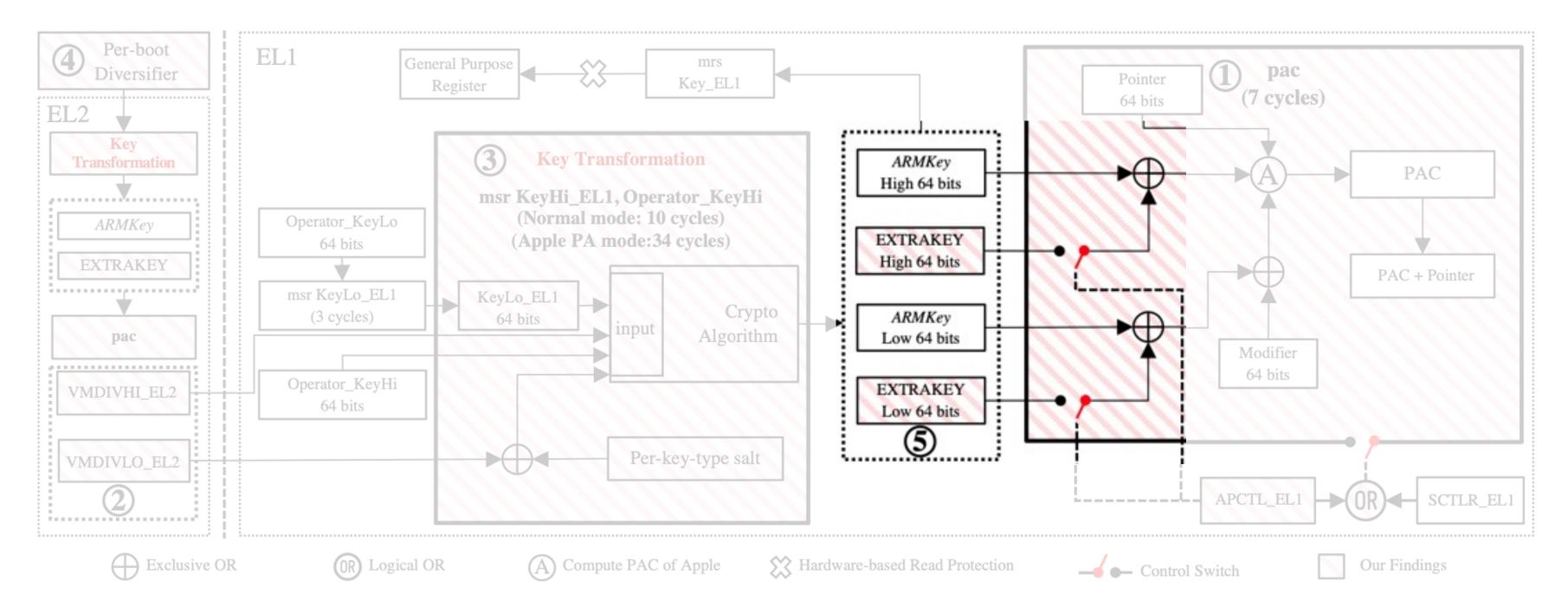


Cross-EL attack mitigation

Apple introduces an EXTRAKEY to differentiate the PAC computation between user and kernel space

(5) EXTRAKEY \bigoplus APKeys (APIA/IB/DA/DB/GA) is the actual key value for PAC computation

Controlled by APCTL_EL1 (bit[1]: Kernel, bit[4]: User (XNU Kernel only enable bit[4]))





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PA-based Kernel Protection

- Sign/Auth Interfaces Analysis
 - We analyze all pac instructions in XNU kernel
 - The result shows that XNU kernel uses
 - 9 types of signing modifiers
 - (5 types in official documentation)
 - 6 policies for generating modifier constant
 - (2 policies in official documentation)

Key Usage	Modifier	Target	
	Hash(function_type) (38564)	Function pointer	
	Hash(function_name) + storage address(384)	Recovery Handler/ Corecrypto related	
		ppl_handler(89)	
IA (762376)	Storage address (27568)	Copy/destroy_helper_block (626)	
		Block_invoke function(115)	
		Ptrs in seg:auth_ptr (25452)	
		Block function pointer (1286)	
	Hash(root_class, function_type, function_name) + Storage address (694596)	Vtable entry(694596)	
	Zero (1264)	chkstk_darwin* func (156)/ BluetoothFamily function (9)/ kext_weak_symbol_referenced/ Parameter func ptr (1098)	
		Proc0 (1)	
	Hash(data_field_type) + Stor-	NSConcreteGlobalBlock(11:	
DA (31225)	age addr (2645)	_Block_descriptor (115)	
DA (31223)		sysctl_oid_list (2414)	
	Hash(data_pointer_name)+ Storage addr (464)	Data pointer	
	Hash(root_class) + Storage addr (28116)	V-table pointer	
IB(110852)	SP	Return address	
DB	-	-	
GA	Storage Address	Thread state/ Exception state/ Data Blob	



PA-based Kernel Protection

- Key Management
 - XNU kernel configures the keys
 - Global: APIA/DA/GA
 - Per-Process: APIB/DB, EXTRAKEY

Key	APIA	APDA	APGA	APIB	APDB	EXTRAKEY
Scope	Global	Global	Global	Per-Process	Per-Process	Per-Process



PA-based Kernel Protection

- Key Management
 - PAC instruction scope
 - pacia/da/ga: global in kernel space, per-process in user space
 - pacib/db: per-process
 - For non-arm64e process, the XNU kernel disable the user space PAC

PAC instructions	pacia	pacda	pacga	pacib	pacdb
User (arm64e)	Per-Process	Per-Process	Per-Process	Per-Process	Per-Process
User (Non-arm64e)	-	_	Per-Process	Per-Process	_
Kernel	Global	Global	Global	Per-Process	Per-Process



- We validate 4 attack surfaces (88 cases) and report them to Apple
 - **Incomplete Sensitive Data Identification** (1)

- Potential Enhancements

(2)**Incomplete Interrupt Context Protection**

- Fixed in a security update and acknowledged publicly

- (3)Signing Gadget
 - Fixed and assigned a CVE (CVE-2023-32424)



- Potential Enhancements

More detail about identification and validation in the paper $\stackrel{(())}{\Rightarrow}$

Security Analysis

Result AS	Identified	Validated	Fixed Cases
0	153	83	6
2	$17 + 18^{*}$	2	2+18*
3	1	1	1
4	2	2	-

18 cases are identified in XNU-7195, and all of them are fixed in XNU-8019.





