Pudica: Toward Near-Zero Queuing Delay in Congestion Control for Cloud Gaming

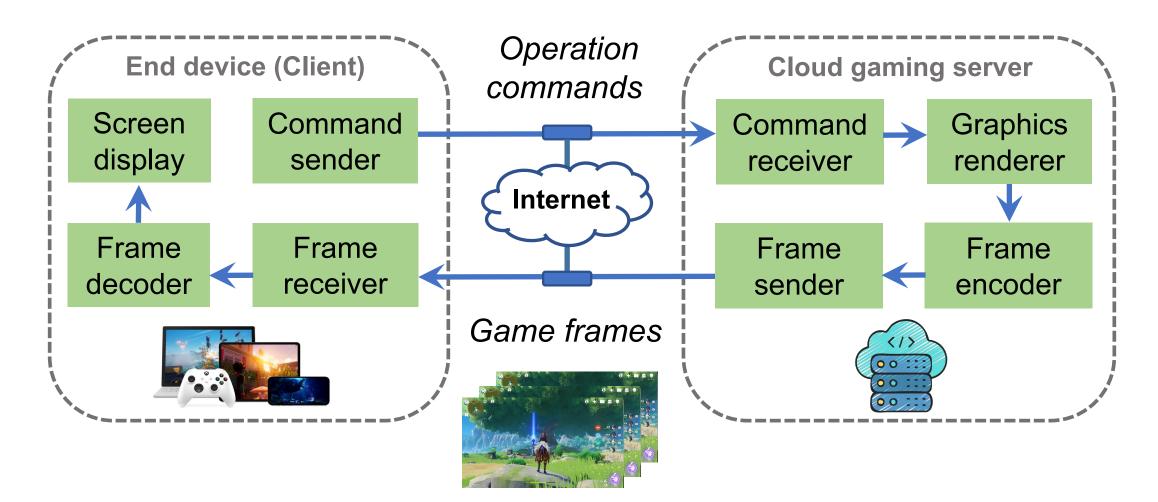
<u>Shibo Wang</u>, Shusen Yang, Xiao Kong, Chenglei Wu, Longwei Jiang, Chenren Xu, Cong Zhao, Xuesong Yang, Jianjun Xiao, Xin Liu, Changxi Zheng, Jing Wang, Honghao Liu

Xi'an Jiaotong University, Tencent Inc., Peking University, Bonree, Tencent America, Columbia University

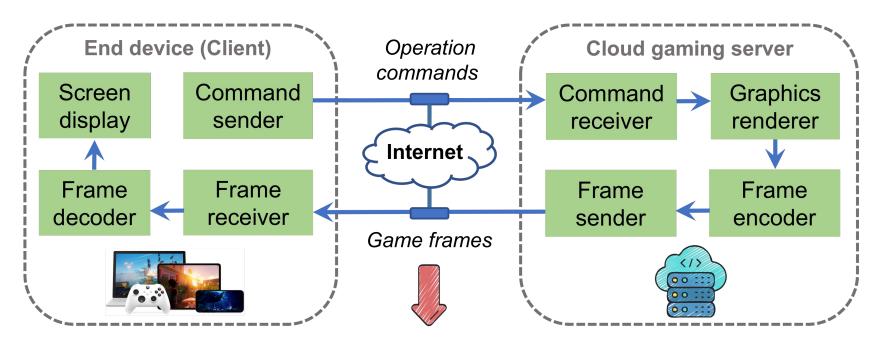
Cloud gaming has already gained world-wide popularity while still under rapid growth



Cloud Gaming System

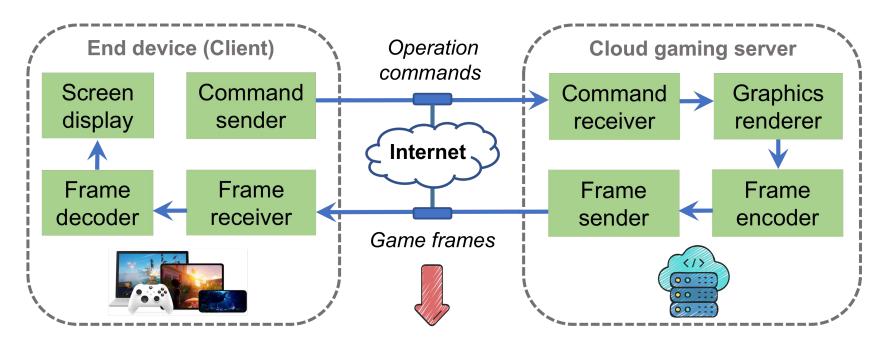


Cloud Gaming System



Consistently demand a low transmission delay for frames

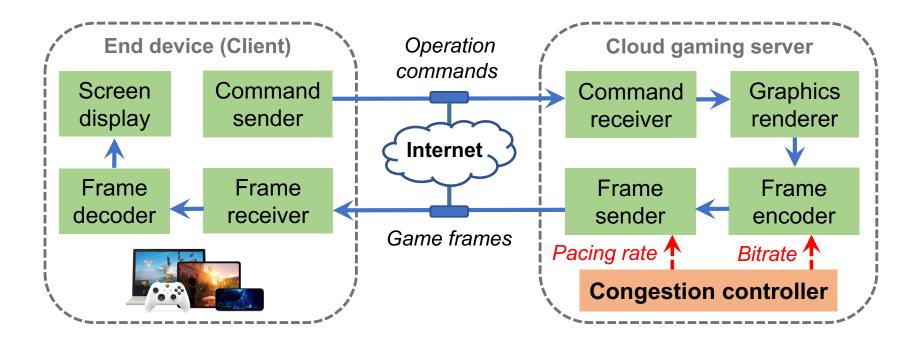
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Consistently demand a low transmission delay for frames

A cloud gaming system needs a carefully designed congestion control (CC) algorithm to manage frame delay.

Congestion Control (CC) for Cloud Gaming

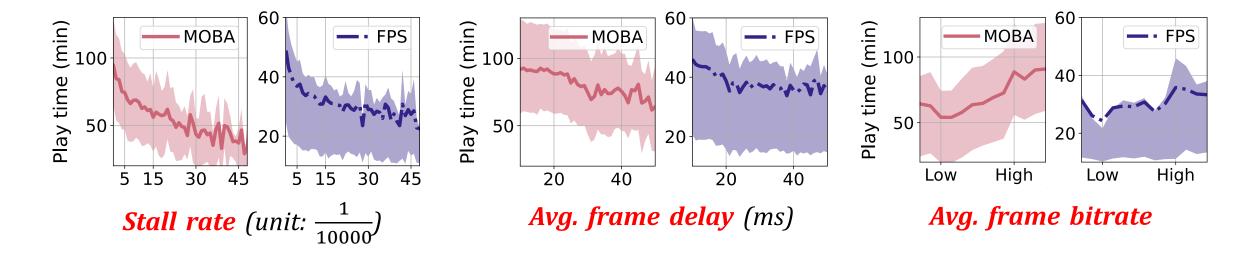


 Cloud gaming CC operates on both application and transport layers, which controls two factors on the fly, namely the *frame bitrate* and *packet sending pace*.

What Cloud Gaming Players Cares?

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- We conducted case studies on two widely-played games¹ to examine the metrics that drive user engagement in cloud gaming:
 - As the stall rate (frame ratio of >100 ms) increases, the play time rapidly decreases.
 - **Frame delay** and **bitrate** also impact the play time, while not as significant as stall rate.



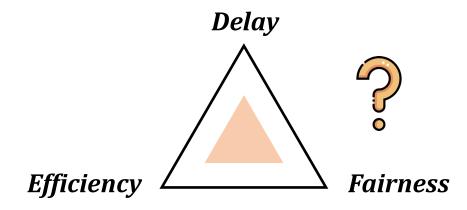
¹A multiplayer online battle arena (MOBA) game and a first-person shooting (FPS) game.

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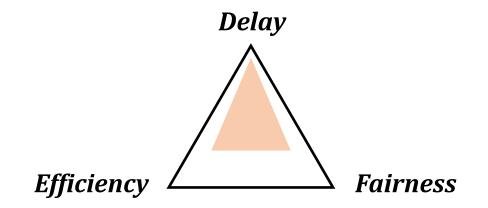
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Frame delay has a higher priority than efficiency or fairness.

• Existing solutions fall short in achieving consistent low frame delay.

Algorithm	Avg. frame delay	95%ile frame delay	Stall rate (>100 ms)
Copa [NSDI'18]	39.8 ms	114 ms	3.2%
Salsify ¹ [NSDI'18]	66.7 <i>ms</i>	186 ms	6.3%
SQP	109.4 ms	287 ms	3.6%
Pudica	23.2 ms	38 ms	0.7%

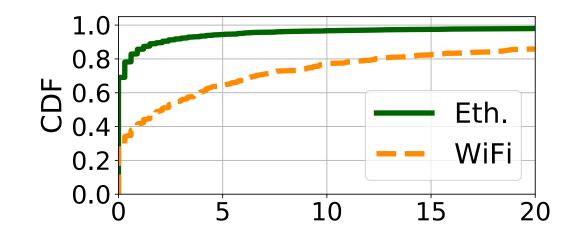
¹Salsify refers to its frame size control part solely.

• Periodical self-induced queue buildups for effective network probing.

Algorithm	Indicator	
Сора	RTTstanding	
Salsify	Packet inter-arrival time	
GCC	Delay gradient	
SQP	Frame transport bandwidth	

When the queue is nearly empty, these indicators fail to provide precise signals.

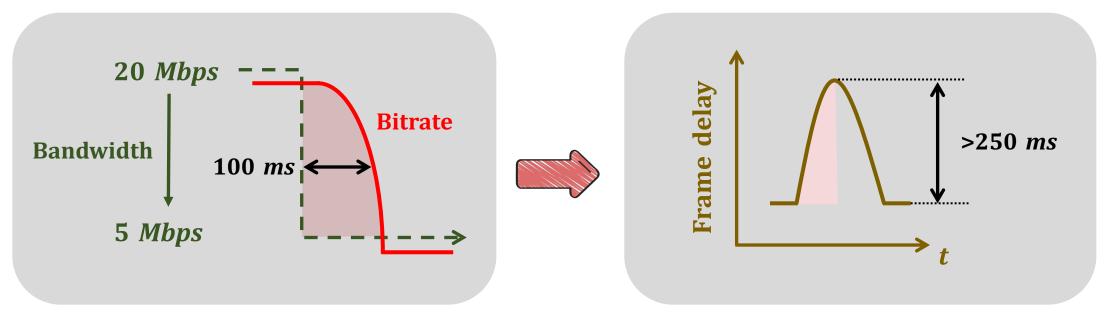
• Slow adaptation to abrupt decreases in available bandwidth.



Num. of bandwidth reductions ($\geq 50\%$) per minute

Internet users frequently encounter significant reductions in available bandwidth.

• Slow adaptation to abrupt decreases in available bandwidth.



A delayed response of only 100 ms

A frame delay spike exceeding 250 ms

Pudica Design: Requirement and Overview

- **Basic requirements for achieving near-zero queuing:**
- Convergence to efficiency (i.e., high link utilization) and fairness without resorting to overshoot-based network probing.
- **Prompt reaction** to abrupt decreases in available bandwidth.

Pudica Design: Requirement and Overview

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- Convergence to efficiency (i.e., high link utilization) and fairness without resorting to overshoot-based network probing.
- **Prompt reaction** to abrupt decreases in available bandwidth.
- Solutions in Pudica:
- Probing the **bandwidth utilization ratio (BUR)** while avoiding frame-level overshooting.
- Bitrate adjustment based on both *smoothed BUR estimations* and *more responsive short-term BUR signals*.

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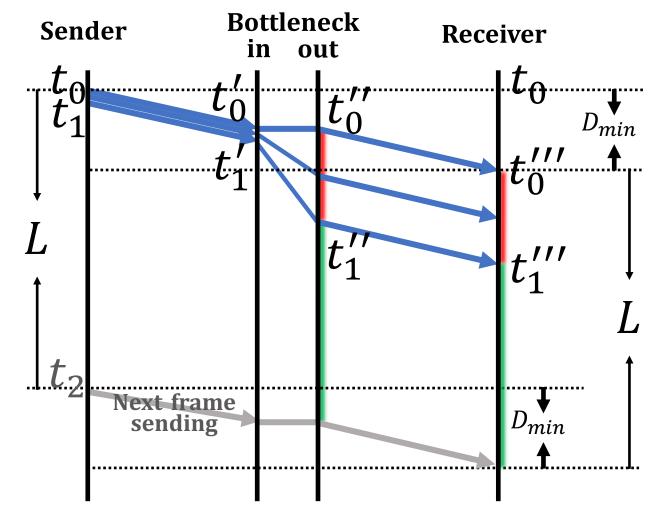
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 - BUR provides an indicator of the precise level of link utilization.
 - BUR has been leveraged by ECN-based CC methods for achieving high link utilization and low bottleneck queuing.

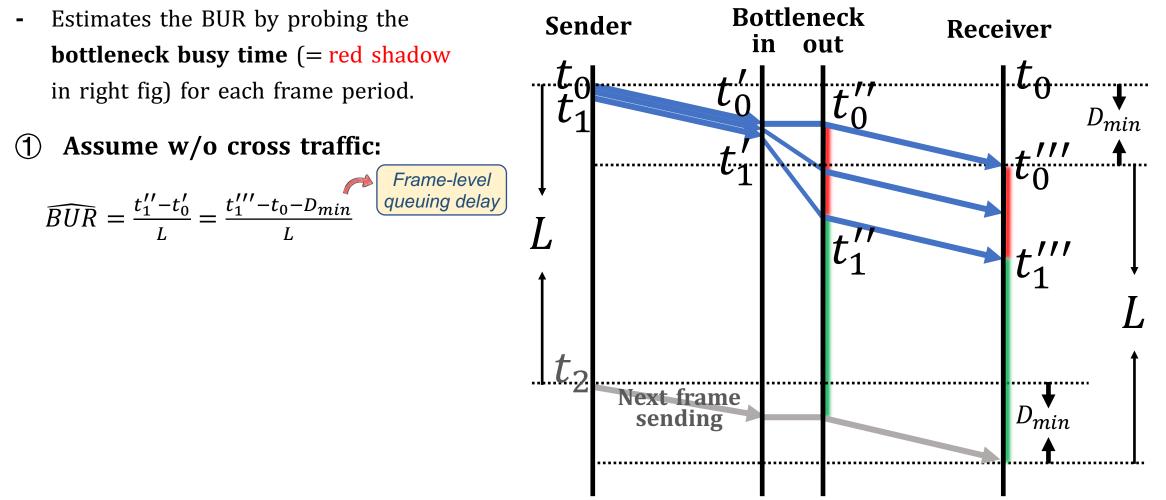
• How to probe?

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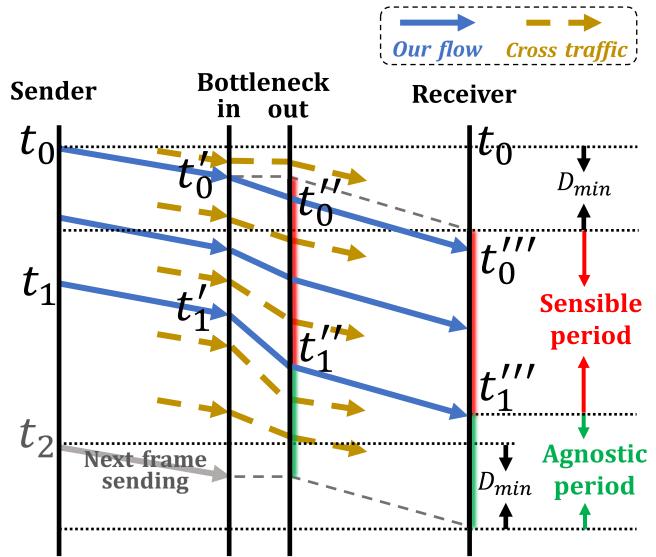
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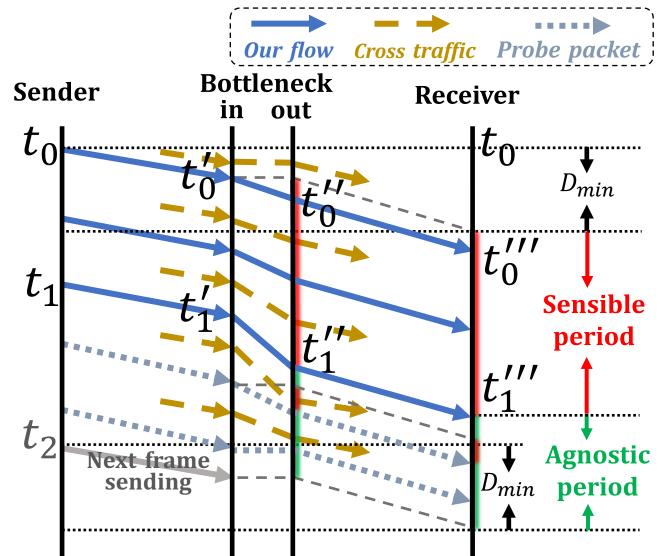
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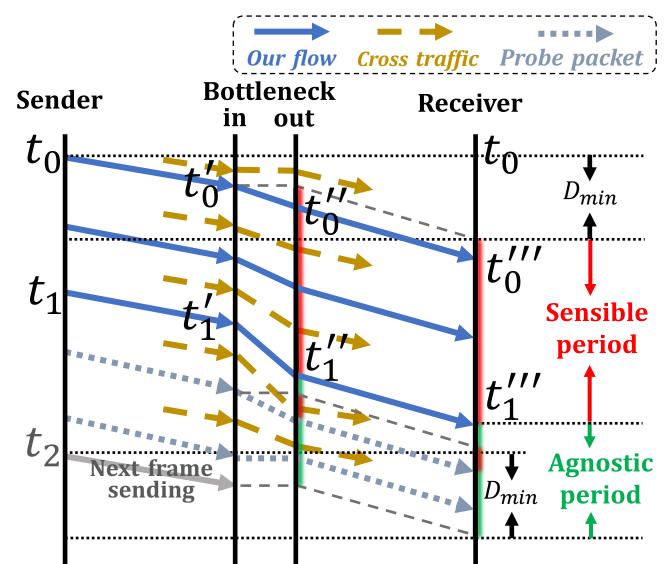
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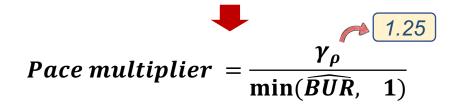
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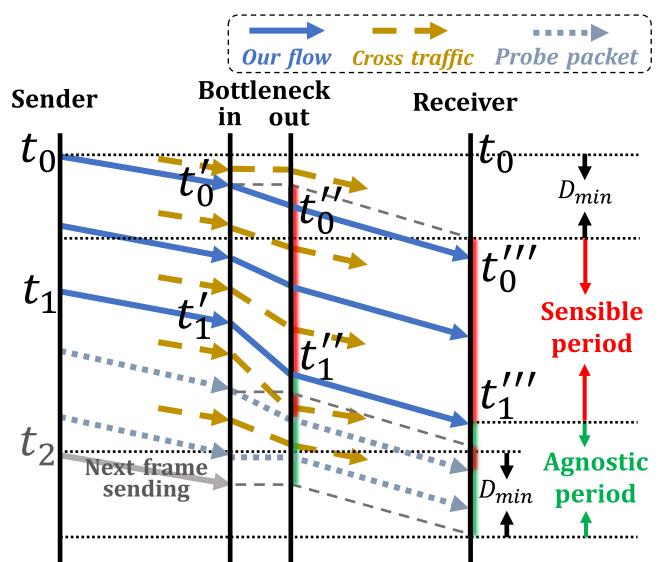


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- When the BURs of **consecutive three frames** exceed one,

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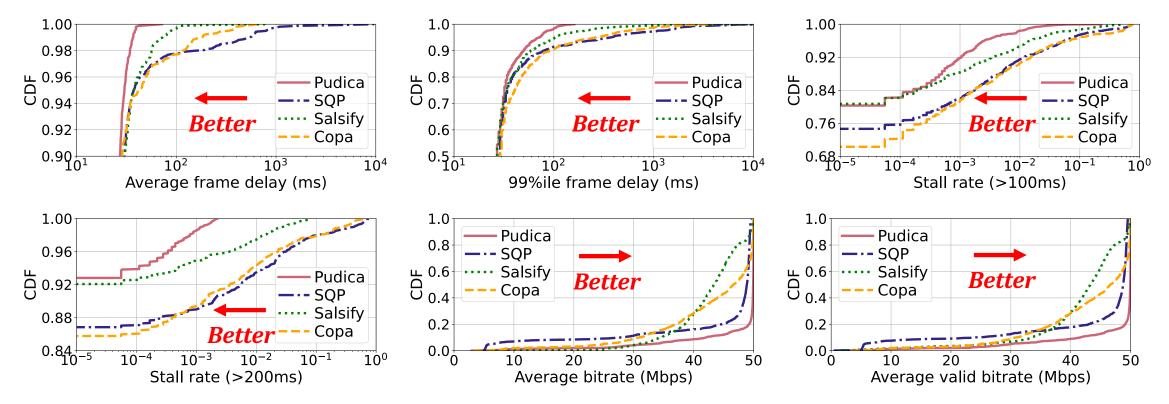
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Evaluation: Methodology

- Deploy four CC methods, including **Pudica, Salsify, Copa, and SQP**, on **Tencent START** cloud gaming platform for large-scale A/B tests.
- The evaluation involved more than **57,000** gaming sessions across **15** cities, **two** network types (**Ethernet and WiFi**), and **three** ISPs over **five** weeks.
- We set the frame rate as **60** and the maximal bitrate as **50 Mbps** for all algorithms.
- **Metrics** (per gaming session):
 - Average, 95%ile, and 99%ile round-trip **frame delay.**
 - Stall rates for frame delays exceeding 100 ms and 200 ms.
 - Average bitrate and valid bitrate (i.e., average bitrate of frames with delay<50 ms).

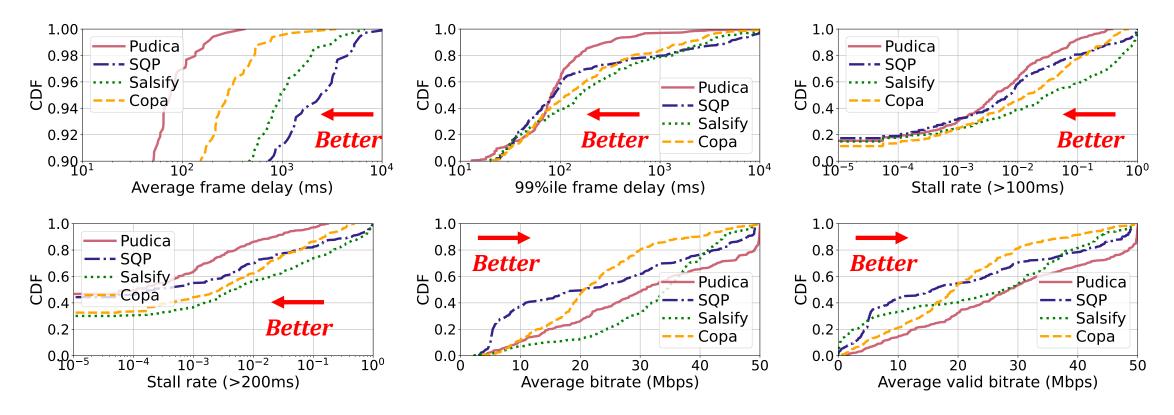
Evaluation: System-Level Performance at Scale

For *Ethernet* networks, Pudica (1) reduces the average and 99%ile frame delay by 1.5× and 3.2×, respectively; (2) reduces the stall rate of 100 *ms* and 200 *ms* by 16.3× and 22.5×, respectively; (3) achieves a slight frame bitrate enhancement.



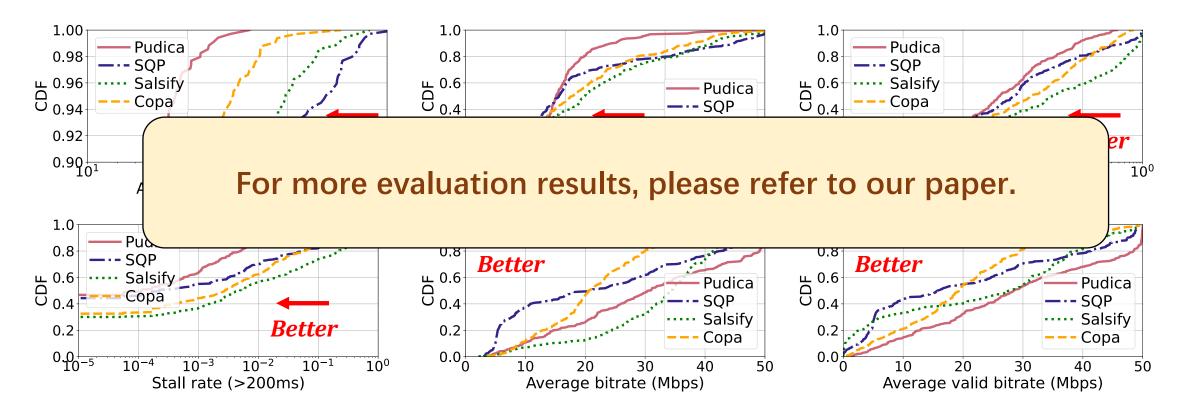
Evaluation: System-Level Performance at Scale

For *WiFi* networks, Pudica (1) reduces the average and 99%ile frame delay by 5.7× and 5.5×, respectively; (2) reduces the stall rate of 100 ms and 200 ms by 5.5× and 12.1×, respectively; (3) achieve the comparable frame bitrate.



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Summary

- We present *Pudica*, an Internet congestion control algorithm designed for cloud gaming:
- **Pudica** proposes a BUR (i.e., *bandwidth utilization ratio*) probing approach and a holistic BUR-based control framework.
- *Pudica* achieves the convergence to efficiency and fairness under the constraint of nearempty bottleneck queues, and the agile adaptation to abrupt bandwidth decreases.
- By conducting large-scale A/B tests on Tencent START cloud gaming services, *Pudica* considerably reduces the frame delay and stall rate while preserving high bitrate.