C2DN: How to Harness Erasure Codes at the Edge for Efficient Content Delivery

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Content Delivery Networks (CDNs)

- Networks of cache clusters close to users
- 72% of Internet traffic





Cache misses are expensive





Unavailabilities at the edge are common

A month-long trace of 2190 clusters

#available servers 9 10 8 8 10 10 10 10 10 10 9 9

> All clusters: unavailability in 45.2% of observations 10-server clusters: unavailability in **30.5%** of observations

Unavailabilities are more common than in datacenters

Reasons: server overload, hardware failure

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Bucket-based routing (coarser load balancing)



Bucket-based routing makes unavailability worse





State-of-the-art solution



Limitations

cannot remove spike

Why?



The problem: write load imbalance

Cache writes come from cache misses

Write load is imbalanced, production: max/min server load = 2.5

- Reduces the effectiveness of replication lacksquare
- SSDs wear out at different rates





State-of-the-art solution

Replication limitations

cannot remove spike waste limited space



Server unavailability mitigation today is costly and ineffective



- removes miss ratio spikes
- reduces bandwidth cost
- near-perfect write load balancing

o spikes h cost load balancing

C2DN Design

Erasure coding to reduce storage overhead Parity rebalance to balance write load

Quick primer on erasure coding: efficient fault tolerance



Use erasure coding in CDN clusters?





Naive use of erasure coding is insufficient

chunks are evicted at different times





Naive use of erasure coding is insufficient

chunks are evicted at different times



- Limited miss ratio reduction
- Cannot eliminate miss ratio spike when unavailability happens
- Runtime overhead (more CPU usage, longer serving latency)

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Goal: rebalance the load

Observation

- flexibility in placing parity
- parity chunks are needed rarely,
 lookup can be slightly more complex

Technique

- data placement: consistent hashing
- parity placement: rebalance write load







Parity assignment problem



Write load of parity Server leftover capacity

Solution: MaxFlow

Balancing write load: similar cache eviction and SSD wear out rates

In the paper

- Several other techniques and optimizations
 - hybrid redundancy
 - sub-chunking
 - transparent coding
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C2DN evaluation

Evaluation setup

- Built C2DN using Apache Trafficserver
- Replayed week-long production traces from Akamai
- Evaluated using three AWS regions



No more miss ratio spike

Reducing normal case miss ratio

Cache size	Miss ratio (bandwidth) reduction
production	21%
2x production	16%
4x production	5%

Erasure coding reduces storage overhead Parity rebalance allows chunks to be evicted at similar time

Near-perfect write load balancing

In the paper

- Time-to-first-byte latency and content download time
 - no noticeable latency change

• CPU and disk usage

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manageable increase

Unavailability and write load imbalance are common in CDN edge clusters

Traditional approach for fault tolerance is not effective in caching

Open sourced at https://github.com

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https://github.com/Thesys-lab/C2DN

Questions?