

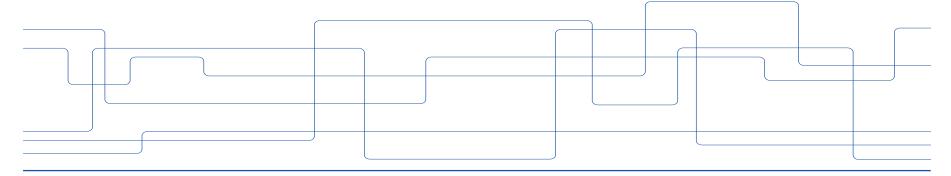
Packet Order Matters!

Improving Application Performance by Deliberately Delaying Packets

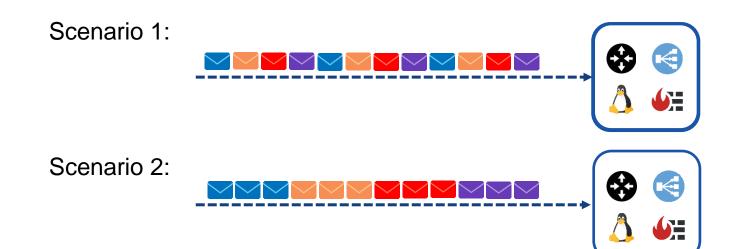
Hamid Ghasemirahni*, Tom Barbette[†], Georgios P. Katsikas*, Alireza Farshin*, Amir Roozbeh*[‡],

Massimo Girondi*, Marco Chiesa*, Gerald Q. Maguire Jr.*, Dejan Kostic*

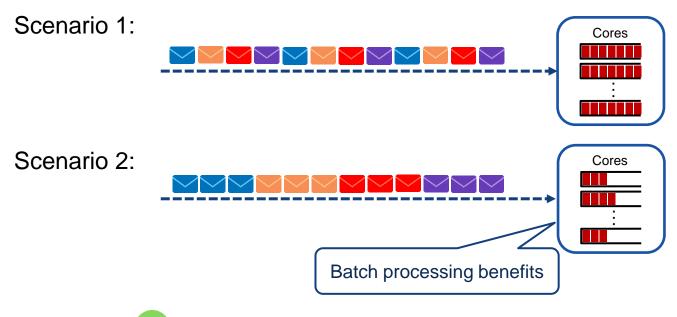
* KTH Royal Institute of Technology, † UCLouvain, ‡ Ericsson Research











Pay the price of re-ordering packets



Main Steps



1- How much does packet order matter? And why?



2- How much is the possibility for reordering packets in a network?



3- Reframer! A software solution to re-order packets!

84% higher throughput with a network function chain

26% improvement in terms of 99.9th percentile latency

How MUCH DOES PACKET ORDER MATTER?



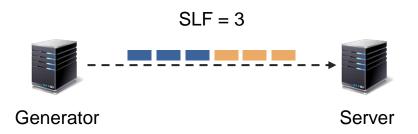


- Spatial Locality Factor (SLF)

Average number of packets, in the **same flow**, that arrive **back-to-back** at a network device

Generate a synthetic train of packets with a given SLF

Measure the performance indicators on the server

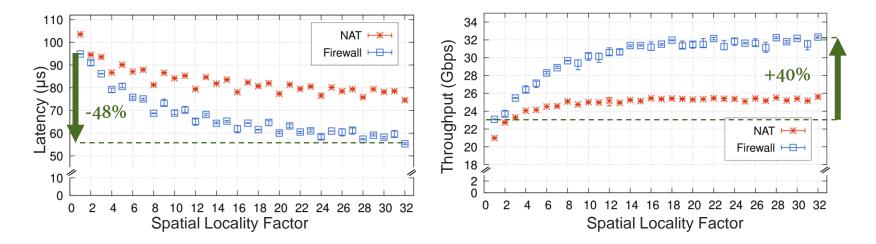




Impact on Network Functions

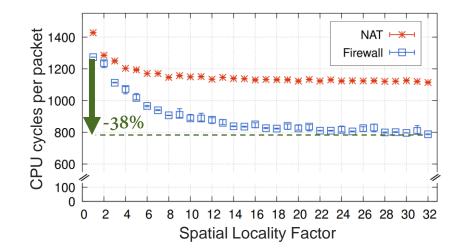
NF use cases: NAT and Firewall

Implementation: FastClick with DPDK



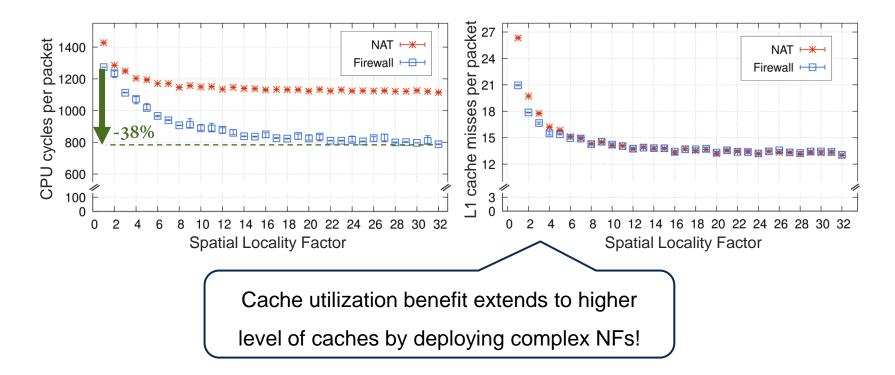


Impact on Network Functions





Impact on Network Functions



HOW MUCH IS THE POSSIBILITY FOR REORDERING PACKETS IN A IN REAL-WORLD NETWORK?



Real-world Trace Analysis

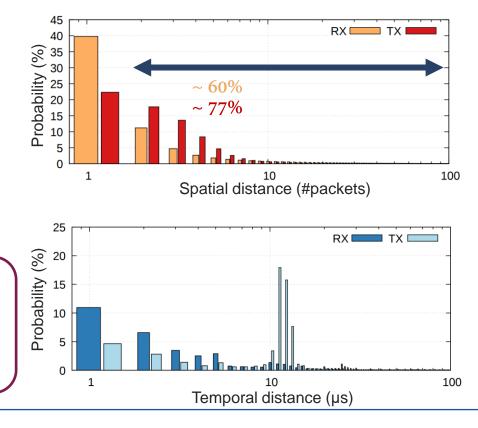
Spatial distance:

Packets gap between two consecutive packets of the same flow

Temporal distance:

Time between two consecutive packets of the same flow

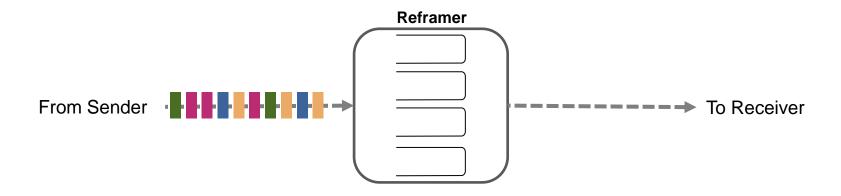
A great opportunity to **reorder** packets by **deliberately delaying** them for a **short period of time**



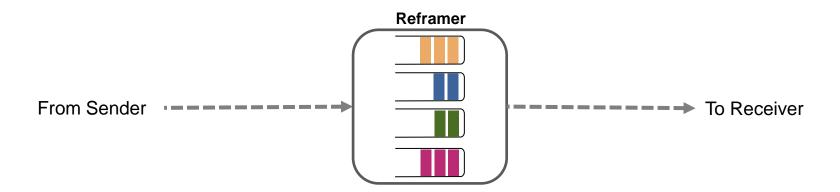
REFRAMER



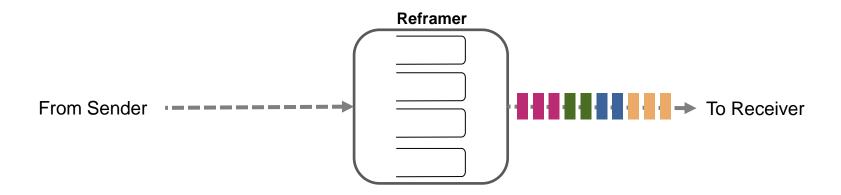




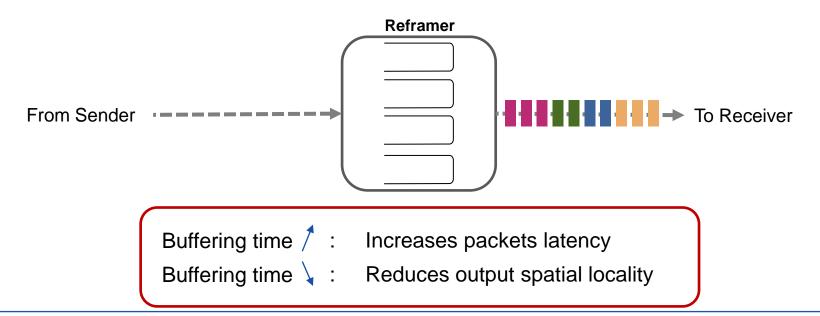
















- Efficiency

Having a small number of operations for buffering and flushing out

Flexibility

Realizing as many scheduling policies as possible

- Scalability

Scaling up linearly by increasing available cores or instances

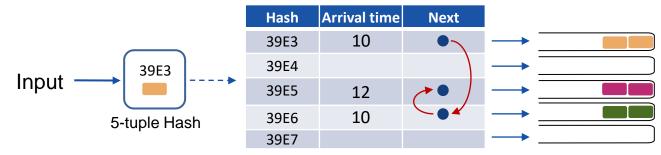




Hash Arrival time Next 39E3 10 ---- 39E4 ---- ---- 39E5 12 ---- 39E6 10 ---- 39E7 ----- -----

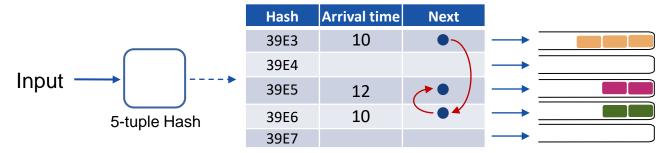
















Hash Arrival time Next 39E3 10 ---- 39E4 ---- ---- 39E5 12 ---- 39E6 10 ---- 39E7 ---- ----





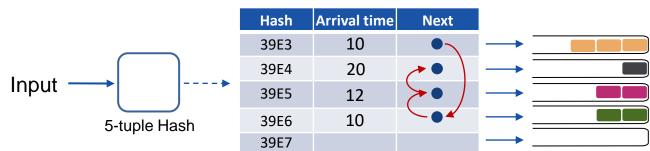




Arrival time Hash Next 10 39E3 39E4 Input 39E5 12 39E6 10 5-tuple Hash 39E7

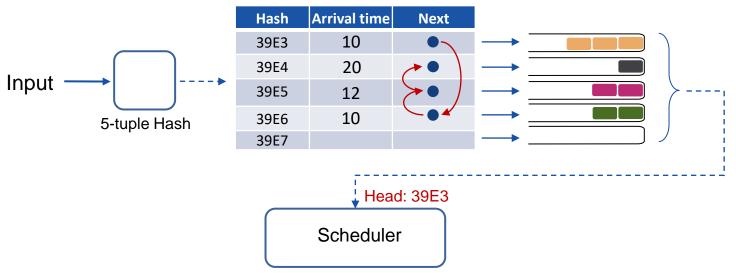






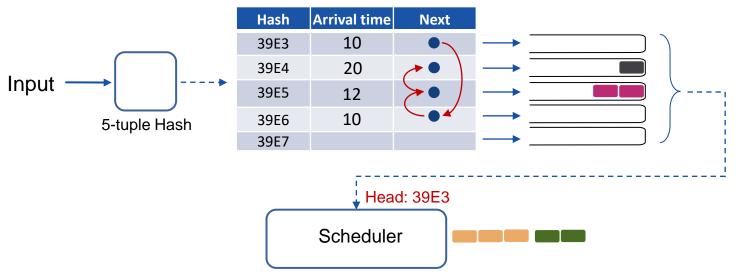






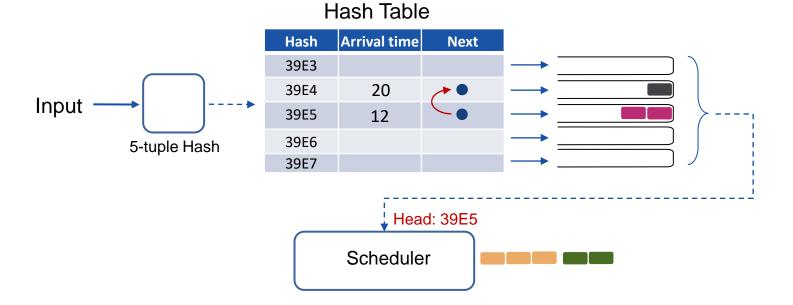
















Hash Table Arrival time Hash Next 39E3 39E4 20 Input 39E5 12 39E6 5-tuple Hash 39E7 L Head: 39E5 Scheduler Protocol Compressor





Hash Table Arrival time Hash Next 39E3 *• 39E4 20 Input 39E5 12 39E6 5-tuple Hash 39E7 L Head: 39E5 Scheduler **Protocol Compressor** Output

EVALUATION



Evaluation

Questions:

Can Reframer increase throughput of an NF chain?

Does the benefit of Reframer persist with the same amount of resources?

How scalable is Reframer?

Testbed:





Evaluation

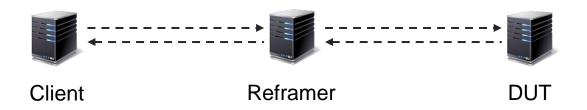
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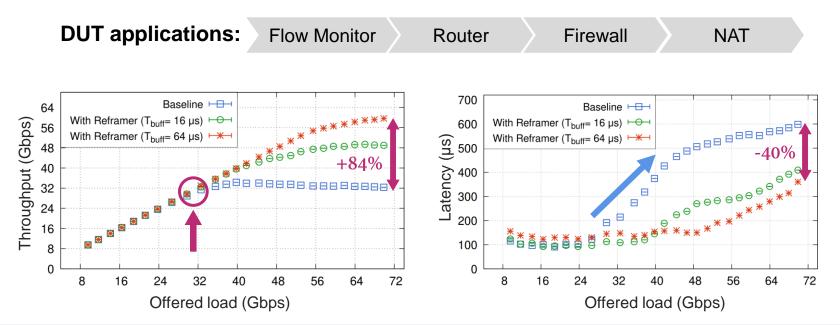
Testbed:







Reframer is deployed in front of a NF chain on a dedicated server

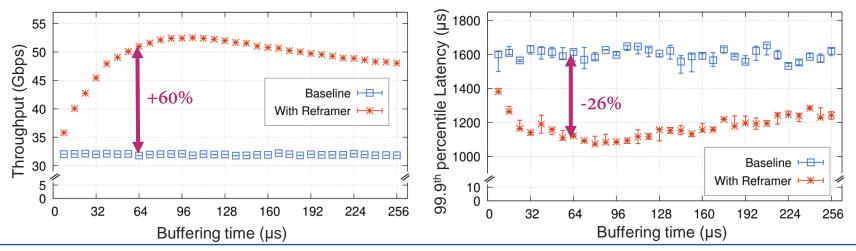






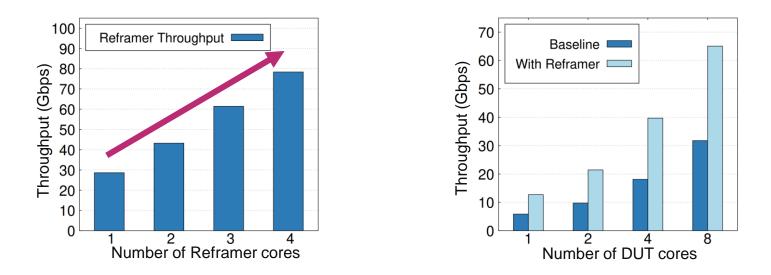
Reframer is deployed on the same server with the NF chain

No additional hardware resource is used to run Reframer!





Reframer scales almost linearly with increasing CPU cores Reframer benefit does not depend on number of DUT cores









Three major steps:

Step1. Packets order impacts network functions and Linux kernel performance

Step2. Our Real-world trace analysis confirms the possibility of ordering packets

Step3. We design, implement, and evaluate Reframer

Evaluation:

Up to 84% benefit in terms of throughput

Up to **26%** improvement in terms of 99.9th percentile latency



Established by the European Commission





STRATEGIC RESEARCH



hamidgh09/Reframer