## TCP-Fuzz: Detecting Memory and Semantic Bugs in TCP Stacks with Fuzzing

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# Background

#### TCP protocol

- Transport-layer network protocol
- Carry over 85% network traffic
- Different implementations
  - Kernel-level TCP stacks
     Linux, FreeBSD, NetBSD
  - User-level TCP stacks
     mTCP, TLDK, F-Stack



## Background

#### • TCP stacks are complex and error-prone

- Rich functionalities: reliable transmission, congestion control
- Complex state model
- Various kinds of possible exceptions
- Common bugs in TCP stacks
  - Memory bugs: null-pointer dereference, use after free
  - Semantic bugs: RFC violation, bad handling of exceptional packet

## Bug study

#### • Study of TCP stack commits

• 87% of bug-fixing commits are related to semantic bugs

Time	FreeBSD		mTCP		TLDK	
	Memory	Semantic	Memory	Semantic	Memory	Semantic
2017	2	26	2	6	1	11
2018	9	51	0	4	0	4
2019	9	65	1	3	2	5
Total	20	142	3	13	3	20

### Example

#### • Semantic bug fix in FreeBSD TCP stack

#### RFC 7323 violation

```
FILE: FreeBSD/sys/netinet/tcp_syncache.c
int syncache expand(...) {
    /* RFC 7323 PAWS: If we have a timestamp on this segment and
      * it is less than ts_recent, drop it.
     if (sc->sc_flags & SCF_TIMESTAMP && to->to_flags & TOF_TS &&
       TSTMP LT(to->to tsval, sc->sc tsreflect)) {
           SCH UNLOCK(sch);
           if ((s = tcp_log_addrs(inc, th, NULL, NULL))) {
                log(LOG_DEBUG, ...);
                free(s, M TCPLOG);
           return (-1); /* Do not send RST */
      .....
```

## **TCP** stack features

#### • F1: two-dimensional inputs with dependencies

- Inputs: syscalls, packets
- Syscalls and packets have dependencies with each other

#### • F2: state model

- Basic model in RFC 793: 11 states and 20 state transitions
- Real-world TCP stacks have many states and state transitions

#### • F3: semantic rules

- Stipulate how syscalls and packets should be handled
- Explicit rules: defined in RFC documents
- Implicit rules: no explicit description in RFC documents

# Key Techniques

#### • F1: two-dimensional inputs with dependencies

- Dependency-based strategy to generate effective inputs
- F2: state model
  - Transition-guided fuzzing approach to improve the coverage of states and state transitions

#### • F3: semantic rules

• Differential checker to detect semantic bugs

## Dependency-based strategy

#### Input sequence mutation

- Item type: packet and syscall
- Select mutation type: *deletion*, *addition*, *replacement* and *change*
- Generate new items: consider dependencies with previous items



## Dependency-based strategy

#### • Dependency

- Refer to RFC documents and syscall-usage conventions
- Type: syscall-syscall, packet-packet, syscall-packet
- Examples
  - Syscall-syscall: socket, bind, listen and accept are called in order when a connection is passive open
  - Packet-packet: the order and control flags of three-way handshake and four-way handshake packets are never changed
  - Syscall-packet: accept can be called only after the three-way handshake when a connection is passive open

# Transition-guided fuzzing approach

- Old coverage metric: code coverage
  - Describe different states
  - Neglect different state transitions



## Transition-guided fuzzing approach

• New coverage metric: **Branch transition** 

 $BrTran_n = \langle BrCov_n, BrCov_n - BrCov_{n-1} \rangle$ 

 $BrTran_n$  is the branch transition of input\_item<sub>n</sub>  $BrCov_n$  is the branch coverage of input\_item<sub>n</sub>



## **Differential checker**

o Basic idea

 Different TCP stacks obey identical semantic rules and produce identical or similar outputs for the same inputs

o Design

 Provide the same input sequences to different TCP stacks and compare their outputs



### Framework

#### • TCP-Fuzz

- Novel fuzzing framework for testing TCP stacks
- Integrate the three key techniques
- Detect memory and semantic bugs



## Deployment

#### Server-client mode

- o Server: generate test cases and validate data
- o Client: run the tested TCP stack with third-party sanitizers



## Evaluation

#### • Experimental setup

- Three user-level and two kernel-level TCP stacks
- Each TCP stack is tested for 48 hours

Туре	TCP stack	Version	LOC
	TLDK	v2.0	15K
User-level	F-Stack	Commit 8d21adc	25K
	mTCP	Commit 0463aad	18K
Korpol lovol	FreeBSD	v12.1	171K
Kernei-levei	Linux	v5.6	169K

# Testing coverage and found bugs

- Branch transitions > branches
- 8 memory bugs and 48 semantic bugs
- 40 bugs have been confirmed

Stock	Testing coverage		Found bugs		
Stack	Branch	Transition	Memory/Semantic	Confirmed/Fixed	
TLDK 1.3K		329.4K	2/26	28/19	
F-Stack	7.5K	46.8K	1/6	6/1	
mTCP	1.2K 47.9K		5/9	0/0	
FreeBSD	-	-	0/6	5/2	
Linux			0/1	1/1	
Total	10.0K	424.1K	8/48	40/23	

## Compared to state-of-the-art fuzzers

• Two classical and widely-used fuzzing approaches

- AFL-like: only generates packet sequences
- Syzkaller-like: only generates syscall sequences
- Three open sourced protocol fuzzing approaches
  - Boofuzz: <u>github.com/jtpereyda/boofuzz</u>
  - Fuzzotron: <u>github.com/denandz/fuzzotron</u>
  - AFLNet: <u>github.com/aflnet/aflnet</u> [ICST'20]



## Compared to state-of-the-art fuzzers

#### • Testing coverage

 TCP-Fuzz covers more branch transitions than AFL-like and Syzkaller-like fuzzers



## Compared to state-of-the-art fuzzers

#### • Bug detection

• TCP-Fuzz finds more bugs than other fuzzers



## Conclusion

- TCP stacks are complex and error-prone
- TCP-Fuzz
  - **Dependency-based strategy** to generate effective inputs
  - Transition-guided fuzzing approach to improve the coverage of states and state transitions
  - Differential checker to detect semantic bugs
- Find 56 real bugs in 5 widely-used TCP stacks
- Outperform state-of-the-art fuzzers



# Thanks

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