SPINFER: Inferring Semantic Patches for the Linux Kernel

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Maintenance of the Linux kernel

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But maintaining the Linux kernel is particularly hard:

- 18M lines of C code
- 13M lines of driver code
- The same kernel API can be used by thousands of files

Even simple API migrations can be difficult to do

Motivating Example

Example of low-resolution timer structure initialization:

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- Originally with the init_timer function
- Since 2006 with setup_timer

Old function was not removed, the migration was not mandatory.

init_timer migration

```
drivers/atm/nicstar.c
@@ -284,10 +284,8 @@ static int __init nicstar_init(void)
- init_timer(&ns_timer);
+ setup_timer(&ns_timer, ns_poll, OUL);
    ns_timer.expires = jiffies + NS_POLL_PERIOD;
- ns_timer.data = OUL;
- ns_timer.function = ns_poll;
```

drivers/gpu/drm/omapdrm/dss/dsi.c

00 -5449,9 +5449,7 00 static int dsi_bind(struct device *dev,

- init_timer(&dsi->te_timer);
- dsi->te_timer.function = dsi_te_timeout;
- dsi->te_timer.data = 0;
- + setup_timer(&dsi->te_timer, dsi_te_timeout, 0);



In 2018 these interfaces were considered insecure and were both replaced.

But at this time API usage was in inconsistent state:

- 60% using the new setup_timer
- 40% using the old init_timer

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Could the transformation have been done automatically?

First contribution: Taxonomy of transformation challenges There are a lot of tools to perform API migration by learning from examples: REFAZER, LASE, AppEvolve, Meditor, ...

But it was hard to know what kind of transformation they could handle.

Our first contribution is to classify transformation challenges.

1. Control-flow dependencies

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- 2. Data-flow dependencies

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- 2. Data-flow dependencies
- 3. Number of variants

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- 2. Data-flow dependencies
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- 5. Presence of unrelated changes

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Both of these constraints are common in Linux kernel transformations.

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Both of these constraints are common in Linux kernel transformations.

And they were necessary for our timer example.

Moreover transformation rules used by these tools are not exposed Meaning that developers cannot check if the transformation will be correct.

Second contribution: Spinfer

To perform API migration in the Linux kernel we want a tool that:

- Learns transformation from examples
- Handles both control-flow dependencies and transformation variants
- Exposes transformation rules to developers

Fortunately, a transformation rules language is already used in the Linux kernel.

Since 2008 Coccinelle rules are used to perform some transformations.

Even used in our motivating example.

Coccinelle



Semantic patch

00

```
expression EO, E1, E2;
@@
```

- init_timer(E0);
- + setup_timer(E0, E1, E2);

• • •

- E0.data = E2;
- EO.function = E1;

00

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Generates diffs like this:

- init_timer(&ns_timer);
- + setup_timer(&ns_timer, ns_poll, OUL);
 ns_timer.expires = jiffies + NS_P_P;
- ns_timer.data = OUL;
- ns_timer.function = ns_poll;

Our approach: Spinfer



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Infering semantic patches

How to convert transformation instances...

- init_timer(&ns_timer);

- + setup_timer(&ns_timer, ns_poll, OUL);
 ns_timer.expires = jiffies + NS_P_P;
- ns_timer.data = OUL;
- ns_timer.function = ns_poll;

... to a semantic patch.

@@
expression E0, E1, E2;
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- init_timer(E0);
- + setup_timer(E0, E1, E2);

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- E0.data = E2;
- E0.function = E1;

1: Extracting modified statements

- init_timer(&ns_timer);
- + setup_timer(&ns_timer, ns_poll, OUL);
 ns timer.expires = jiffies + NS POLL PERIOD;
- ns_timer.data = OUL;
- ns_timer.function = ns_poll;
- init_timer(&dsi->te_timer);
- dsi->te_timer.function = dsi_te_timeout;
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- + setup_timer(&dsi->te_timer, dsi_te_timeout, 0);

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- dsi->te_timer.function = dsi_te_timeout;
- dsi->te_timer.data = 0;
- + setup_timer(&dsi->te_timer, dsi_te_timeout, 0);

2: Clustering similar statements

- init_timer(&ns_timer);
- init_timer(&dsi->te_timer);
- + setup_timer(&ns_timer, ns_poll, OUL);
- + setup_timer(&dsi->te_timer, dsi_te_timeout, 0);
- ns_timer.data = OUL;
- dsi->te_timer.data = 0;
- ns_timer.function = ns_poll;
- dsi->te_timer.function = dsi_te_timeout;

3: Abstracting clusters

- init_timer(*Expr*); init_timer(*Masi-*);

+ setup_timer(*Expr, Expr, Expr*); + setup_timer(*a*dsi->te_timer, dsi_te_timeout, 0);

- *Expr*.data = *Expr*;
- *Expr*.function = *Expr*;^{pol1};
- dsi->te_timer.function = dsi_te_timeout;

4: Assembling abstractions

- init_timer(Expr);
- *Expr*.function = *Expr*;

- Expr.data = Expr;
- + setup_timer(Expr, Expr, Expr);

- init_timer(Expr);
- *Expr*.function = *Expr*;

Spinfer takes a first abstraction

- init_timer(Expr);

- Expr.data = Expr;
- + setup_timer(Expr, Expr, Expr);

- init_timer(Expr);
- Expr.function = Expr;

- Expr.data = Expr;
- + setup_timer(Expr, Expr, Expr);

It extends rules using control-flow dependencies

- init_timer(Expr);
- . . .
- *Expr*.function = *Expr*;

When there are inconsistencies in control-flow, rules are split:

- init_timer(Expr);
- • •
- Expr.data = Expr;
- *Expr*.function = *Expr*;

- init_timer(Expr);
- Expr.function = Expr;
- Expr.data = Expr;

. . .

This allows Spinfer to discover transformation variants.

This process goes on until all abstractions are exhausted.

- init_timer(Expr);
- + setup_timer(*Expr*, *Expr*, *Expr*);
- • •
- Expr.data = Expr;
- *Expr*.function = *Expr*;

- init_timer(Expr);
- + setup_timer(Expr, Expr, Expr);
- • •
- *Expr*.function = *Expr*;
- *Expr*.data = *Expr*;

To obtain a valid rule Spinfer transforms abstractions into metavariables:

A unique name is chosen for each set of terms found in the examples.

```
@@
expression E0, E1, E2;
@@
```

- init_timer(Expr); init_timer(E0);
- + setup_timer(Expr, Expr, Expr); + setup_timer(E0, E1, E2);
- • •
- **Expr**.data = **Expr**;
- **Expr**.function = **Expr**;

- • •
- E0.data = E2;
- EO.function = E1;

Obtained semantic patch

Spinfer obtained these two rules:

00

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```

- init_timer(E0);
- + setup_timer(E0, E1, E2);

. . .

- **E0**.data = **E2**;
- EO.function = E1;

@@
expression E0, E1, E2;
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- init_timer(E0);
- + setup_timer(E0, E1, E2);

. . .

- E0.function = E1;
- E0.data = E2;

Evaluation

We evaluated Spinfer by learning real Linux kernel transformations.

We extracted two datasets of 40 groups of transformation each:

- One selected to be challenging
- Another randomly sampled from changes in 2018

We compared the results produced by Spinfer generated semantic patches to the results produced by a human written semantic patch.

Results on the randomly sampled dataset

Spinfer was learning on one part of the changes and evaluated on the other part.

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Two metrics:

- Precision: fraction of changes produced that were correct
- Recall: fraction of needed changes that were produced

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- Precision: fraction of changes produced that were correct
- Recall: fraction of needed changes that were produced

Spinfer obtained 87% precision and 62% recall in average.

In 8 cases Spinfer obtained a perfect semantic patch.

More experiments on the paper

Conclusion

Spinfer learns semantic patches from examples.

It can learn transformations variants with many constraints such as:

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It uses code clustering to find similar pieces of code and abstract them.

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Spinfer learns semantic patches from examples.

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It uses code clustering to find similar pieces of code and abstract them.

Abstractions are assembled using control-flow information.

Produced semantic patches can be checked and fixed by developers.

Thank you

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