

vCRIB: Virtual Cloud Rule Information Base

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Introduction

Datacenters use rules to implement **management policies**

- Access control
- Rate limiting
- Traffic measurement
- Traffic engineering

Introduction

Datacenters use **rules** to implement management policies

An **action** on a hypercube of **flow space**

Examples:

- Deny
- Normal
- Enqueue

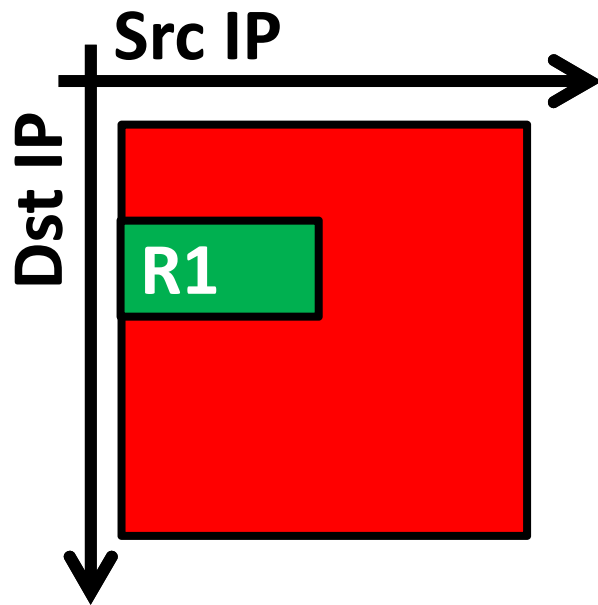
Flow fields examples:

- Src IP / Dst IP
- Protocol
- Src Port / Dst Port

Introduction

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An **action** on a hypercube of **flow space**

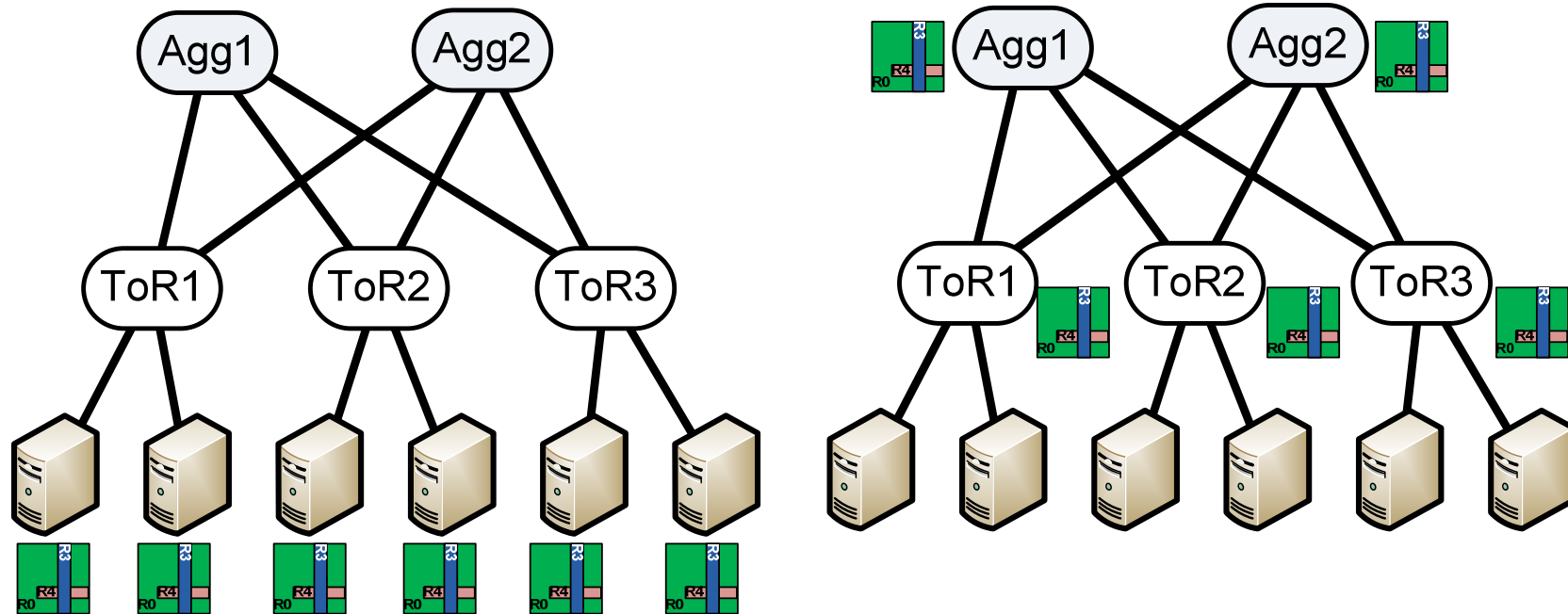


R1=Accept

- SrcIP: 12.0.0.0/8
- DstIP: 10.0.0.0/16

Current Practice

Rules are saved on **predefined fixed** machines



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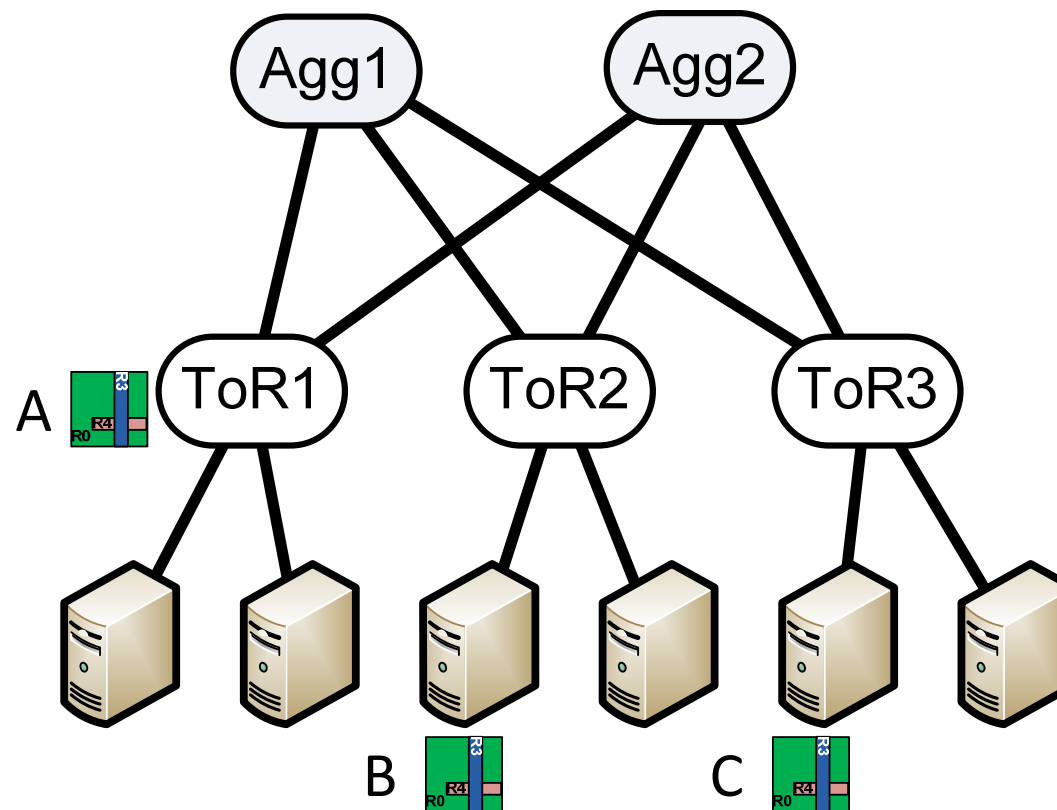
1 Machines have **limited resources**

2 Datacenters have **different resource constraints**

3 Multiple policies may **compete for resources**

vCRIB Goal: Flexible Rule Placement

Find the best feasible rule placement based on resource constraints



Future Datacenters will have many fine-grained rules

Regulating VM pair communication

- Access control (CloudPolice)
- Bandwidth allocation (Seawall)

100K – 1M

Per flow decision

- Flow measurement for traffic engineering (MicroTE, Hedera)

10 – 100M

VLAN per server

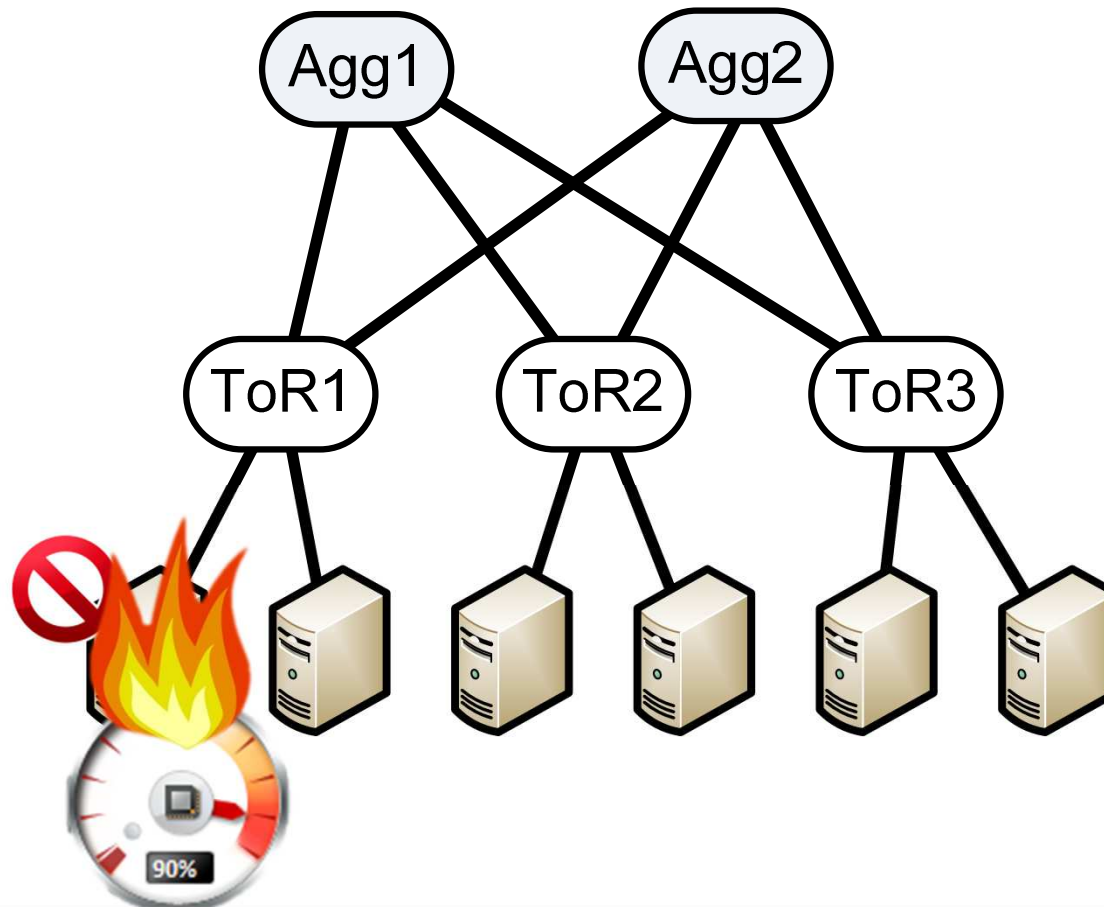
- Traffic management (NetLord, Spain)

1M

Where to place rules? Hypervisors vs. Switches

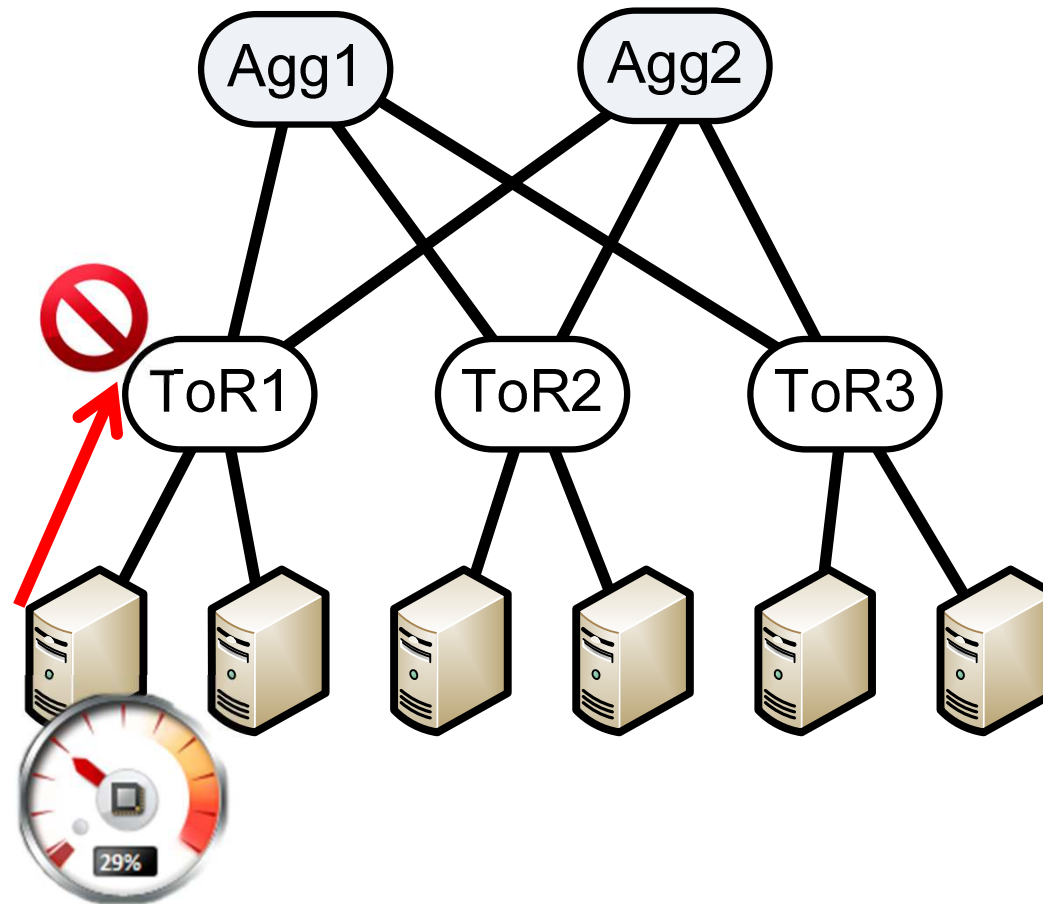
	Hypervisor	Switch
Performance	Software, Slow	Hardware, Fast
Flexibility	Complex rules	OpenFlow rules
Entry point	Close to VMs	External traffic, Aggregate traffic
Resources	Limited CPU budget	# TCAM entries

Rule Location Trade-off (Resource vs. Bandwidth Usage)



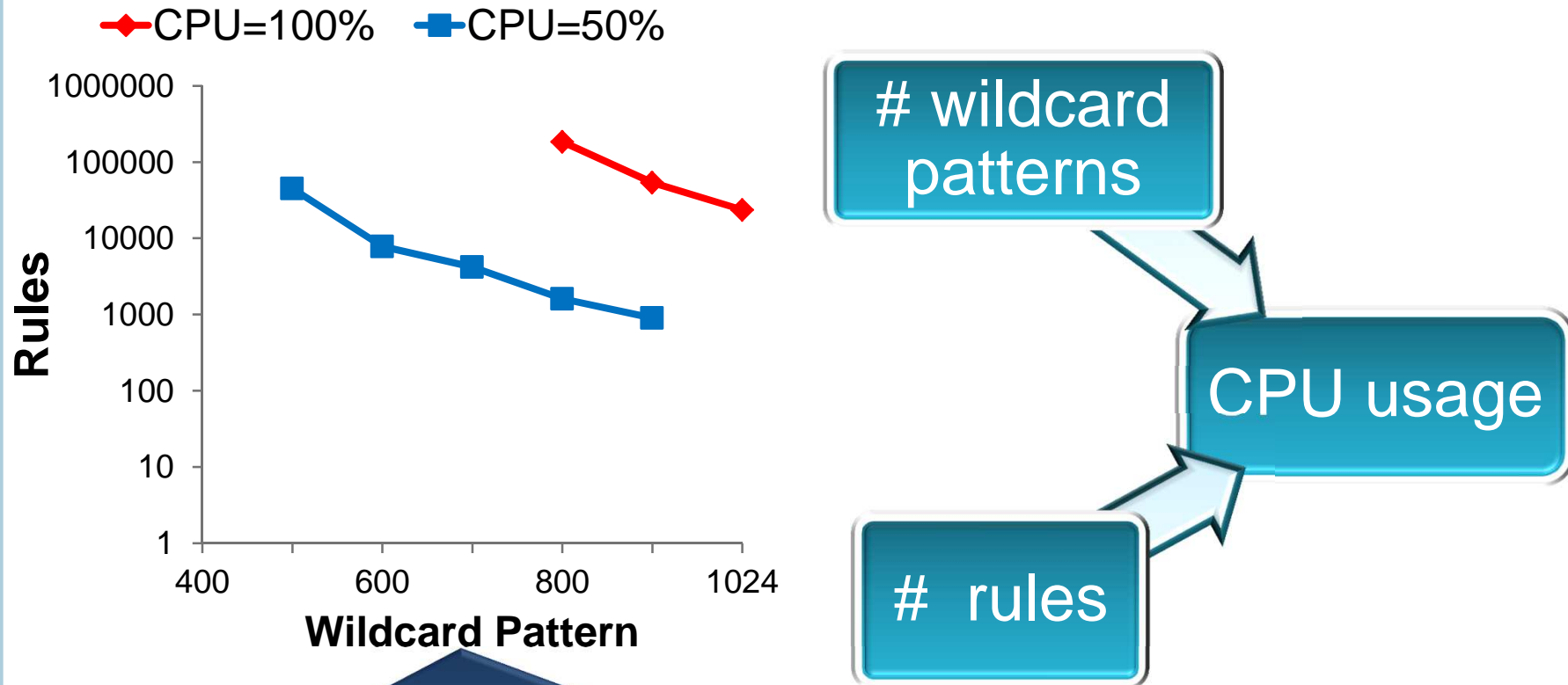
Storing rules at hypervisor incurs CPU processing overhead

Rule Location Trade-off (Resource vs. Bandwidth Usage)



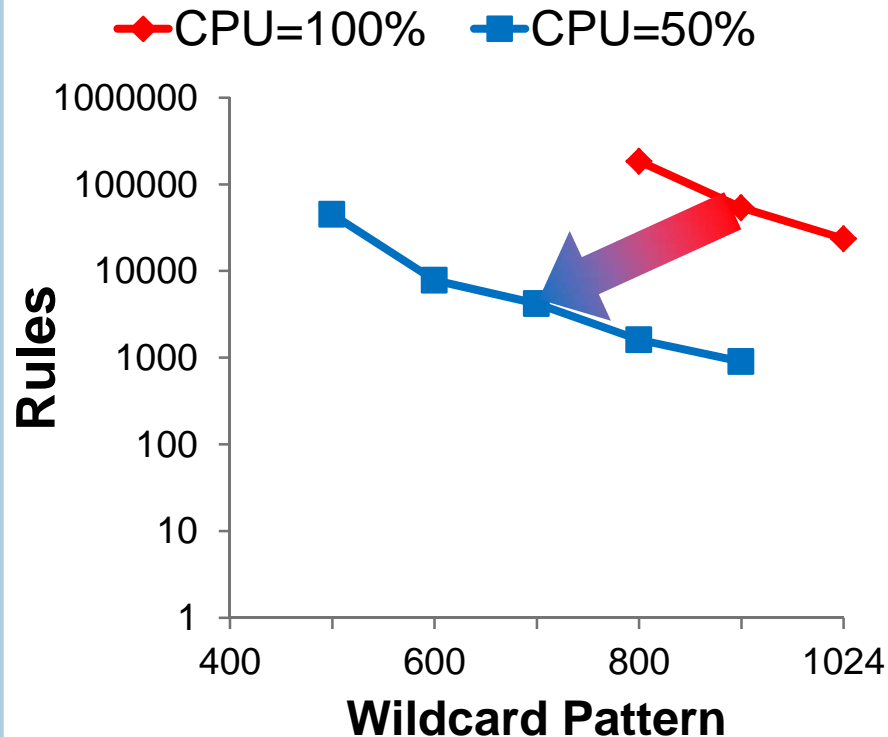
Move the rule to ToR switch and forward traffic

Can we reduce Open vSwitch CPU usage?



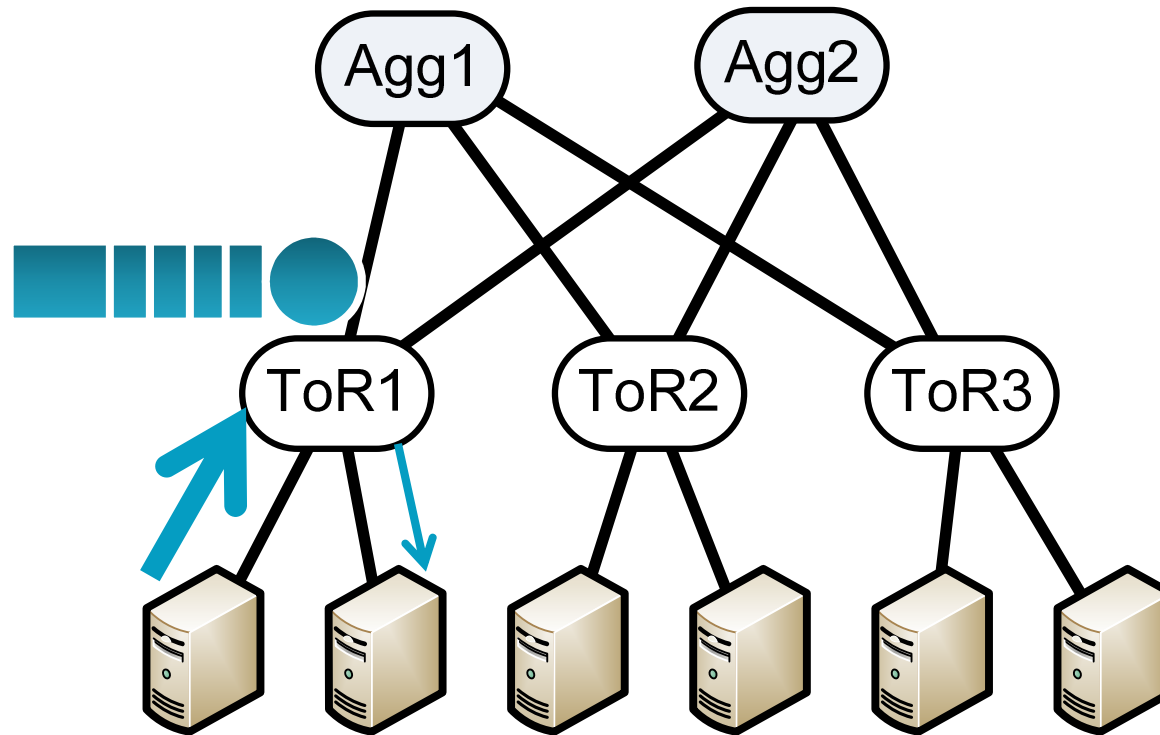
The set of ignore bits in the mask
R1=Accept, DstIP: 10.0.0.0/16, SrcIP: 12.0.0.0/8
11111111111111111111*****
11111111*****

Can we reduce Open vSwitch CPU usage?



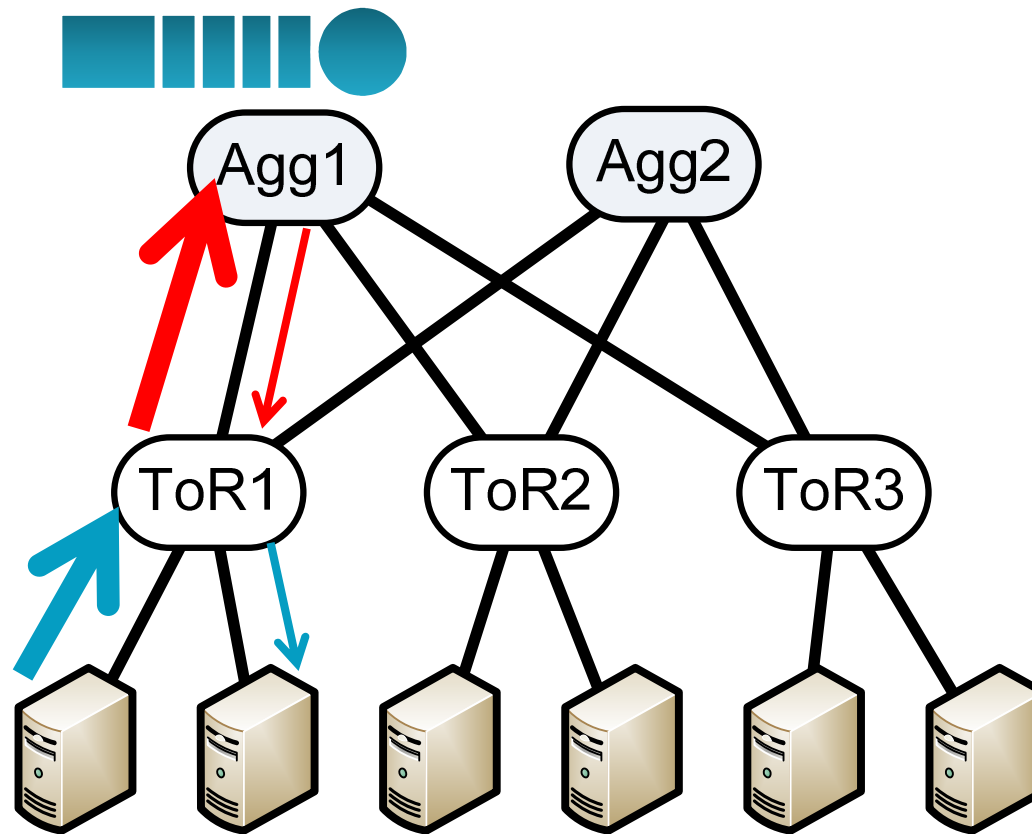
Handle same number of new flows with lower CPU budget

Rule Location Trade-off (Resource vs. Bandwidth Usage)



If rule memory is limited in one switch

Rule Location Trade-off (Resource vs. Bandwidth Usage)

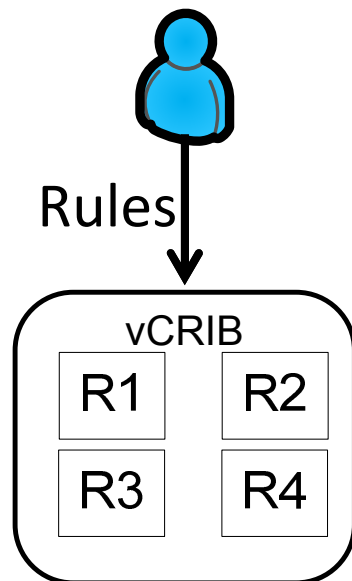


Can tradeoff bandwidth within the switch fabric,
in addition to trading-off bandwidth between
hypervisors and switches

Our Approach: vCRIB, a Virtual Cloud Rule Information Base

Proactive rule placement abstraction layer

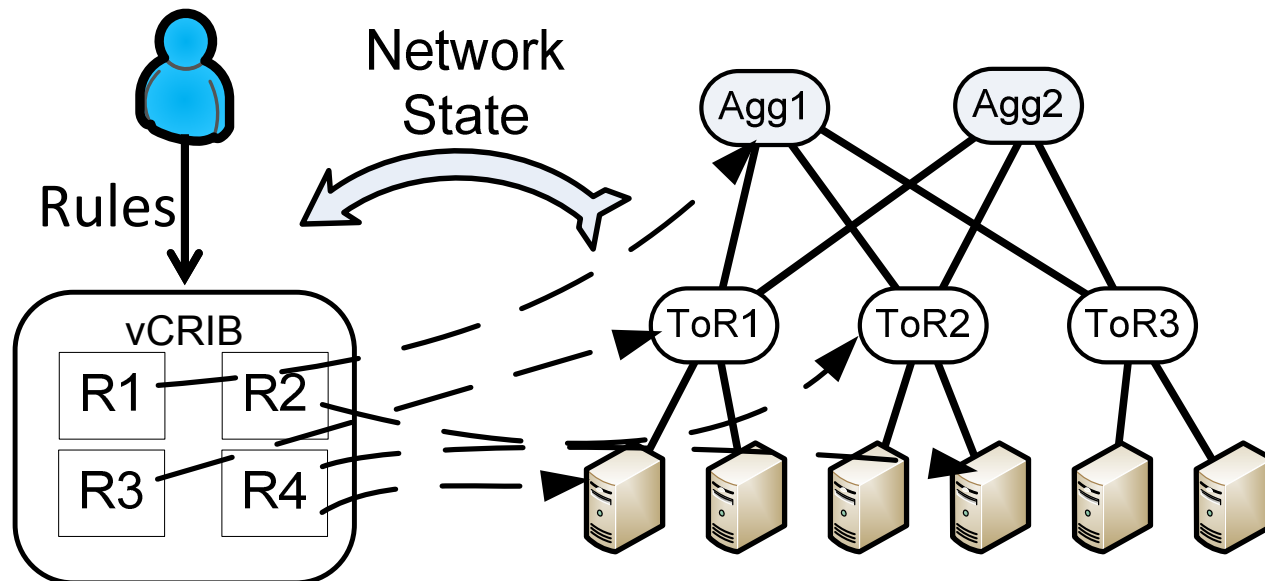
Allow operators to define fine-grained rules without worrying about placement



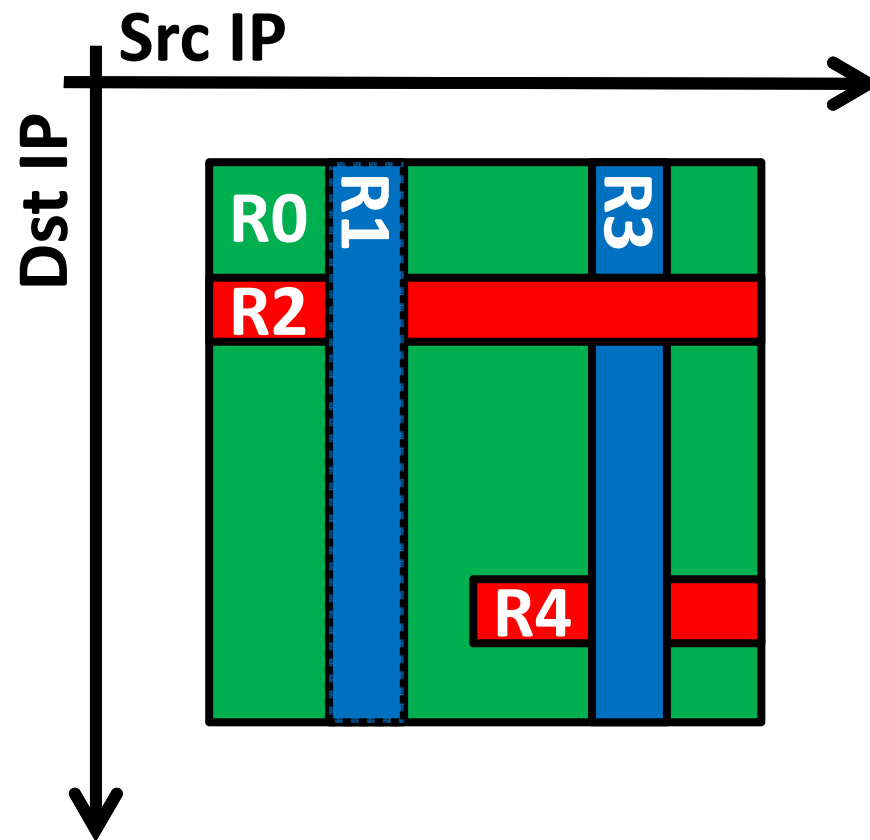
Our Approach: vCRIB, a Virtual Cloud Rule Information Base

Flexible rule placement at hypervisors and switches

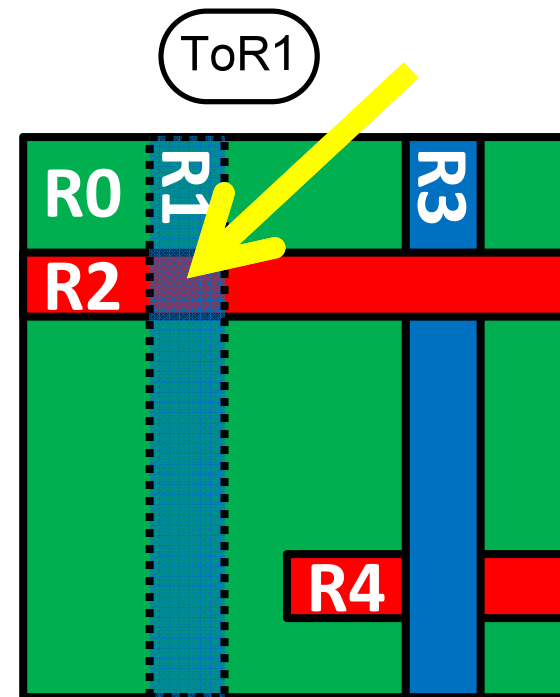
Optimize performance given resource constraints



Challenges: Overlapping Rules

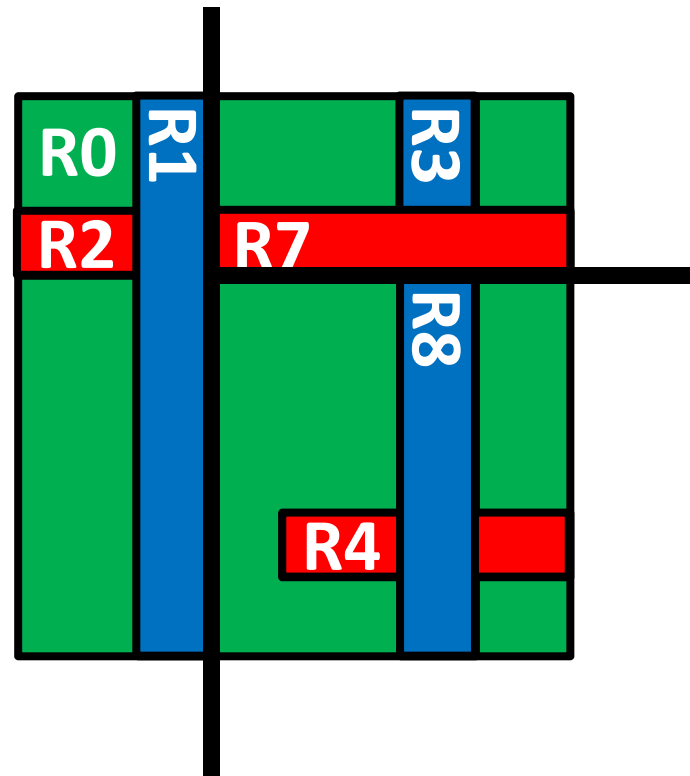


Challenges: Overlapping Rules



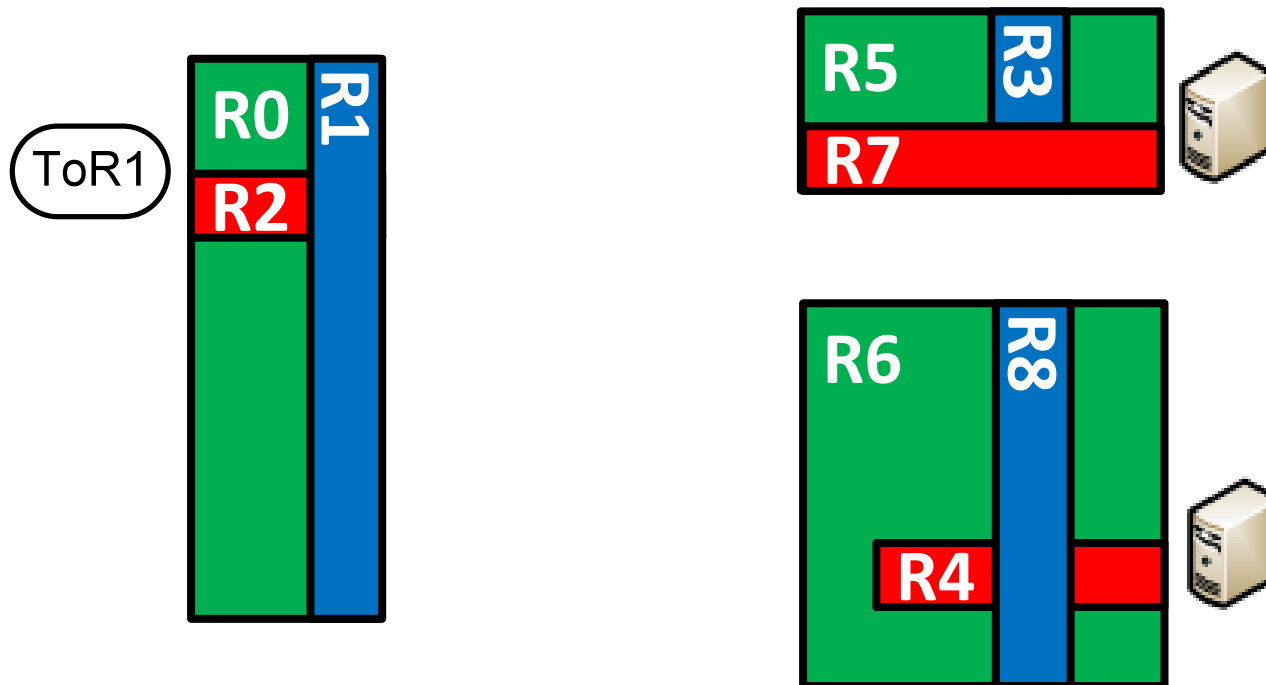
Challenges: Overlapping Rules

Partitions rules to reduce overlapping rules dependency



Challenges: Overlapping Rules

Partitions rules to reduce overlapping rules dependency



Splitting rules covering multiple partitions causes inflation

vCRIB: Partitioning

Recursively cut partitions to create a BSP tree

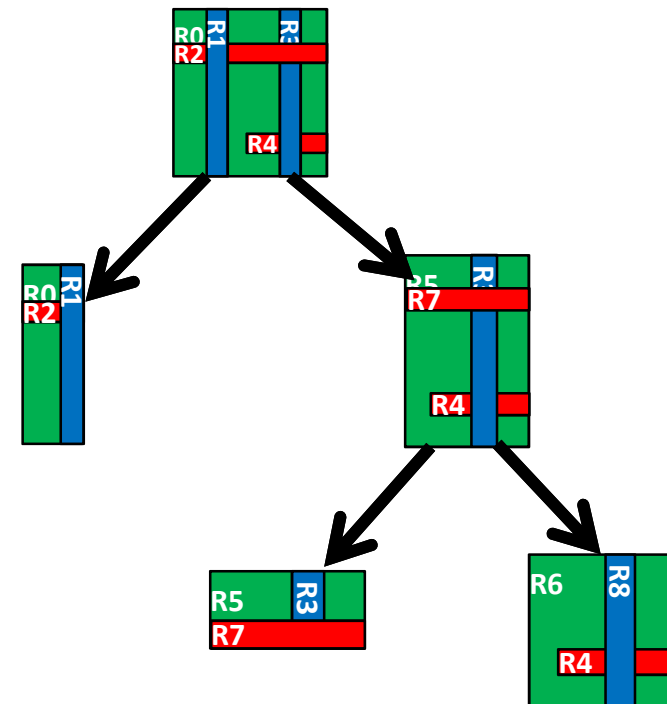
Select a cut that

- balances two partitions
- creates fewest number of new rules

Smaller partitions

- are more flexible to place
- match fewer communicating VMs

Stop whenever a resource at a node is exhausted



Challenges: Placement Complexity

Constraints

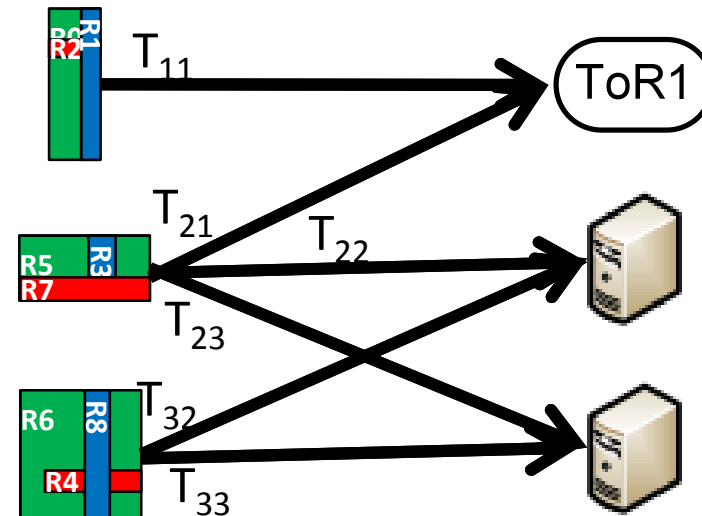
- Functionality
- Machine resources

Goal

- Minimize traffic overhead
- Minimize delay
- Minimize cost of bandwidth usage vs. saved CPU

- Different partition sizes
- Different machine capacities
- Different traffic overhead for each partition location

Generalized Assignment Problem



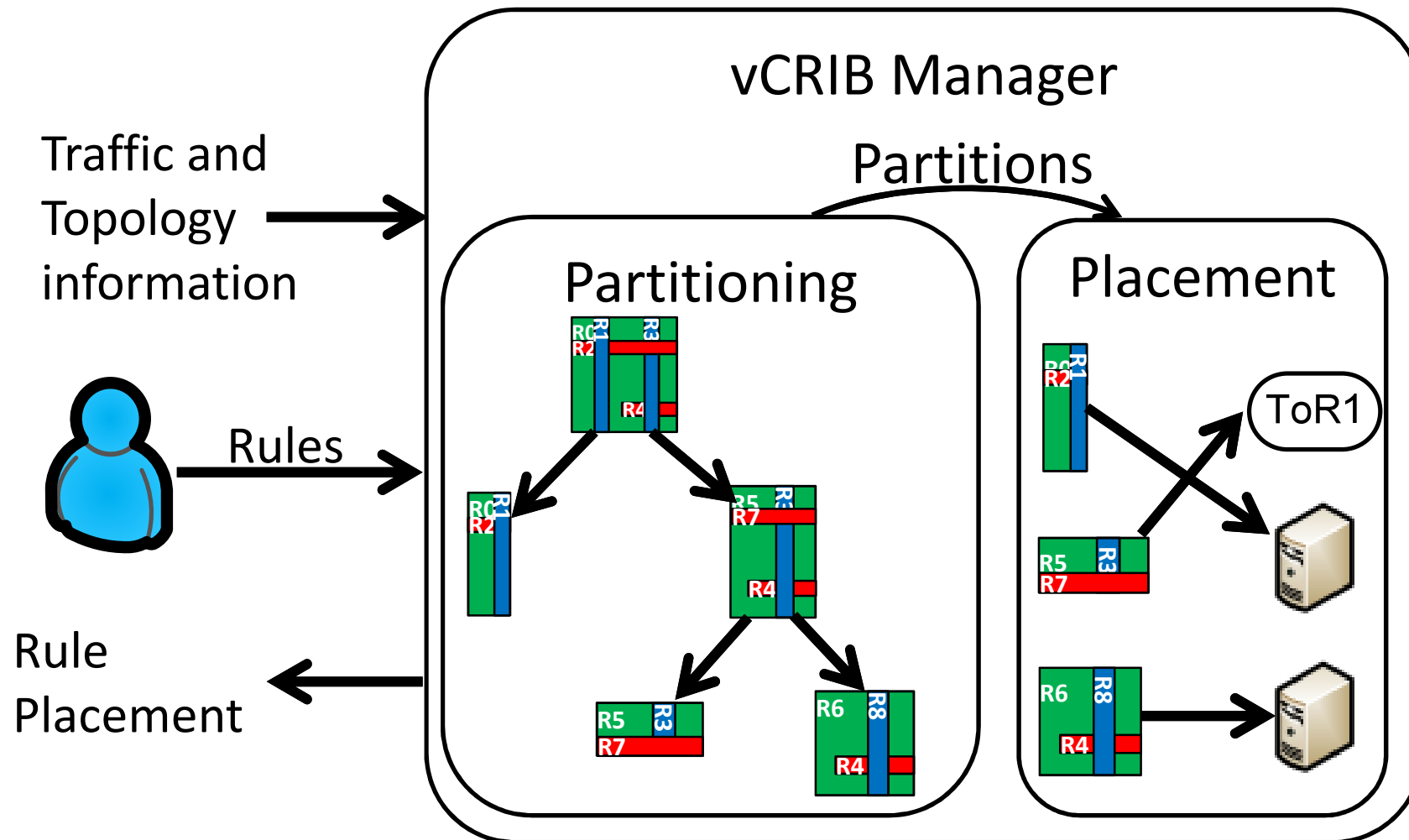
vCRIB: Placement (Branch and Bound)

Select the largest unassigned partition

Place it on a switch/hypervisor

- Capable of handling its rules
 - Functionality
 - Resources
- Make minimum traffic overhead

vCRIB Architecture



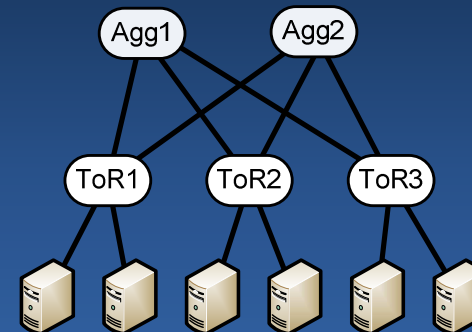
Evaluation: Goal

Can partitioning algorithm achieve small partitions?

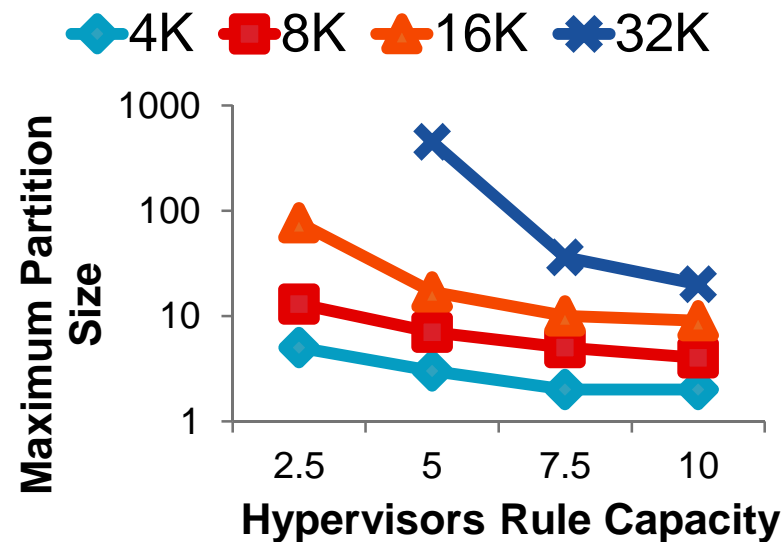
Can placement algorithm leverage resource availability to decrease traffic overhead?

Configuration

- 100 VMs per machine
- 10K flows (10KB) per machine
- ClassBench rules
- 1K rule capacity per switches



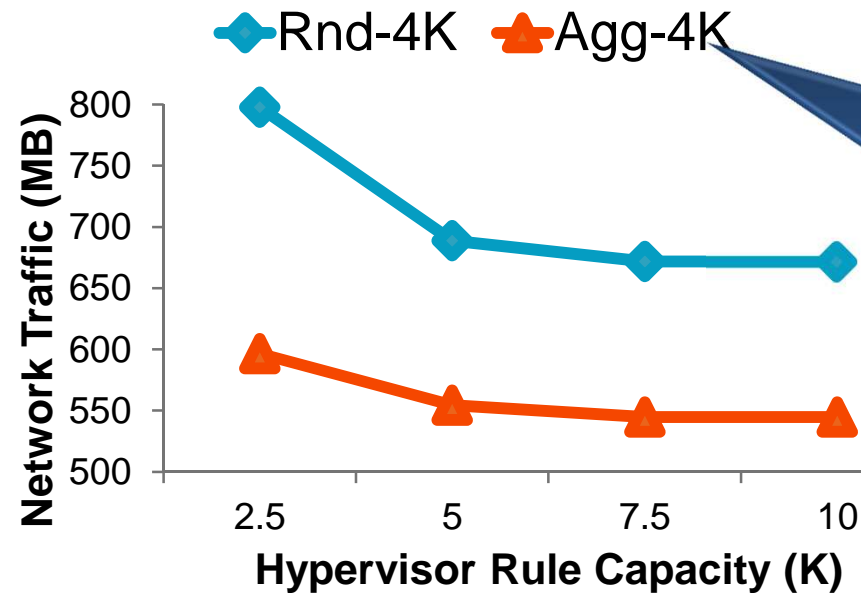
Evaluation: Partitioning



Change rule capacity to show the effect of different CPU budgets

Maximum size of partitions goes down as resources increase

Evaluation: Placement

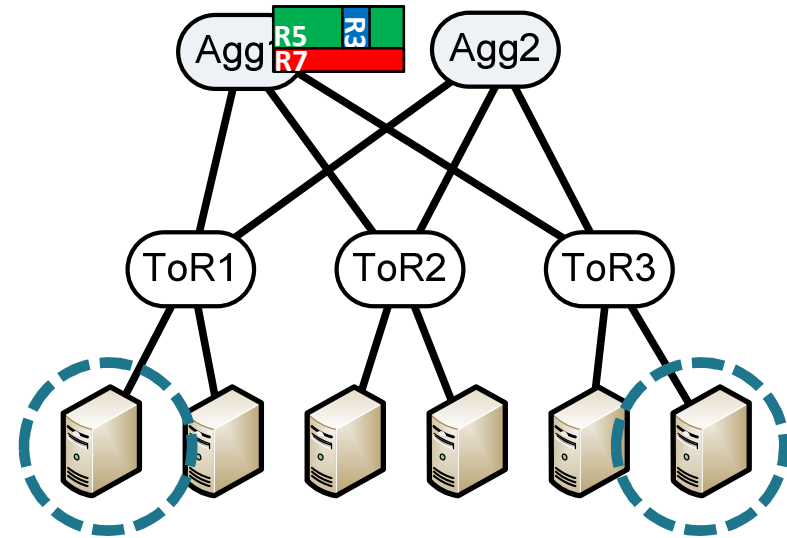
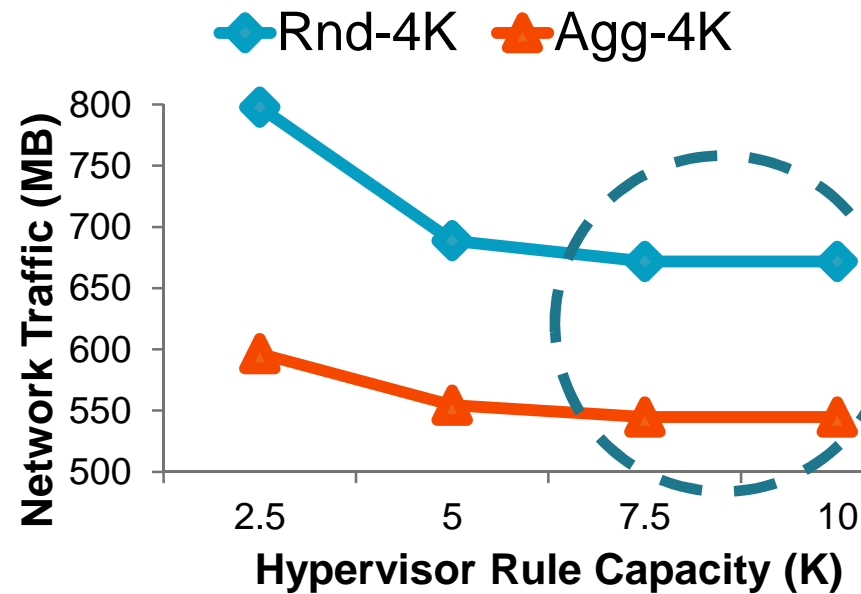


For each machine select VM addresses from a contiguous IP range

Traffic decreases as resources increase

Aggregated addresses make lower traffic overhead

Evaluation: Placement



No traffic decrease

Replication

Conclusion

vCRIB provides an abstraction layer for placement of rules in datacenters

Places the rules on both hypervisors and switches to achieve the best performance given the resource constraints

Future Work

Exploit performance model of hypervisors & switches

Online Algorithm adjusting to traffic changes

Replication in the partitioning and placement algorithm

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