# MemProf: a Memory Profiler for NUMA Multicore Systems

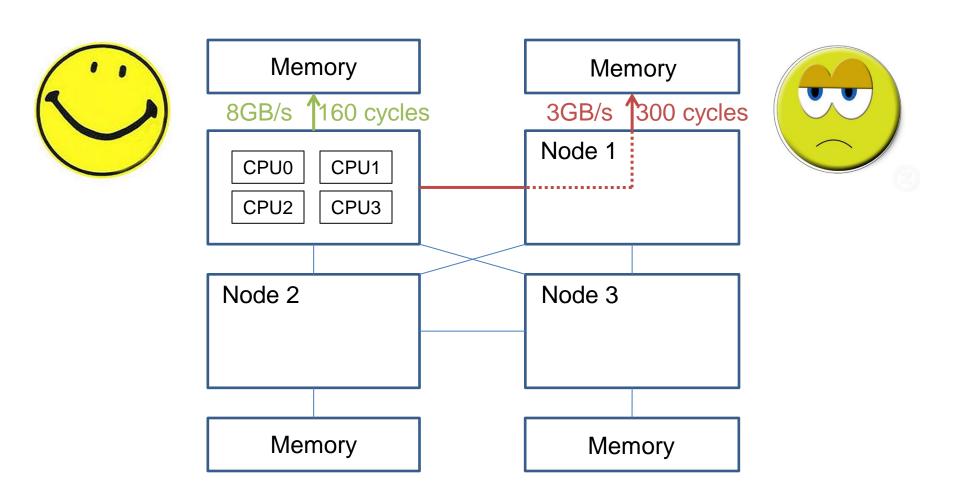
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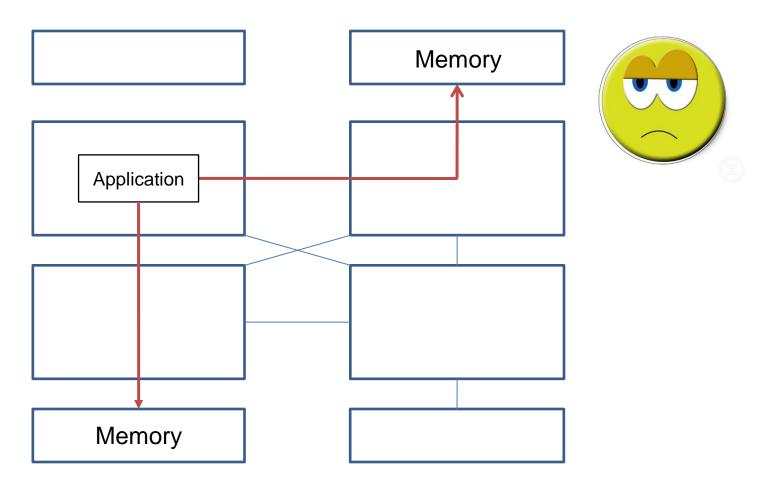




#### Machines are NUMA



#### Applications ignore NUMA



#### That is problematic

Application	% remote memory accesses in default version
FaceRec (ALPBench)	63%
Streamcluster (Parsec)	75%
Psearchy (Mosbench)	13%
Apache	75%

#### That is problematic

Application	% remote memory accesses in default version	% performance improvement provided by an adequate NUMA optimization
FaceRec (ALPBench)	63%	42%
Streamcluster (Parsec)	75%	161%
Psearchy (Mosbench)	13%	8.2%
Apache	75%	20%

#### Application-Agnostic Heuristics exist

- Thread scheduling and page migration (USENIX ATC'11)
- Thread Clustering (EuroSys'07)
- Page replication (ASPLOS'96)
- Etc.

# ... But they do not always improve performance

Example: Apache

	% remote memory accesses	% performance impact over default version
On default Linux	75%	<u>-</u>
With thread scheduling and migration (USENIX'11)	75%	-5%

# We want to understand the causes of remote memory accesses

# ... In order to select an adequate optimization

- Custom allocation policy
- Memory replication
- Memory migration
- Memory interleaving
- Custom thread scheduling policy

# Can we understand the causes of remote memory accesses using existing profilers?

### Let's take an example

### FaceRec (1)

- Facial recognition engine
- 63% of DRAM accesses are remote
- 42% gain when modified based on MemProf output

### Existing profilers point out

The functions that perform remote accesses

The memory pages that are remotely accessed

The global static objects that are remotely accessed

# Existing profilers point out (FaceRec)

- The functions that perform remote accesses
  - transposeMultiplyMatrixL = 98%
- The memory pages that are remotely accessed
  - 1/3 of the allocated pages
- The global static objects that are remotely accessed
  - No such object

#### What can we conclude?

- Should we change the allocation policy?
  - No idea
- Should we migrate memory pages?
  - No idea
- Should we replicate memory pages?
  - No idea
- Etc.



#### So... We need a new profiler!

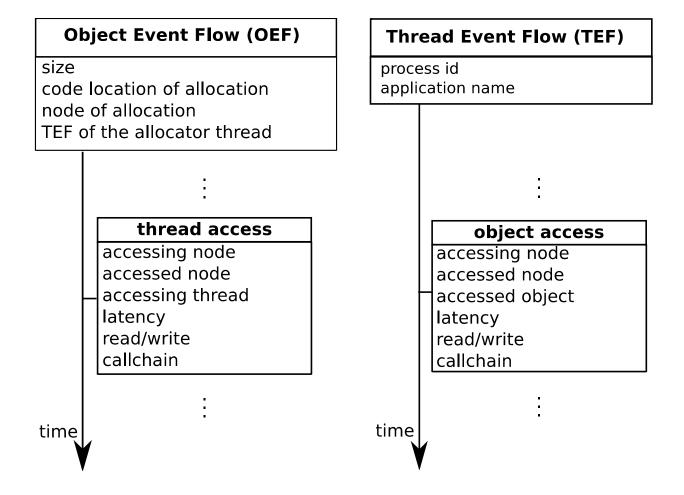
# We designed MemProf, a profiler that points out

- Remotely accessed objects
- Thread-Object interaction patterns

#### Objects

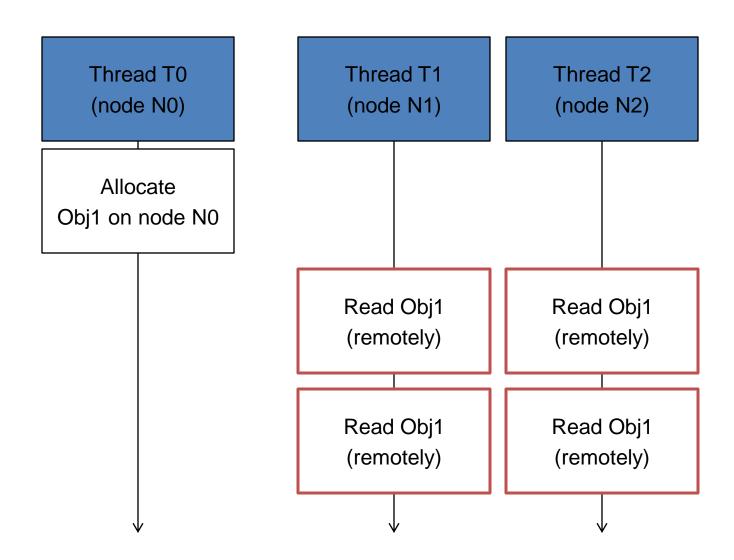
- Global statically allocated objects
- Dynamically allocated objects
- Memory-mapped files
- Code sections mapped by the OS
- Thread stacks

### Thread-Object interaction patterns

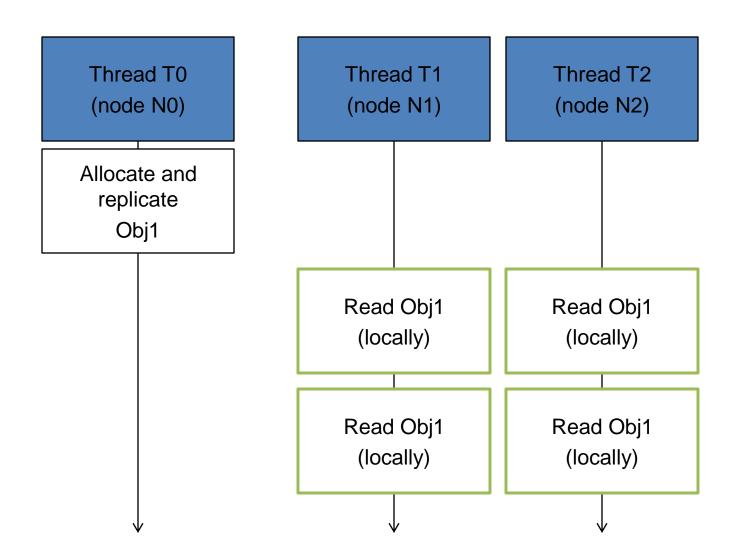


#### What can we do with MemProf?

### We can detect that an object is simultaneously read by several remote threads...



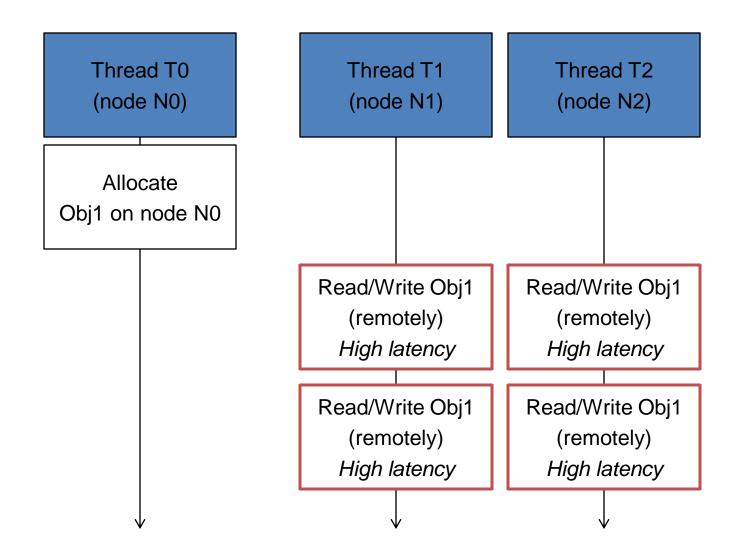
#### And thus decide to replicate this object on several nodes



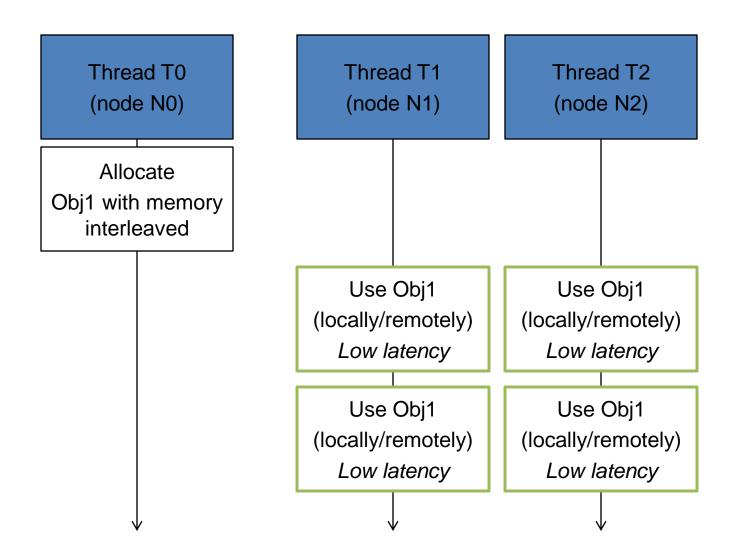
## This is the pattern observed in FaceRec

- 193 matrices
- 1 matrix induces 98% of the remote accesses
- This matrix is first written and then read by all threads
- We replicate the matrix (10 lines of code)
- Performance improvement: 42%

### We can detect that an object is simultaneously read and written by several threads with a high latency



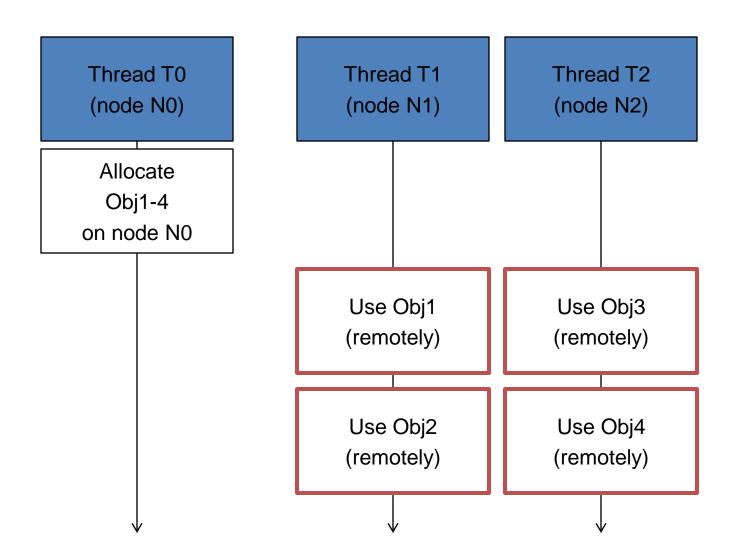
#### And thus decide to interleave this object



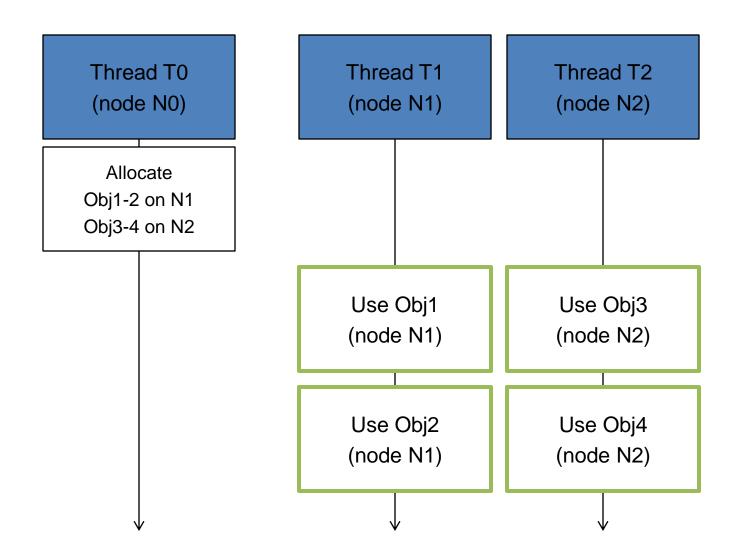
### This is the pattern observed in Streamcluster

- 1000 objects allocated
- 1 represents 80% of remote memory accesses
- It is accessed read/write by all threads
- We interleave this object (1 line of code)
- Performance improvement: 161%

#### We can detect that threads do not share objects



#### And thus decide to change the allocation policy



# This is the pattern observed in Psearchy

- Remote accesses are done on private variables
- We forced local allocations (2 lines of code)
- Performance improvement: 8.2%

#### As a summary

- MemProf allows finding memory access patterns
- Knowing memory access patterns allows designing simple and efficient optimizations

### A word on the implementation

### MemProf – Online Profiling

- Memory access tracking
  - IBS samples
- Object lifecycle tracking
  - Overloading of allocation functions
  - Kernel hooks
- Threads lifecycle tracking
  - Kernel hooks

#### MemProf – Offline Analysis

- Sort samples by time
- Match memory addresses with objects
  - Leverages object lifecycle tracking
  - Leverages thread lifecycle tracking
- Create object-thread interaction flows
  - Leverages thread lifecycle tracking

#### Overhead

- 5% slowdown
- 2 sources of overhead:
  - IBS sampling collection: one interrupt every 20K cycles
  - Object lifecycle tracking

#### Conclusion

- Remote memory accesses are a major source of inefficiency
- Existing profilers do not pinpoint the causes of remote memory accesses
- We propose MemProf, a memory profiler that allows:
  - Finding which objects are accessed remotely
  - Understanding the memory access patterns to these objects
- Using MemProf, we profiled and optimized 4 applications on 3 machines
  - Optimizations are simple: less than 10 lines of code
  - Optimizations are efficient: up to 161% improvement

