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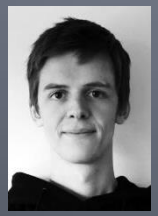
REYKJAVIK UNIVERSITY

Microsoft Research

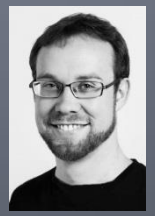
harmonium

Elastic Cloud Storage via File Motifs

Helgi Kr. Sigurbjarnarson **RU**



Pétur O. Ragnarsson **RU**



Ymir Vigfusson **RU**

Mahesh Balakrishnan **MSR**

Elasticity

VM

VM

VM

Host cannot reclaim space from virtual disk

Virtual disk cannot expand beyond its fixed capacity

Hardware

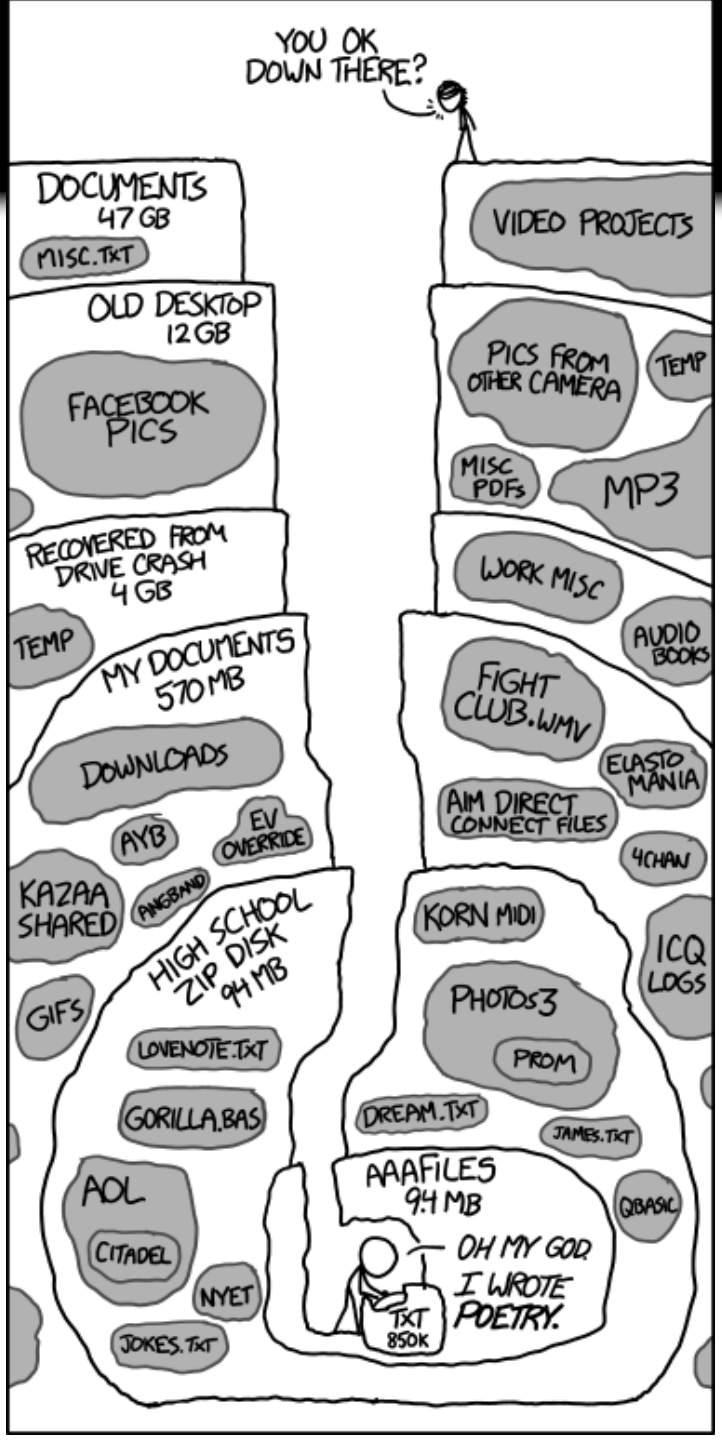
RAM

CPU

Storage

No storage elasticity

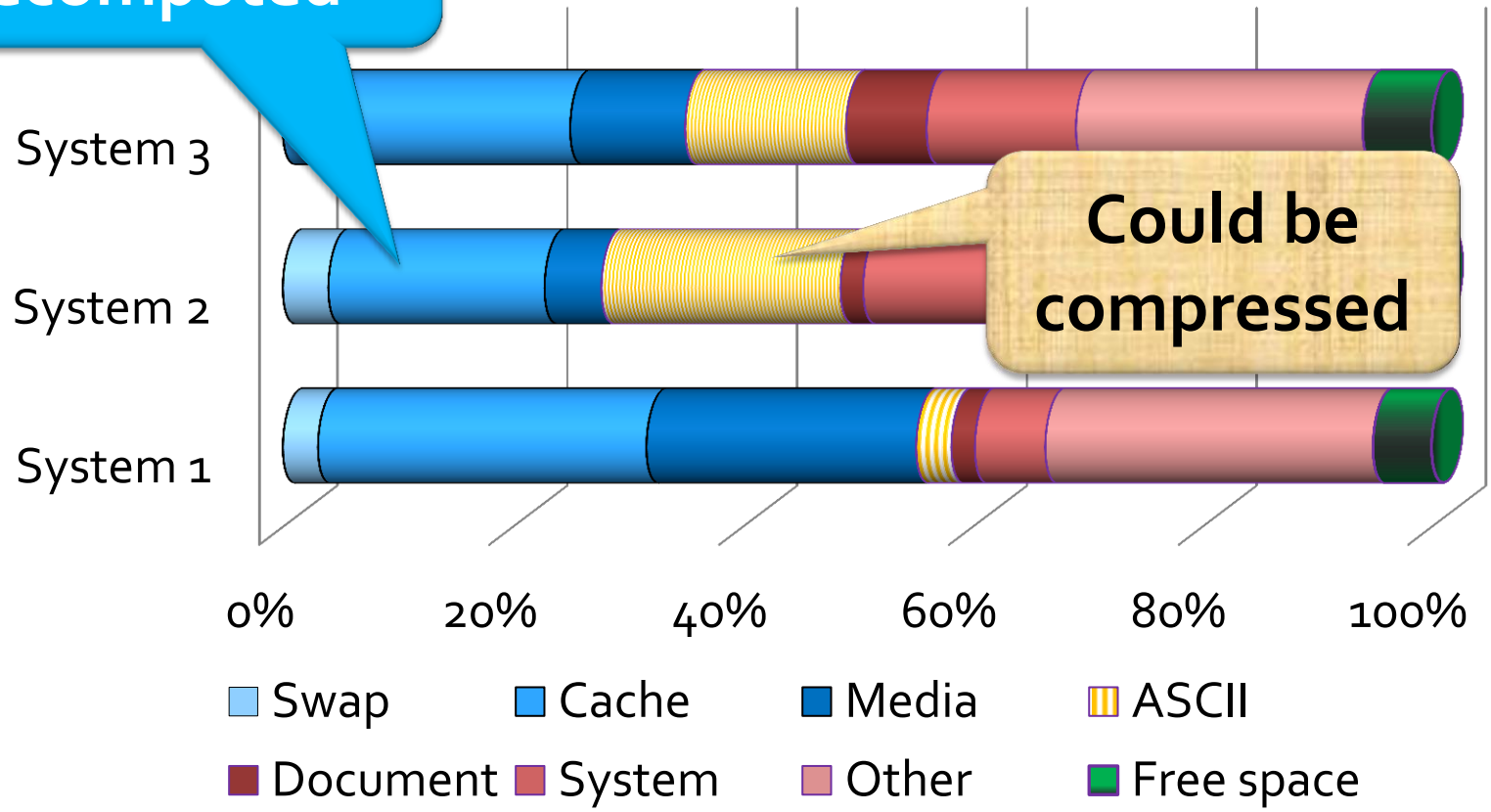
http://xkcd.com/1360/



Not all data requires durability

Could be recomputed

ccupied by file types (3 developer laptops)

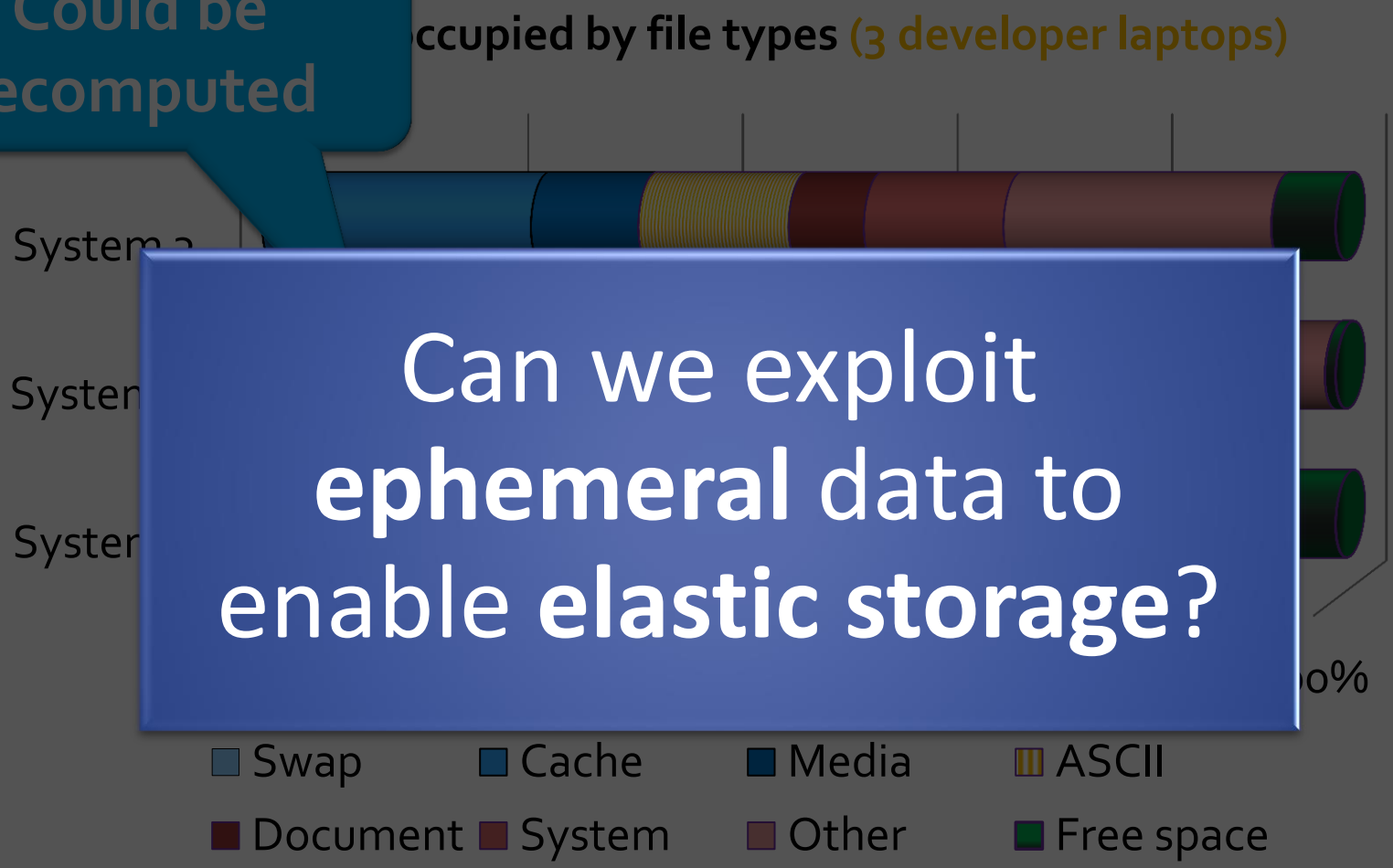


- **Significant ephemeral data**

Representative?

Not all data requires durability

Could be recomputed



Can we exploit **ephemeral** data to enable **elastic storage**?

- **Significant ephemeral data**

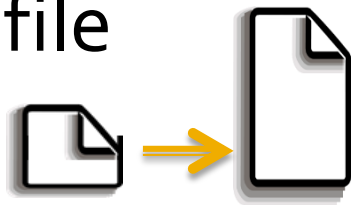
Representative?

What if files could **contract** and **expand**?

The motif abstraction

A **motif** is **code** to generate the **data** in a file

- **expand()** obtains the raw bytes of the file
  - Run **computations** (e.g. **compression**)
  - **Fetch data** across a network
  - Operate over **other files** on the FS

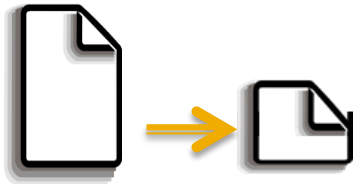


- **contract()** deletes raw bytes, retains motif code

Motifs can be recursive

Motifs can have circular dependencies

Files can have multiple motifs

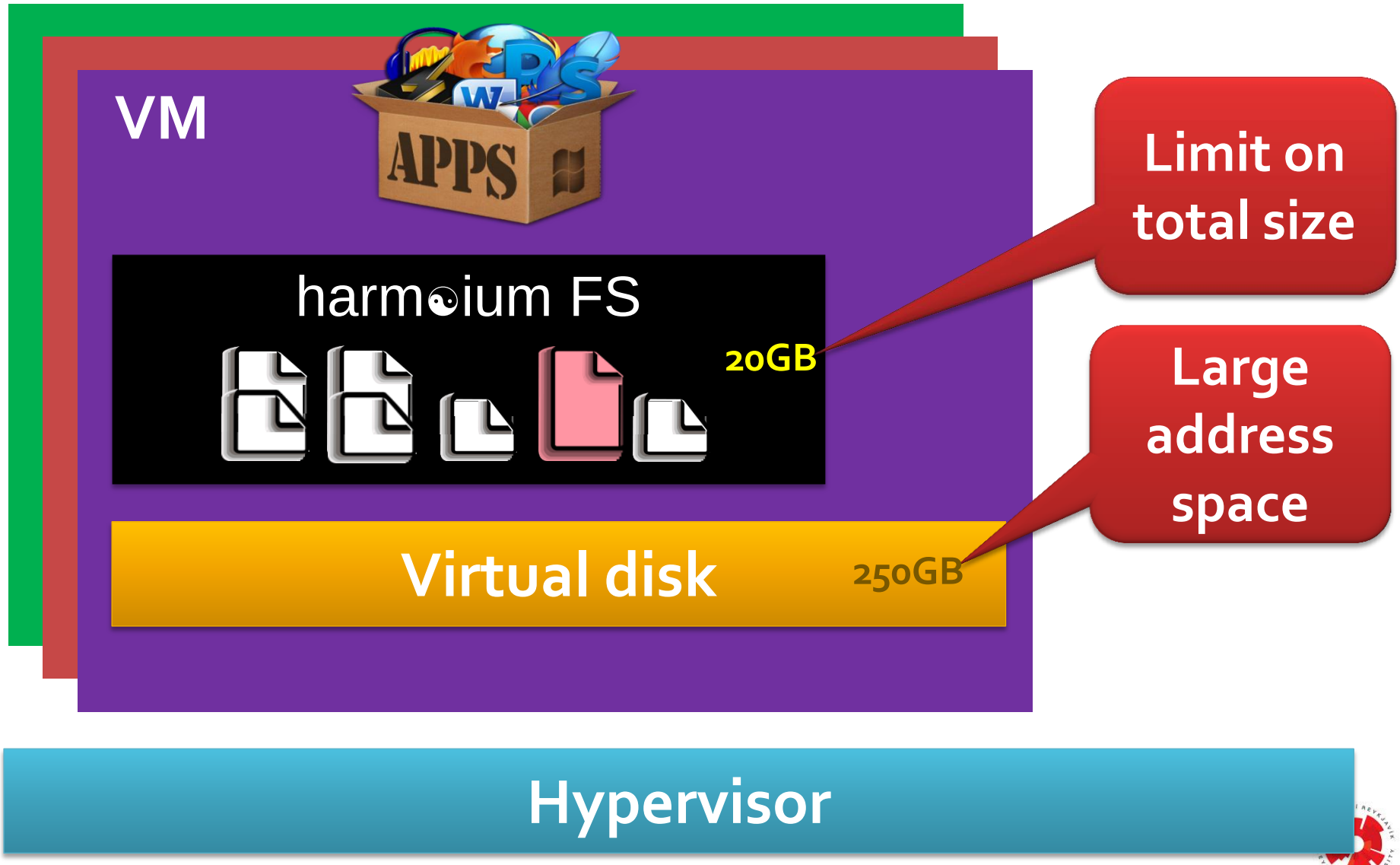


Example *network storage* motif

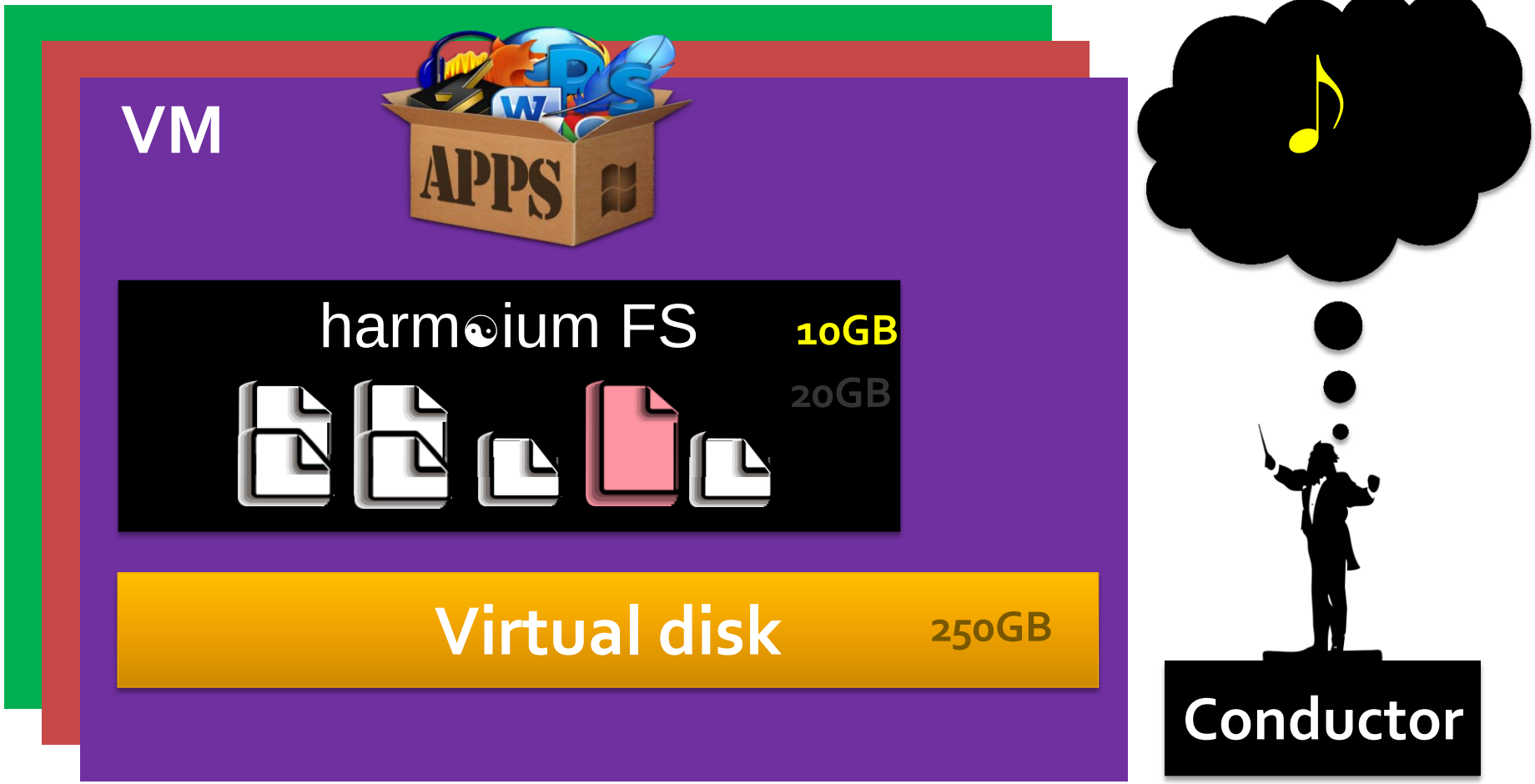
```
class SCPMotif(object):
  def expand(self, fname, meta=None):
    p = fname.bypass()
    os.popen('scp \
       fileserver1:storage%s "%s"' % (p,p))

  def contract(self, fname, meta=None):
    p = fname.bypass()
    os.popen('ssh fileserver1\
      "mkdir -p storage%s"' %\
      os.path.dirname(p))
    if os.popen('scp "%s"\
       fileserver1:storage%s' % (p,p)) == 0:
      open(p, 'w').close()
```

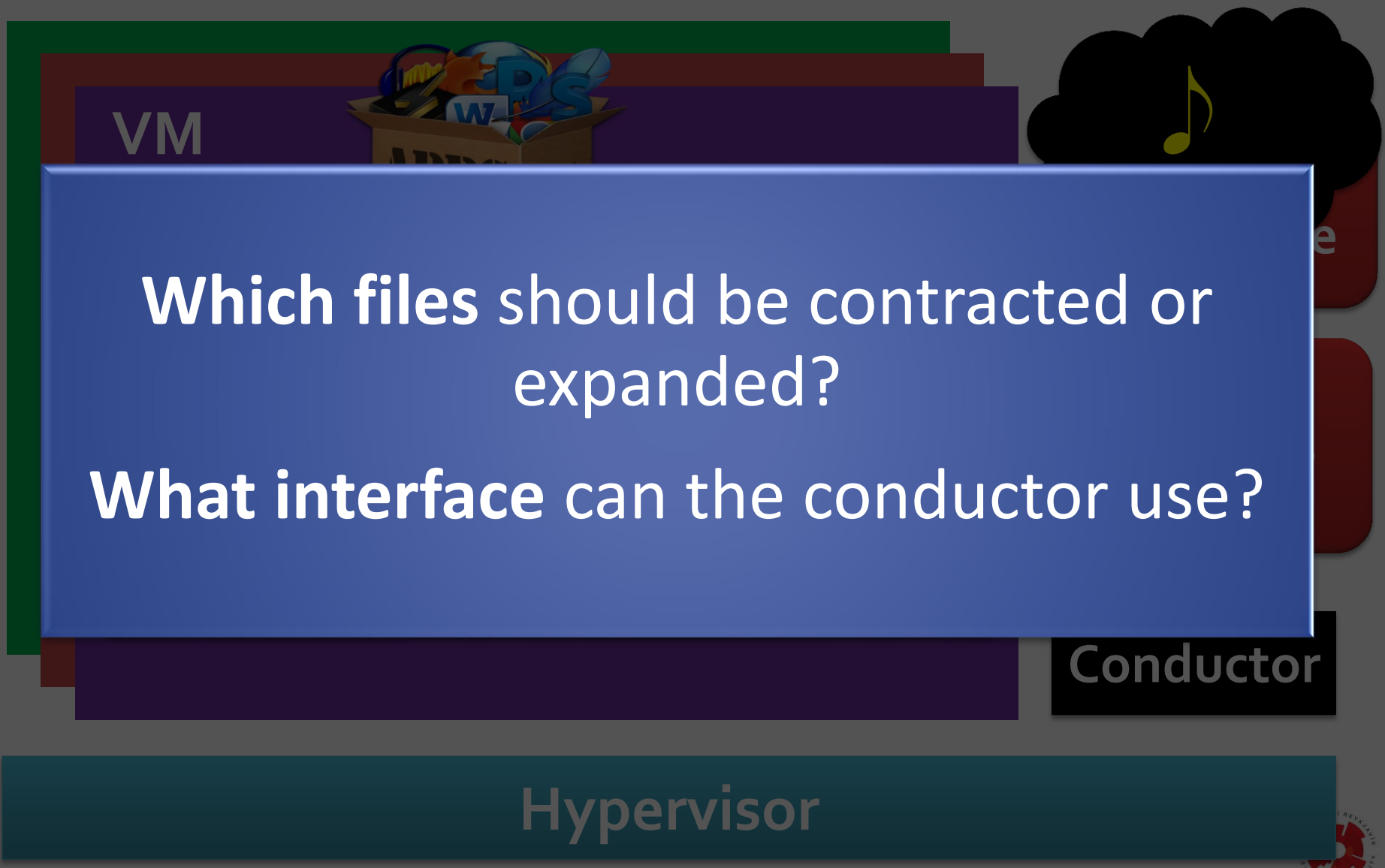
The harmonium file system



The harmonium file system



The harmonium file system



Which files should be contracted or expanded?

What interface can the conductor use?

What files to contract/expand?

harmonium FS

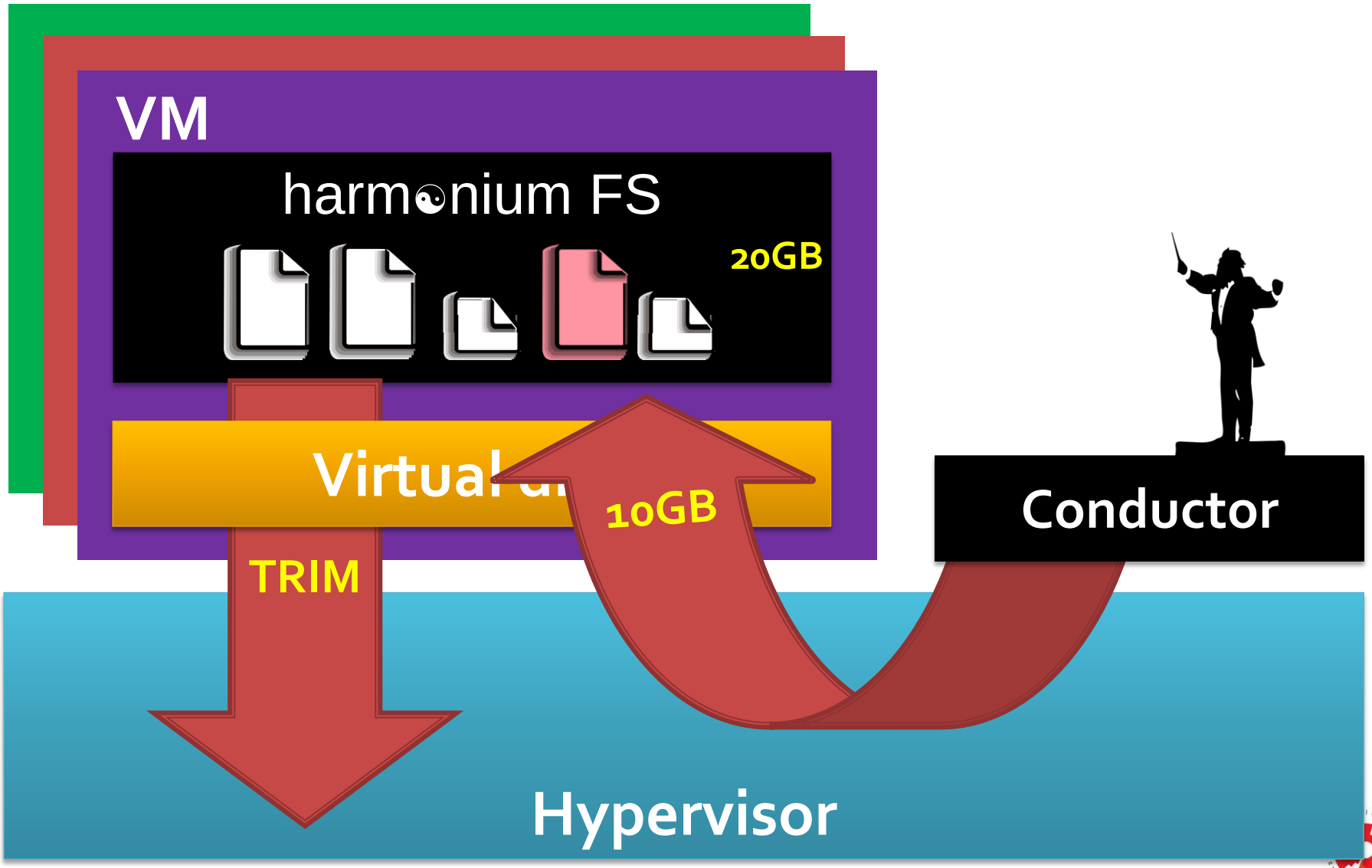
x GB limit

When will the file be accessed next?

How much time will it take to expand?

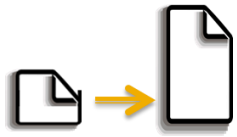
- Feed info into optimization mechanism
  - Choose files with minimum total expansion latency s.t. contraction saves sufficient space.
  - Most competitive algorithm: Greedily choose files on LRU list to maximize ratio of space savings to expansion latency

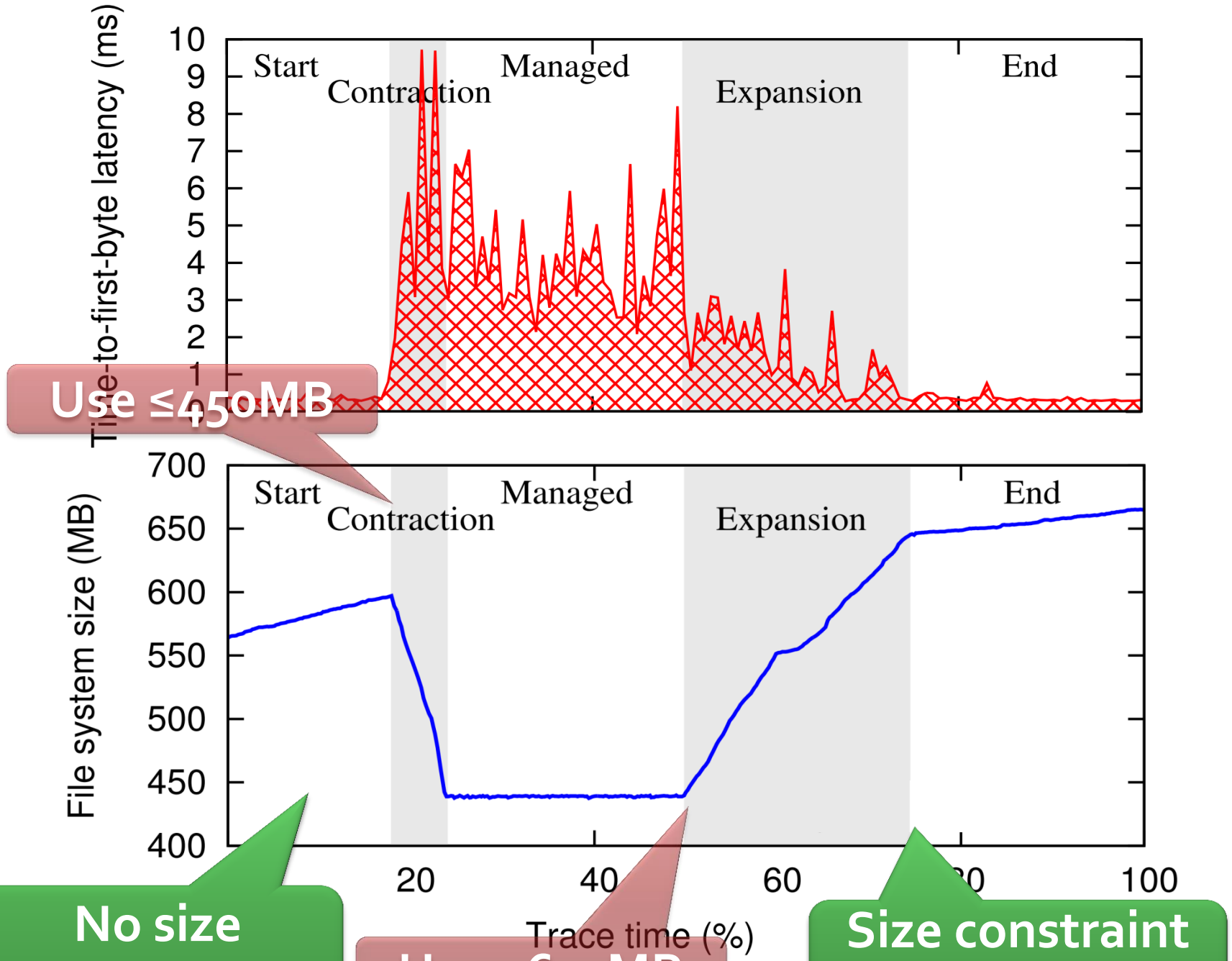
Conductor interface



Evaluation

- User-space FUSE prototype
  - Conductor: Python, via UNIX sockets
- Workload
  - Set of 54,000 patch files applied in chronological order to the Linux kernel

- Motif 
  - Network storage via scp
- Measure latency to access first byte


Time-to-first-byte latency (ms)
10
9
8
7
6
5
4
3
2
1
Start
Contraction
Managed
Expansion
End
Use ≤450MB
File system size (MB)
700
650
600
550
500
450
400
Start
Contraction
Managed
Expansion
End
20
40
60
100
Trace time (%)
No size constraints
Use ≤650MB
Size constraint lifted

Conclusion

Elastic performance/capacity trade-off for storage in VMs

Problem

Applications store **ephemeral** data on secondary storage

But storage stacks provide **durability** for every file

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Associate ***motifs*** with every file, allowing reconstruction

Contract/expand files to minimize access latency

Results

Best algorithm: LRU **greedily** maximizing space/latency

Fully functional **FUSE** prototype

Extra slides

Theoretical formulation

- **0-1 Knapsack (NP-complete)**
  - S = space needed to save
  - e_i = expected expansion latency for file *i*
  - s_i = expected storage savings for file *I*

$$\min \quad \sum_{i=1}^{n} e_i x_i$$

$$\text{s.t.} \quad \sum_{i=1}^{n} s_i x_i \geq S,$$

$$x_i \in \{0,1\} \quad , \quad 1 \leq i \leq n$$

- **Reduction:** Choose the files ***not*** to include
- **APX-knapsack** takes ***O(nW)*** time, prohibitive

Related work

- Trade-off between storage footprint and performance
  - Usually in distributed settings
  - Sierra (EuroSys 2011), Rabbit (SOCC 2011), Springfs (FAST 2014), …
- These systems maintain 1 to N copies of each file
- harm☯ium, however, maintains "0 to 1" copies of each file

But isn't this just ...

- A compression file system?
- A glorified cache?
- A de-duplication system?
  - Harm☯ium can support arbitrary motifs:
    - Compression of rarely accessed files
    - Remote network storage [scp/rsync/nfs/...]
    - Pull from repositories [git/svn/...]
    - Re-wget data in Downloads folder
    - Resume torrent download [remove partial files]
    - System packages [retrieve from apt/debian/...]
    - Regeneration of data sets [ala Nectar]
    - ...

Security concerns

- Motifs are really arbitrary code
  - Can cause system to hang, crash, corrupt data or consume resources wastefully
- Our current implementation is vulnerable
  - Motifs execute within the same process as the FS
  - Isolation by virtualization too coarse-grained
  - Sandboxing great for security, may be slow
- Ongoing work: require authorization
  - Users specifically approve running of motifs generated by those of lesser privilege or fewer capabilities

Computation vs. storage

- *"But isn't computation more expensive than storage?"*
- Underlying principle of our work:

Computation, Network and Storage are fungible

- Harmonium allows use of Computation or Network when Storage is scarce
- Other parts of the trade-off interesting as well!