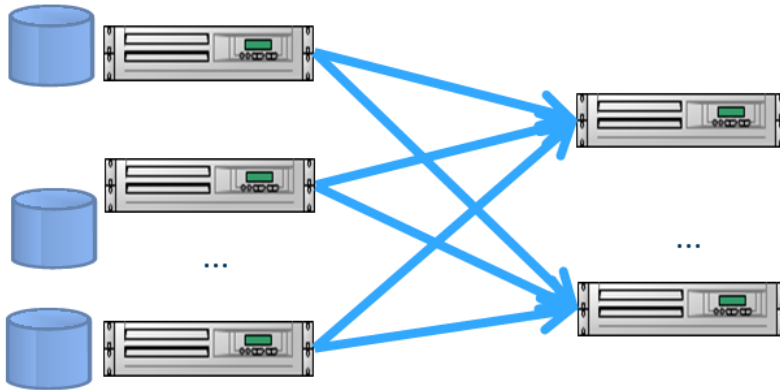


Rhea: Automatic Filtering for Unstructured Cloud Storage

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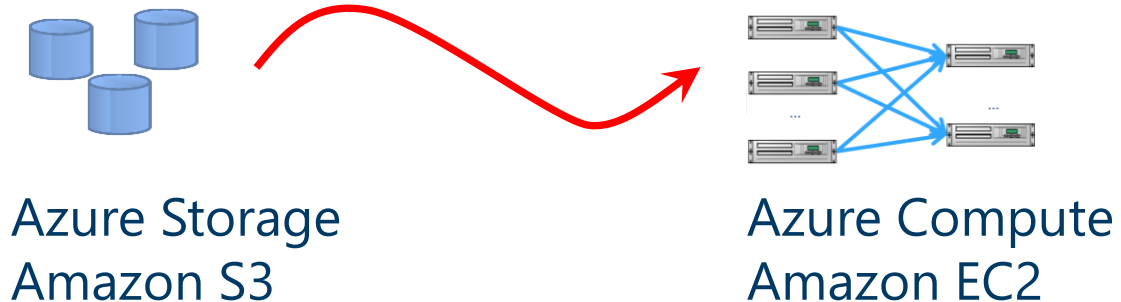
Cluster design for data analytics: [Traditional] Collocate storage & compute



- Hadoop & MapReduce, Dryad/DryadLinq, Scope, etc

Cloud Analytics: Hadoop in the Cloud

Separate storage and compute

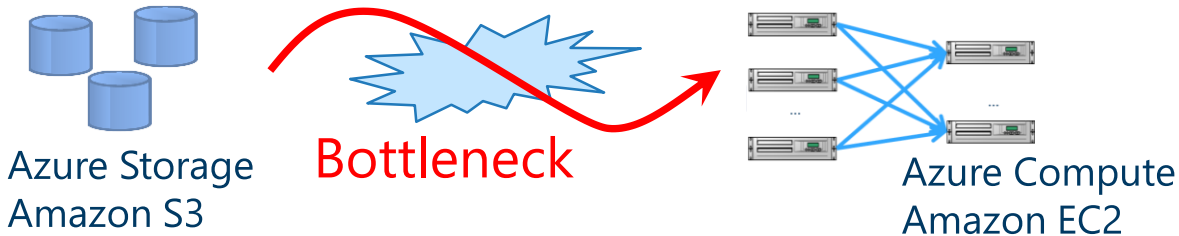


Examples:

- Hadoop on Azure
- Amazon's Elastic MapReduce

Cloud Analytics: Hadoop in the Cloud

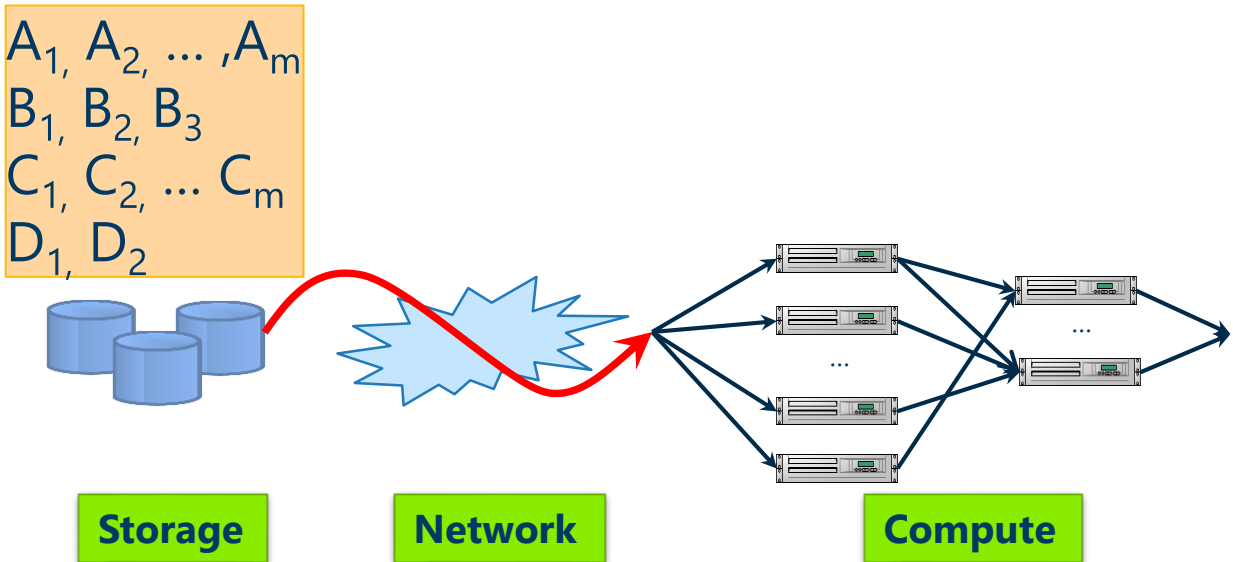
Separate storage and compute



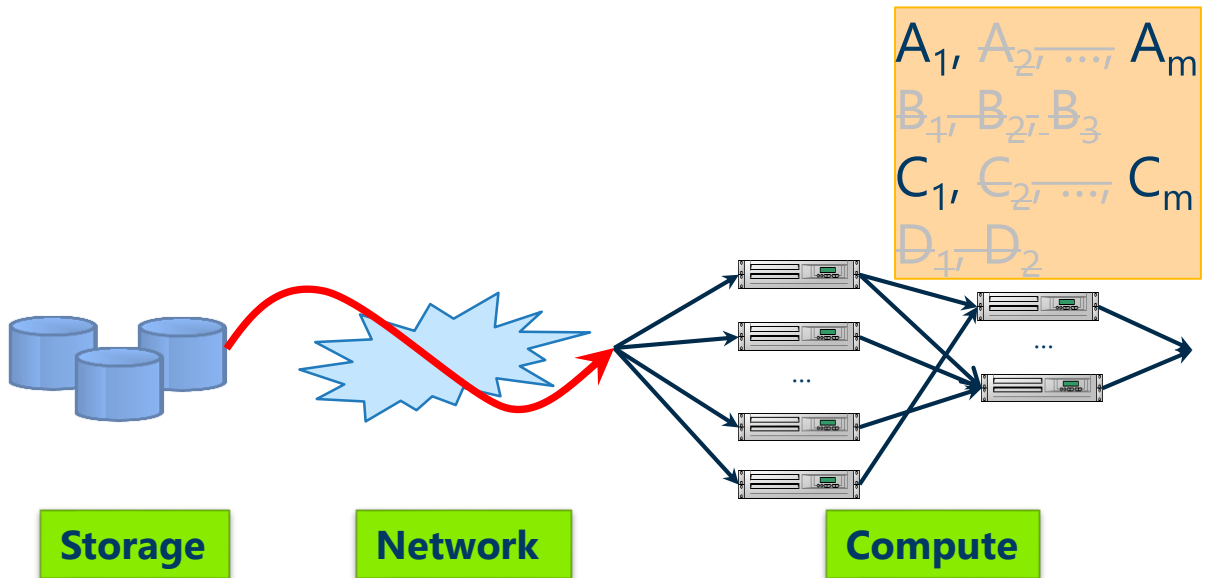
Why separate storage from compute?

- + (User) Don't pay for compute just to keep data alive
- + (User) Offload storage management to operator
- + (Operator) Evolve compute & storage independently
- + (Operator) Offer services that do not require both
- **Network between storage and compute is limited**
(see paper for details)

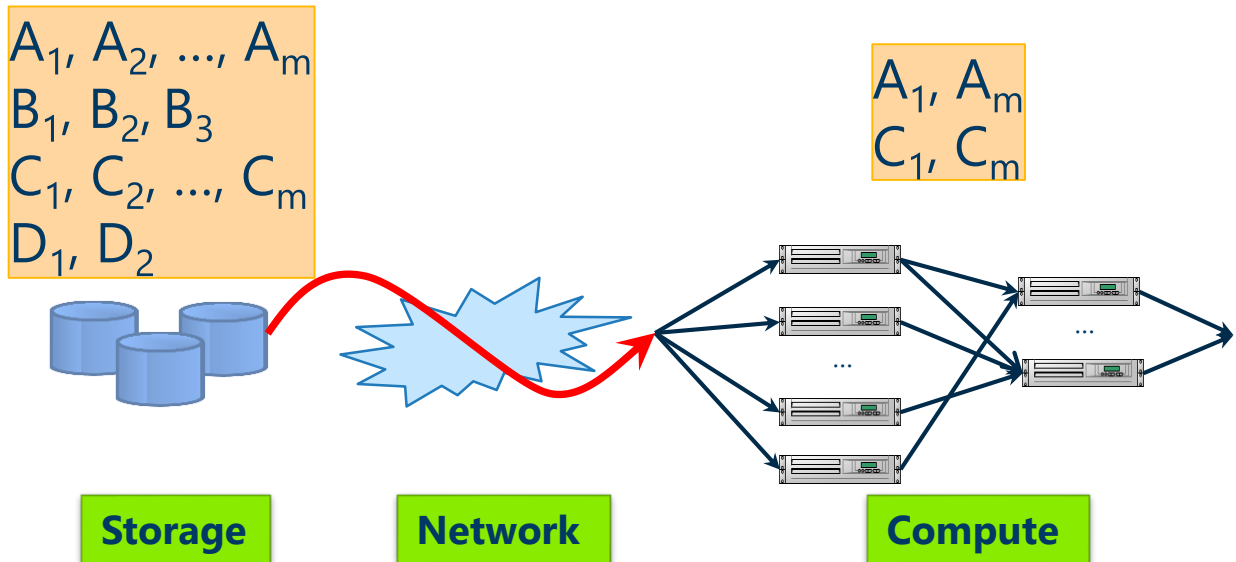
Problem: Transfer lots of data ...



Problem: Transfer lots of data ...
... even when only a subset is needed



Problem: Transfer lots of data ...
... even when only a subset is needed



Scenario

- Apache Hadoop (Map/Reduce)
- Input data in storage service
- Hadoop running in compute service
- Unstructured data:
 - text, log files, etc

Goal

Transparently reduce data transfers
from storage to compute

How to minimize transfers?

- Strawman: Can we execute mappers on storage nodes?
 - Intuition: Mappers throw away a lot of data
 - ☹ Data reduction not guaranteed
 - ☹ Difficult to stop mappers during storage overload
 - ☹ Storage nodes have to execute complicated logic (Hadoop system & protocol)
 - ☹ Dependencies on runtime environment, libraries, etc
- Better approach: Filter unnecessary data at storage nodes
 - Filters need to be **opportunistic and transparent**
i.e. can kill/restart at any time (e.g. during overload)
 - Filters need to be **correct**
i.e. always preserve correctness of computation

Challenge: How to filter the data?

Recall: data are typically unstructured text

- No external source of structure/schema

Insight:

- The data analytic job knows structure
- ... and what needs to be filtered

Idea: static analysis of job bytecode

```
public void map(... value ...)
```

Input Value

```
{  
String[] entries = value.toString().split("\t");  
String articleName = entries[0];  
String pointType = entries[1];  
String geoPoint = entries[2];
```

Projection operation

- 3 "columns" interesting (out of 4 for this job)

```
if (GEO_RSS_URI.equals(pointType)) {
```

Selection operation

- roughly 1/3 of rows are of the interesting type

```
StringTokenizer st = new StringTokenizer(geoPoint);  
String strLat = st.nextToken();  
String strLong = st.nextToken();  
double lat = Double.parseDouble(strLat);  
double lang = Double.parseDouble(strLong);  
String locationKey = .....  
String locationName = .....  
geoLocationKey.set(locationKey);  
geoLocationName.set(locationName);
```

"selects"/"projects"
implicit in Java byte code

```
outputCollector.collect(geoLocationKey, geoLocationName);
```

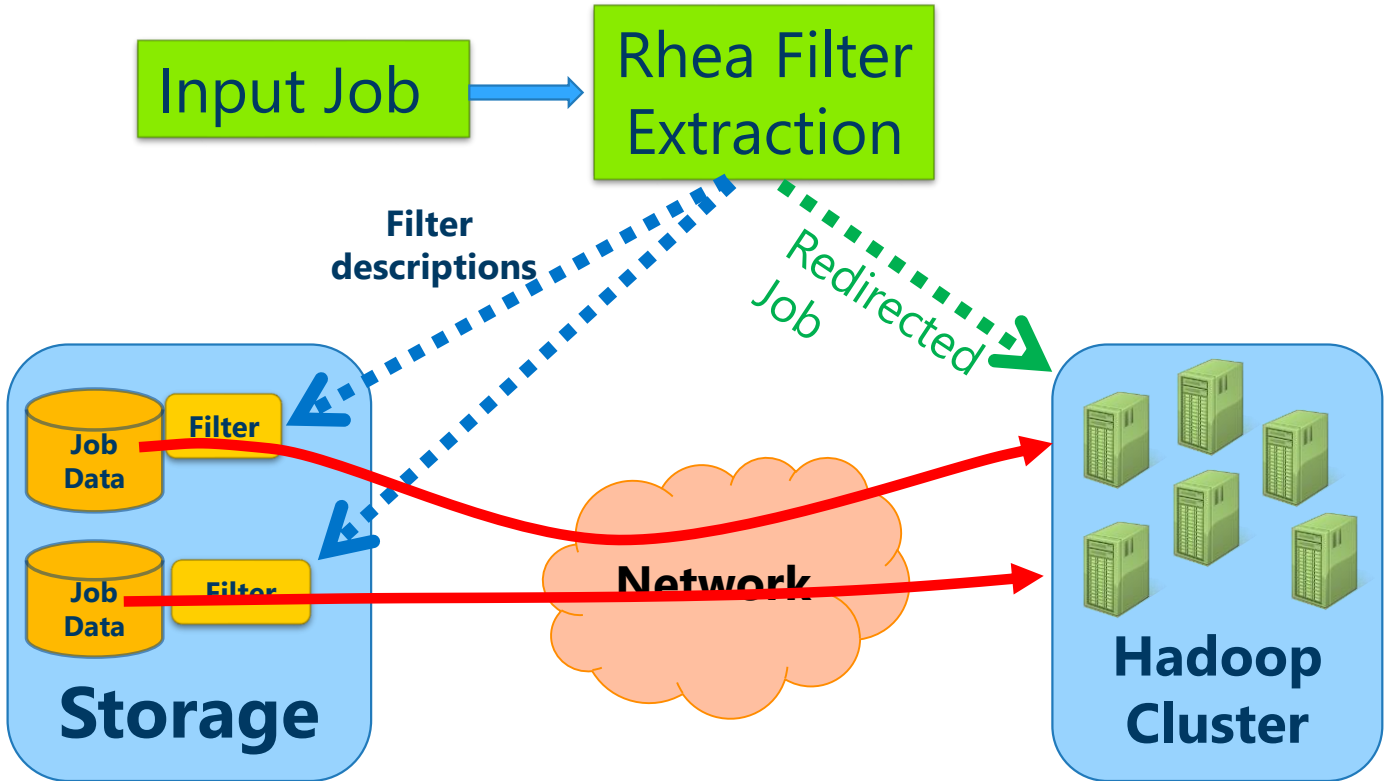
Output operation

```
}}
```

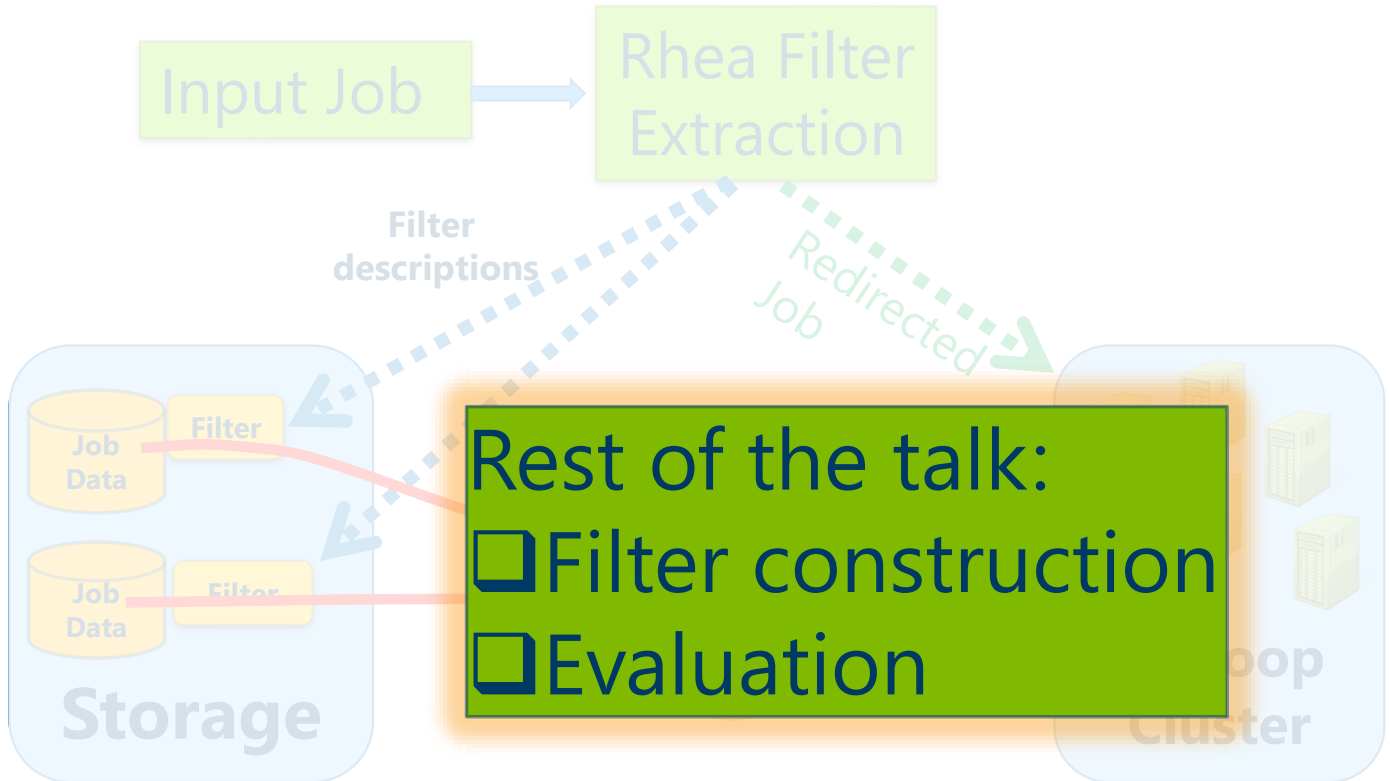
Rhea

- Static analysis of Java byte code
- Extract row (select) & column (project) filters
 - as **executable** Java methods
 - column filters can also be C, regular expressions, etc.
- Filters are **conservative**:
 - May accept more data than strictly necessary
- Filters are **opportunistic**
 - kill/restart at any time (e.g. during storage overload)
- Filters are **transparent**
 - no change to Hadoop job

Rhea's Architecture



Rhea's Architecture



Filters: Identify bits of data that affect output of mapper

- Row Filters:

- Given an input row:

Does it lead to output?

- Row corresponds to one invocation of map

- Approach: Path Slicing

- Challenge: Deal with mutable state

- Column Filters:

- Given a row that leads to output:

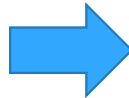
Which substrings of the row affect output?

- Approach: Abstract interpretation

- Challenge: Deal with loops

Row Filter Generation via Path Slicing

```
public void map(... value ...)  
{  
    String[] entries = value.toString().split("\\t");  
    String articleName = entries[0];  
    String pointType = entries[1];  
    String geoPoint = entries[2];  
  
    if (GEO_RSS_URI.equals(pointType)) {  
        StringTokenizer st = new  
            StringTokenizer(geoPoint, " ");  
        String strLat = st.nextToken();  
        String strLong = st.nextToken();  
        double lat = Double.parseDouble(strLat);  
        double lang = Double.parseDouble(strLong);  
        String locationKey = .....  
        String locationName = .....  
        geoLocationKey.set(locationKey);  
        geoLocationName.set(locationName);  
        outputCollector.collect(geoLocationKey,  
            geoLocationName);  
    }  
}
```



1. Tag "**observable**" instructions
2. Identify **path conditions** that lead to observable instructions
3. Perform dataflow analysis to identify all instructions that affect path conditions
4. Emit code



```
public boolean filter(Text bcvar2) {  
    String[] bcvar5 = bcvar2.toString().split("\\t");  
    String bcvar7 = bcvar5[1];  
    boolean irvar0_1 =  
        GEO_RSS_URI.equals(bcvar7);  
    if (irvar0_1 == 1) { return true; }  
    return false;  
}
```


Challenge: Taming State

- Map-Reduce programs are often NOT pure functions
 - ➔ M/R programmers use state (i.e. objects in heap):
 - ... to avoid frequent initializations
 - ... to pass job parameters
 - ... to optimize temporary storage (e.g. with dictionaries)
- Filters cannot rely on mutable state:
 - Recall: output of filtered data = output of original data
- Solution: Tag all access to mutable fields as “observable” (i.e. output) instructions.

Column Filter Generation (aka projects)

- Goal: Identify substrings that affect output
- Based on *abstract interpretation*
 - Captures common patterns for “reading” fields: e.g. string tokenizers, regular expressions, etc.
 - Guarantees termination by using numerical constraints
 - Important to deal with loops
- Output:
 - Tokenization method and separator
 - List of indices of interesting tokens

Rest of the talk:

Filter construction

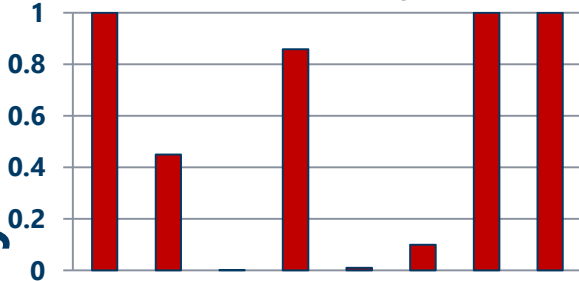
Evaluation

Experimental setup

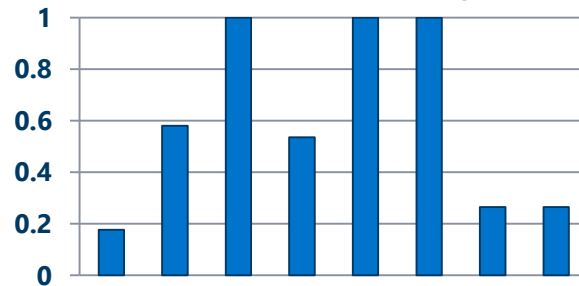
- Hadoop on Azure:
- Input data in Azure Storage
- Compute on Azure Compute
- 8 jobs with both code and data
 - 200 jobs code only (in paper)
- Same data-center
 - Also, cross data-center (in paper)

Job Selectivity

Rows Only



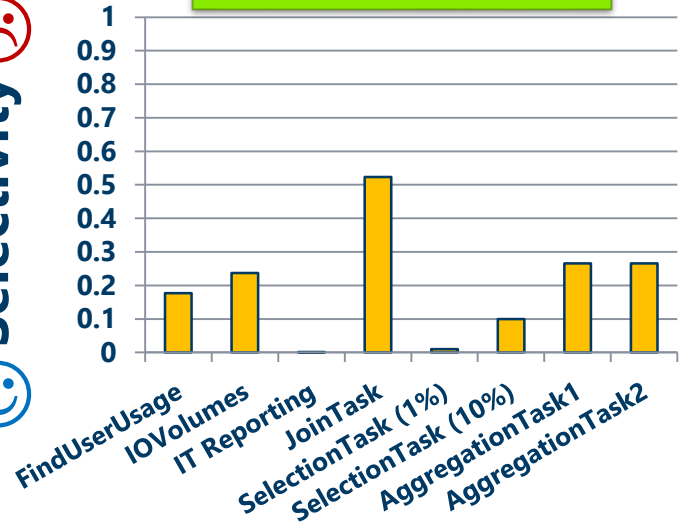
Columns Only



Selectivity



Rows + Columns



Many jobs are very selective ...
either on rows, columns, or both

Job Selectivity

High selectivity →
less bytes to transfer

- ✓ Good for operators
- ✓ Cheaper for users for cross-data centers scenarios [see paper]

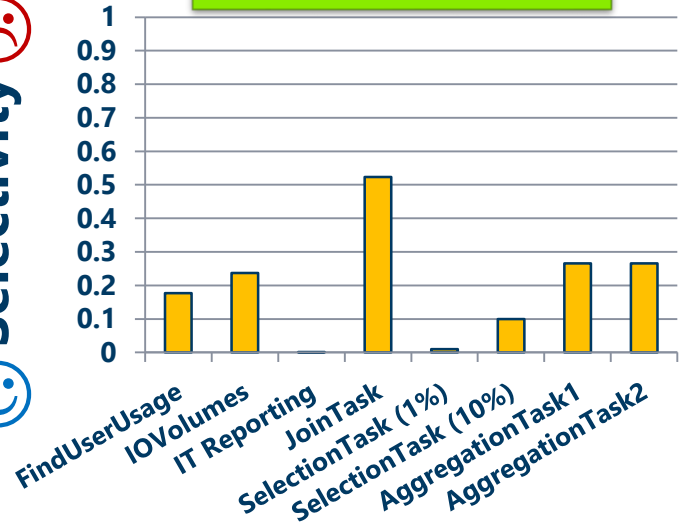
reduce runtime

- ✓ Good for users

Selectivity



Rows + Columns



Many jobs are very selective ...
either on rows, columns, or both

Measuring runtime benefits

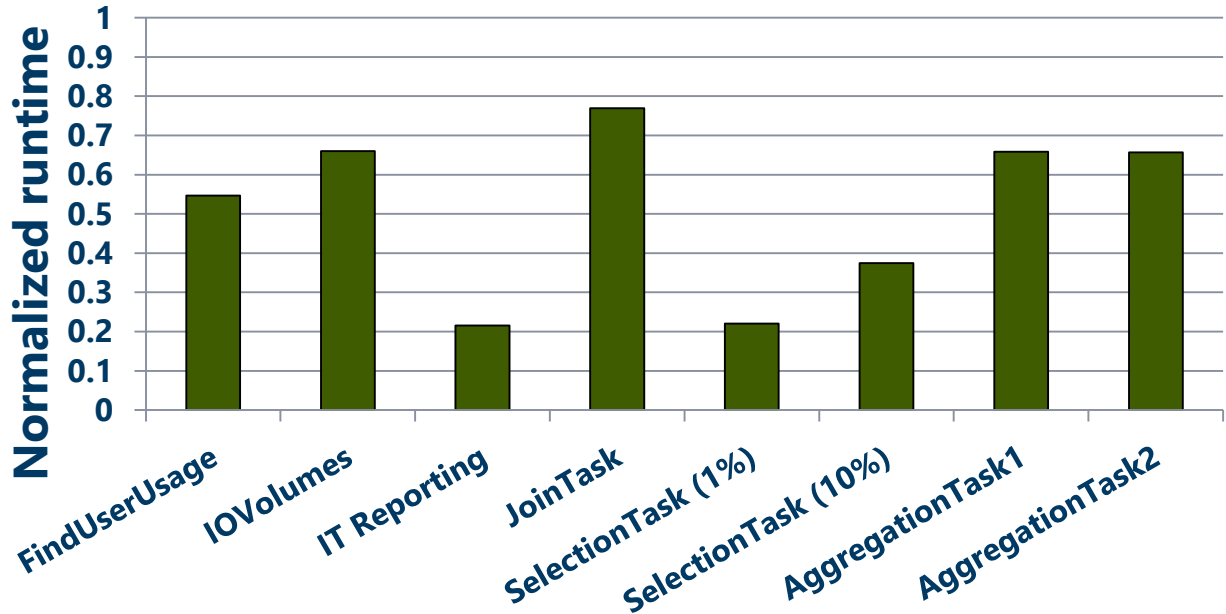
- We cannot extend Azure Storage or Amazon S3 with filters ☹️
- Instead, **we use pre-filtered data** and compare with unfiltered data
- We assume storage with: (a) scalable I/O, and (b) enough processing power for filtering

Diversion:

Do we have enough processing power?

- Row & Column filtering in **Java**: ~100MBytes/sec per core
- Scales linearly with multiple cores
- ≤ 2 cores for filtering enough for all but 1 job
- Runtime **always reduces** runtime, even with fewer cores
- Performance dominated by string input/output, not filter
- Column filtering in optimized **C**: **5-17x faster** than Java

Runtime benefits



- 30-80% reduction in runtime
- Runtime reductions less than selectivity due to Hadoop overheads

Conclusions

- Hadoop in the cloud:
separation of storage and compute.
- Rhea minimizes transfers from
storage to compute
 - Uses static analysis on the job bytecode
 - Extracts **selection** and **projection** operators from code
 - Generates filters to run in the storage layer
 - Runs transparently to user (and is safe for provider)
 - Potential benefits to the user (time, money) and
cloud provider (bandwidth)

