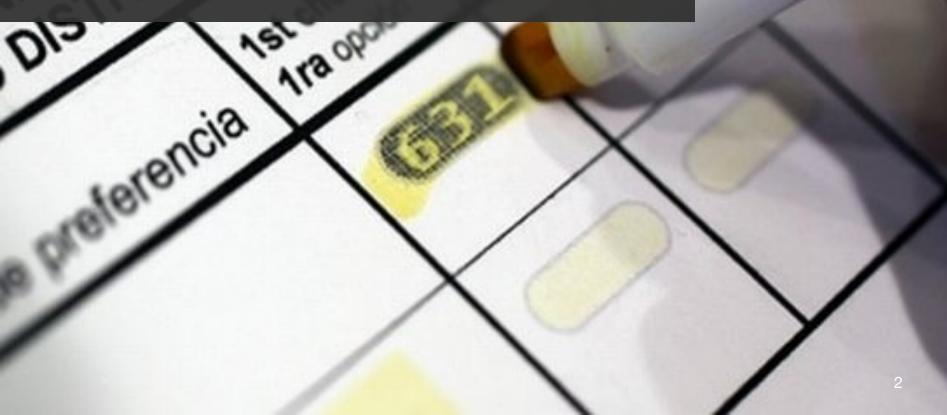
Cobra: Toward Concurrent Ballot Authorization for Internet Voting

Aleksander Essex, Jeremy Clark, & Urs Hengartner



End-to-End Verifiable VotingVerifiably correct tallyBallot secrecy



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Internet Voting Coercion & vote selling Untrustworthy platform Denial of service

JCJ, Civitas, Selections, Araujo et al., Spycher et al.

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Untrustworthy Platform

SureVote, Code Voting, Pretty Good Democracy, Remotegrity

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Denial of Service Attacks Application layer flooding

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Coercion-Resistance

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Contra la

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A voter can convince an adversary she voted for Alice while actually voting for Bob





- Voters choose a password that allows them to vote during registration
- The password scheme has a simple cognitive rule for creating fake ("panic") passwords
- Fake passwords can be sold or supplied under coercion
- The system will accept votes with fake passwords, but these votes will be obliviously canceled out
- Voters can vote with their real password any time. This ballot is unlikable to any ballots they cast with fake passwords

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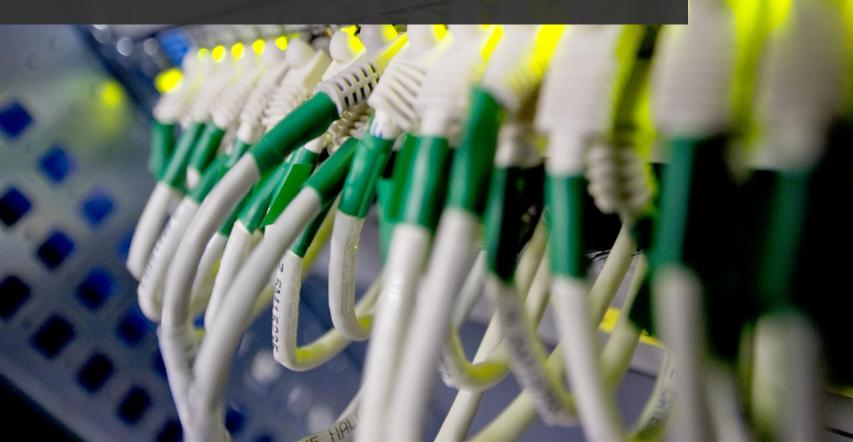


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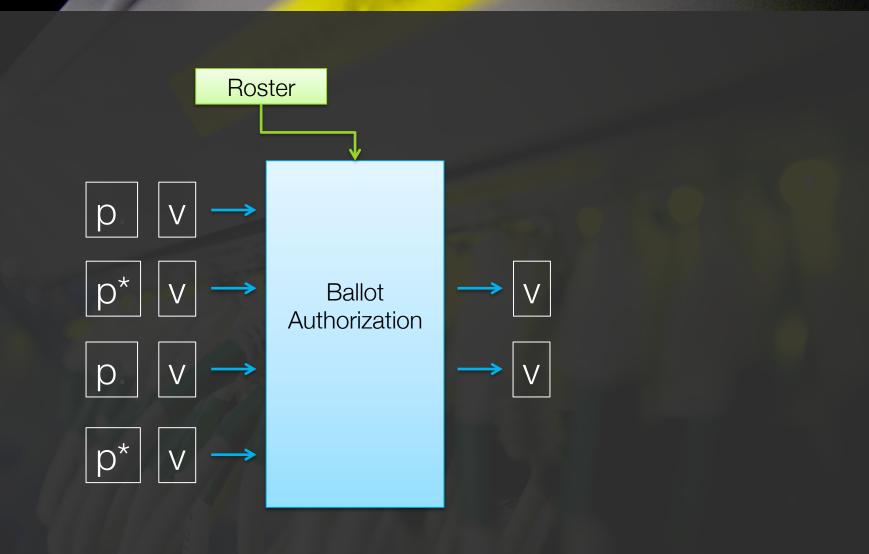
HB 24 B2

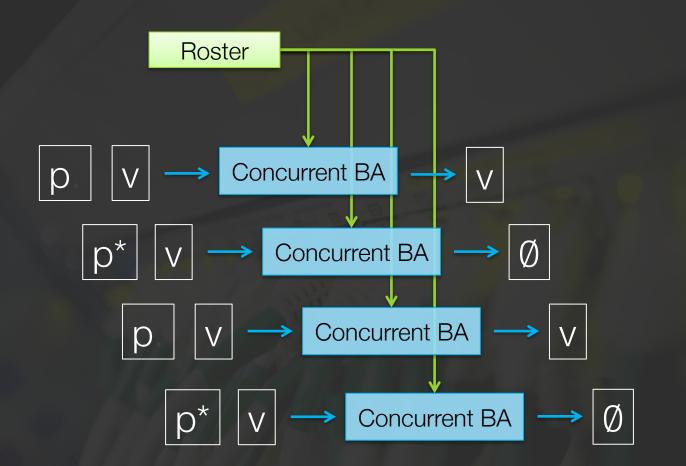
Denial of Service

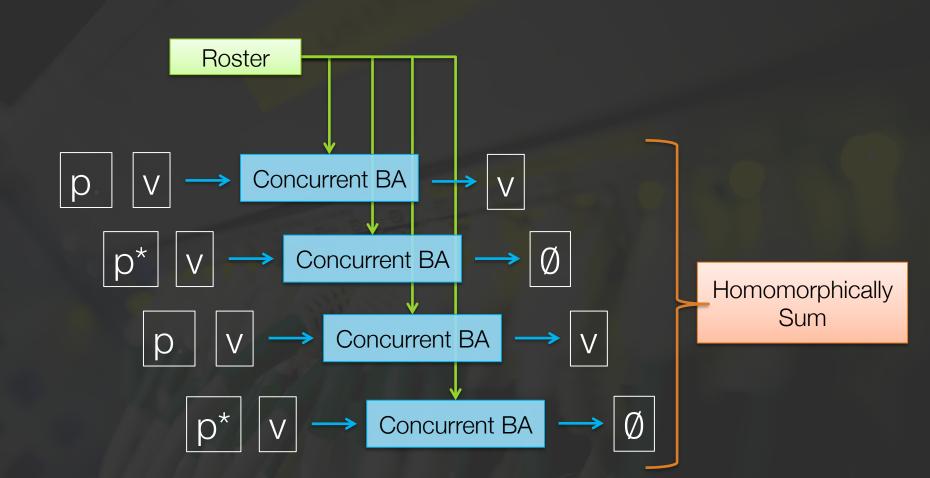
- Application-layer flooding
- Concurrent ballot authorization



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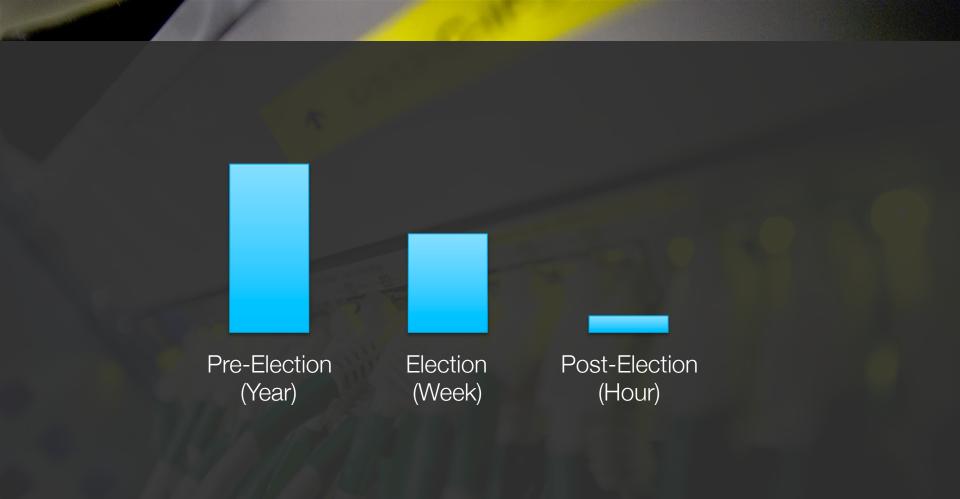




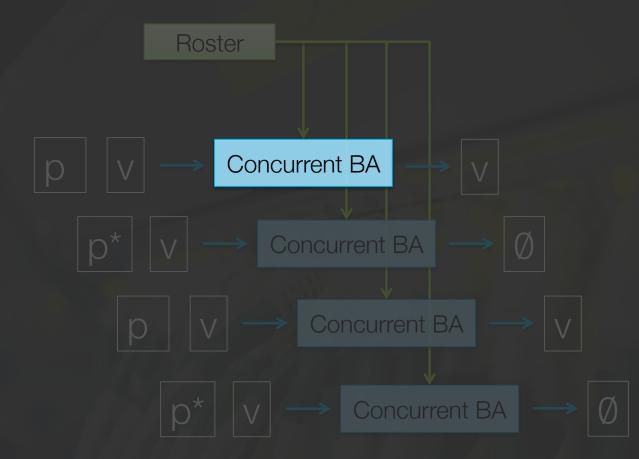




Work done by election authority

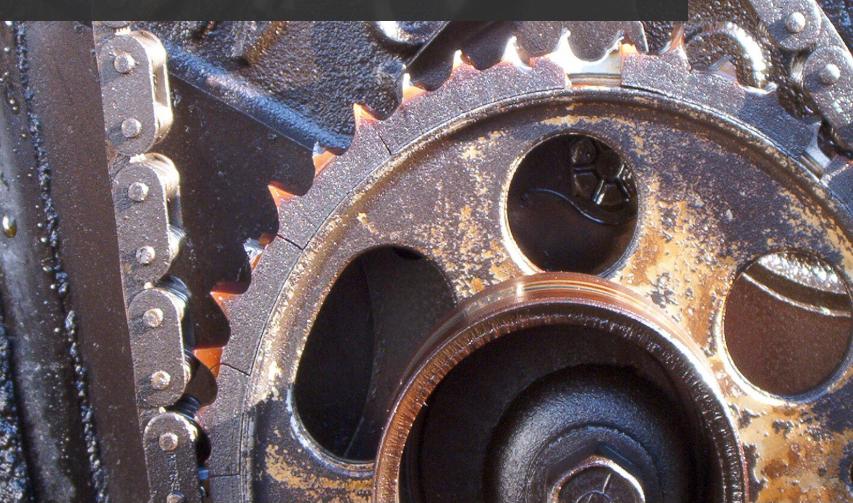


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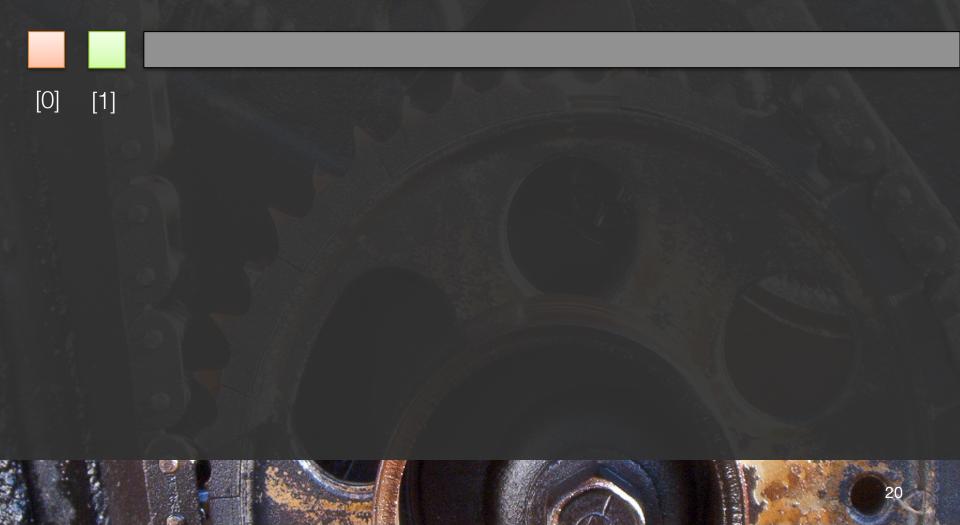


E HA

Fundamental MechanismPrivate Set Membership



Is encrypted password [p] on the roster?



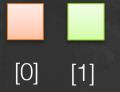
Is encrypted password [p] on the roster?

[0]

[1]

No: $[0] \times [v] = [0]$ Yes: $[1] \times [v] = [v]$

Is encrypted password [p] on the roster?



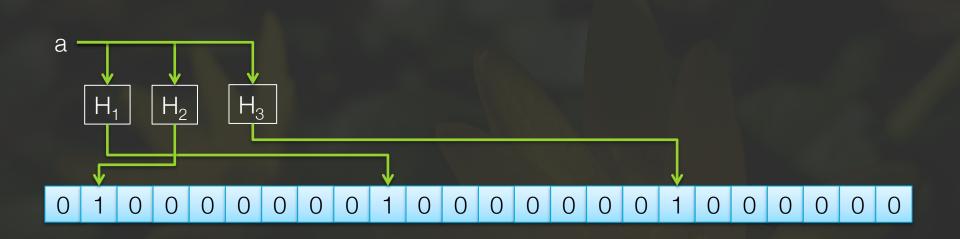
[0]

Plaintext equality tests & polynomials:

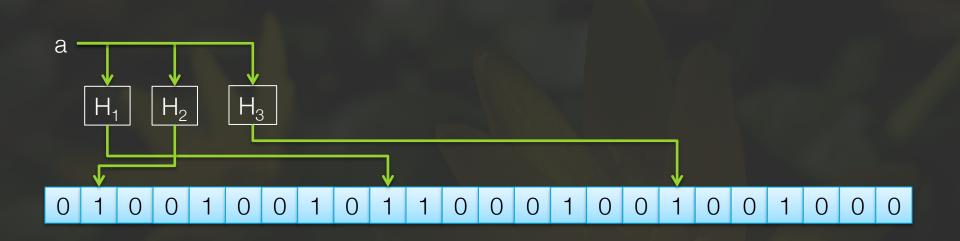
[r]

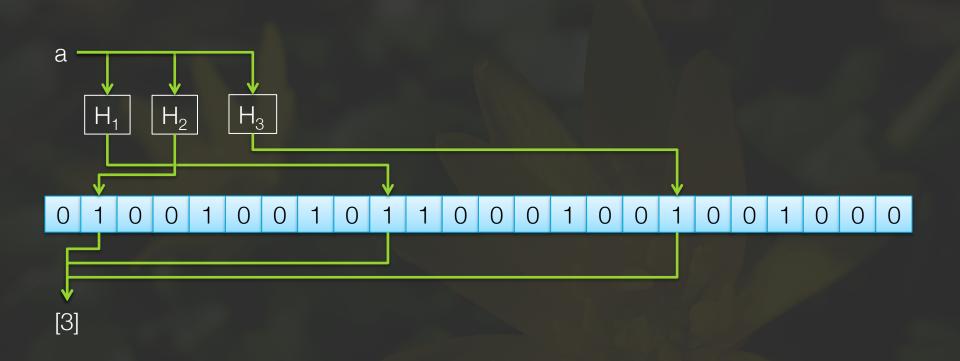
Encrypted Bloom Filters

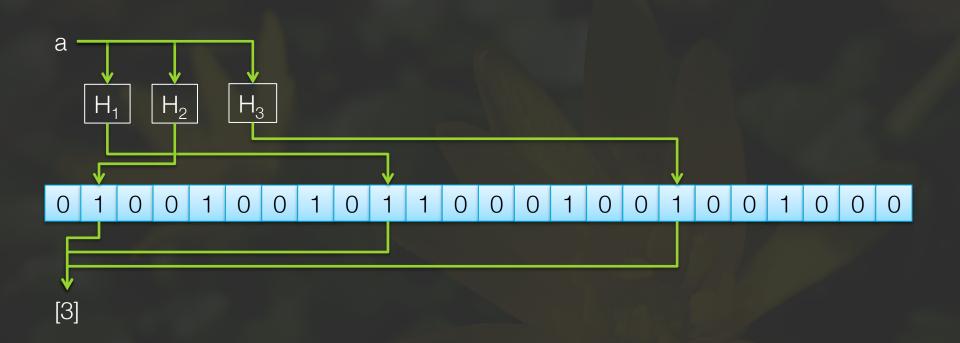
0 0 0 0 0 0



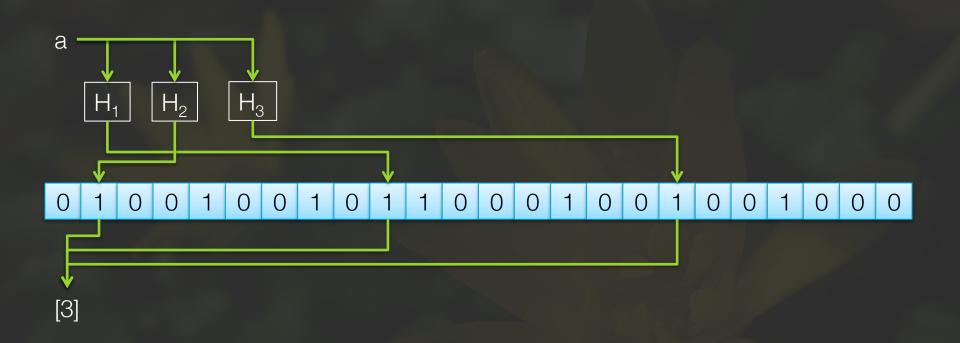
0 1 0 0



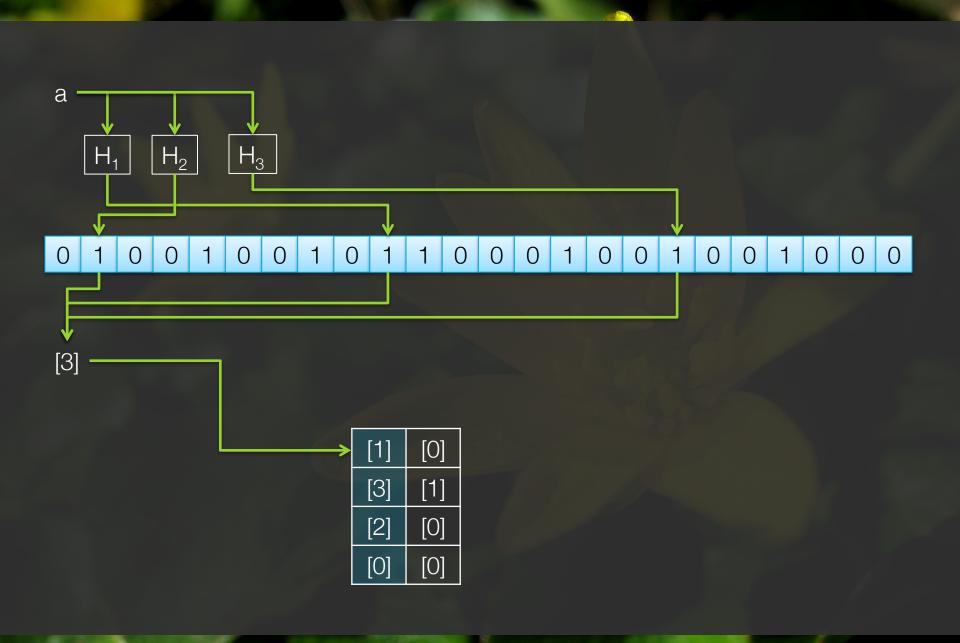


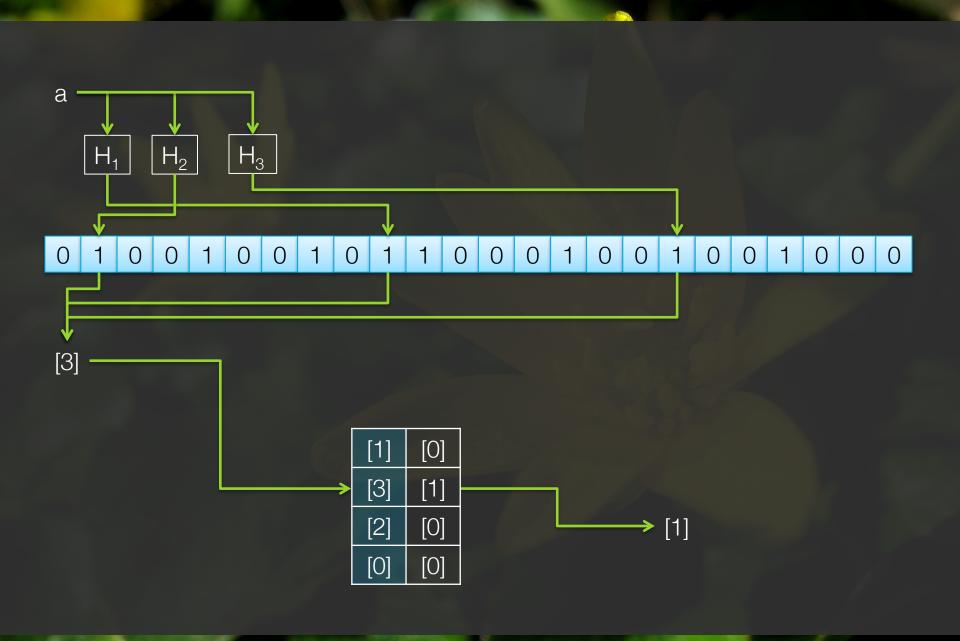


[0]	[0]
[1]	[0]
[2]	[0]
[3]	[1]



[0]	[0]	[1]	[0]
[1]	[O]	[3]	[1]
[2]	[O]	[2]	[0]
[3]	[1]	[0]	[0]







TO

- Voters' (obfuscated) passwords are added to an encrypted Bloom filter during registration
- See paper for details
- Properties:
 - Registrar does not see obfuscated password
 - Publicly verifiable proof that each voter added only a single entry
 - Coercion-resistant

$< g^{p}$, [v] , PoK(p: g^{p}) , PoM(v) >

$< g^{\rm p}$, [v] , PoK(p: g^{\rm p}) , PoM(v) > ^{-1}

Check Proofs

 $< \mathrm{g}^{\mathrm{p}}$, $[\mathrm{v}] > \checkmark$

< gp , [v] , PoK(p: gp) , PoM(v) >

$< g^{p}$, [v] > \leftarrow Duplicate Check

< g^p , [v] , PoK(p: g^p) , PoM(v) >

$< g^{p}$, [v] > -

Query(g^p)

< [1] , [V] > <

< g^p , [v] , PoK(p: g^p) , PoM(v) >

$| < g^p$, [v] >

< [1] , [V] > -

< [V] > <

Mix & Match:

[0]	[0]
[1]	[V]

$< g^{p^{\star}}$, [v] , PoK(p: g^p) , PoM(v) >

$|<\mathrm{g}^{\mathrm{p}^{\star}}$, [V] >

< [O] , [V] > -

لـ ◄ < [0] > ◄

Mix & Match:

[0]	[0]
[1]	[V]

See paper for more:

- Registration: setting up the Bloom filter (*expensive!*); setting false positive rate
- Optimizations: using BGN to eliminate steps
- Security analysis: eligibility verification, integrity, coercion-resistance
- A blueprint that might be useful for concurrent ballot authorization other ways

Performance



Registration (Before Election)

Voter	11	9	39	11	55,680,006
Registrar	8	10	20	8	37, 120, 006

Casting (During Election)

Submit Ballot	42	42	42	42	42
Submit Credential	3	13	202	3	2

Processing (During Election)

Check Ballots	240,000	240,000	240,000	240,000	240,000
Ballot Authorization	0	0	2,000,000	0	10,790,000

Processing & Tallying (After Election)

Ballot Authorization	3,000,960,000	4,080,000	2,010,000	100,710,000	0
Tally Ballots	45	45	45	45	45

Table 1: Performance comparison in number of modular exponentiations for a moderately-sized election scenario: 5 candidates, 10,000 registered voters, 20,000 submitted ballots, and 3 trustees.

Concluding Remarks

- DOS on internet voting is a reality
- Common properties of coercion-resistance systems (anonymous ballot submission, intensive post-tally processing) make protocol-level DOS a threat
- We have shown in principle ballots can be authorized concurrently (and incidentally post the fastest tally with Cobra)
- Future work: speed-up registration

Questions?

MITA

@AleksEssex
 @PulpSpy
 @uhengart

TICKET VENDING MACHINES

LINE DEPARTURES	NEW HAVE
HOUNT VERNON - IST STOP	TIME TRK DESTINATION
RYE + IST STOP Hount verhon - Ist stop	4:15 1028 NEW HINES 4:25 110 HINESSON
GREENWICH - IST STOP	4:34 77 NEW NEWEN 114 4:40 108 NEW NEWEN
STAMFORD - IST STOP	4:43 104 NEW CONTRA

INCLUSION Kiewan neo ny Lint onaechan ne fa shi uo ne anna fi ar a sa na bi te o n u te o n u te o n

INE DEBARTURE