

#### Zanzibar: Google's Consistent, Global Authorization System

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## Authorization checks are central to preserving privacy



### Zanzibar stores access control lists (ACLs)



...and performs authorization checks based on stored ACLs



## Zanzibar is...

- **Consistent**: Respects causal ordering of updates to ACLs and objects
- **Flexible**: Supports a rich variety of access control policies
- **Scalable**: Trillions of ACL entries, millions of checks/second
- Fast: Less than 10ms @ 95%, less than 100ms @ 99.9%
- **Available**: 99.999% over the past 3 years

## Namespaces, relations, usersets, and tuples

Namespace: videos

Object	Relation	Userset
video X	viewer	user A
video Y	viewer	All Users

System-defined value

### Namespaces, relations, usersets, and tuples

#### Namespace: videos

Object	Relation	Userset
video X	viewer	user A
video Y	viewer	All Users

#### **Check results:**

- video X, viewer, user A? Yes
- video X, viewer, user B? No
- video Y, viewer, user A? Yes
- video Y, viewer, user B? Yes

## Userset indirection can create deep/wide hierarchies

Namespace: videos

Namespace: groups

Object	Relation	Userset		Object	Relation	Userset
video X	viewer	user A		group 1	member	user B
video X	viewer	(group 1, member)		group 1	member	user C

# Userset indirection can create deep/wide hierarchies

Namespace: videos

Namespace: groups

Object	Relation	Userset		Object	Relation	Userset
video X	viewer	user A		group 1	member	user B
video X	viewer	(group 1, member)		group 1	member	user C

#### **Check results:**

- video X, viewer, user B? Yes
- video X, viewer, user D? No





# **Consistency protocol**

#### ACL update by Alice:



UpdateACL(doc X, viewer, remove Bob)



Leverages Spanner's TrueTime mechanism [Corbett et al. 2012]

time

# **Consistency protocol**

ACL update by Alice:

UpdateACL(doc X, viewer, remove Bob)



#### Content update by Charlie:

CheckContentUpdate(doc X, writer, Charlie)



timestamp T0

time

# **Consistency protocol**

ACL update by Alice:







ACL check for Bob:		
CheckACL(doc X, viewer, Bob, T1)		
	<b>No</b> [at T2 ≥ T1 > T0]	>



# Implementation techniques

- Timestamps chosen to reduce latency
- Hot-spot mitigation to increase availability
- Request hedging to reduce tail latency
- Isolation to protect against misbehaving clients
- Optimized processing of large and deeply nested sets

# Deployment

- Zanzibar has been in production use for > 5 years
- > 1,500 namespaces defined by hundreds of clients
- > 2 trillion relation tuples replicated in several dozen locations worldwide
- > 10 million client queries per second, mostly read-only
- > 10,000 servers in several dozen clusters worldwide

Check queries per second



Checks peak at 4.2M QPS, Reads at 8.2M, Expands at 760K, Writes at 25K



95th-percentile latency is below 10 ms, 99.9th-percentile below 100 ms

# Availability



Availability over the last 3 years has remained above 99.999%

# Summary

- Robust authorization checks are central to preserving privacy
- Zanzibar is a unified authorization system for Google services
  - Respects causal ordering of user actions
  - Supports a rich variety of access control policies
  - Offers low latency and high availability
  - Scales to trillions of ACL entries and millions of checks per second
  - Supports hundreds of services used by billions of people

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