

Truly Non-blocking Writes

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HotStorage Workshop, 2011

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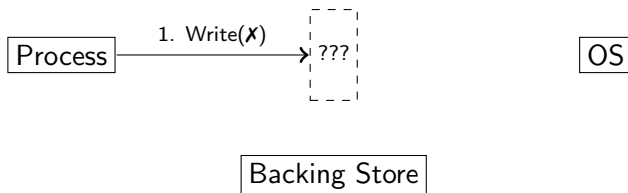
Process

OS

Backing Store

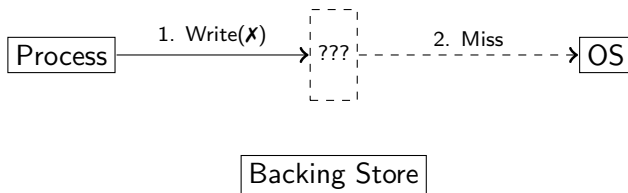
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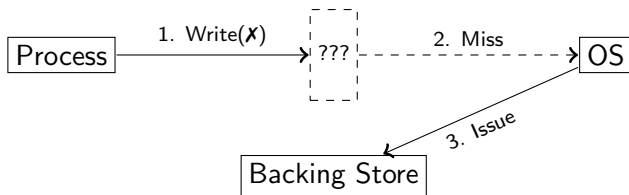
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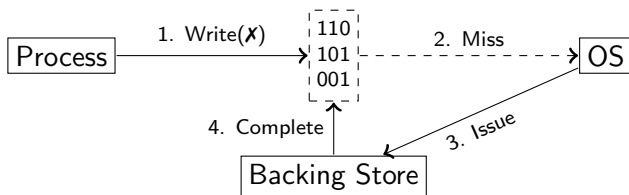
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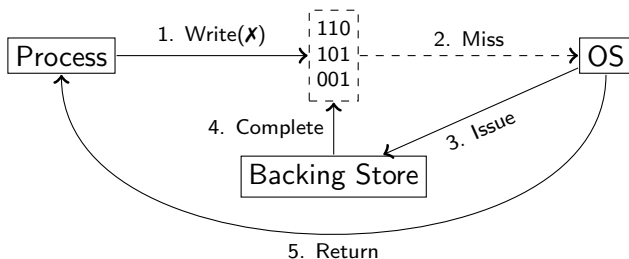
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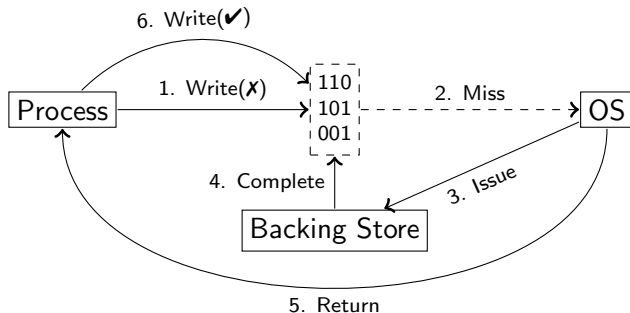
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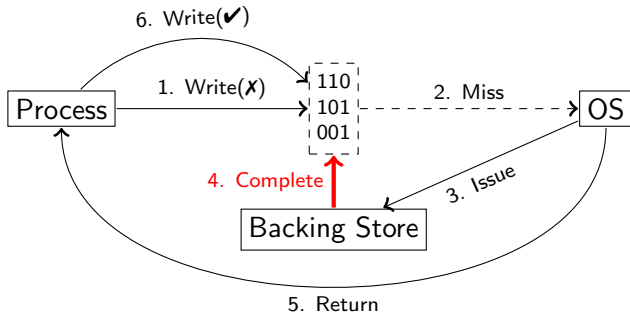
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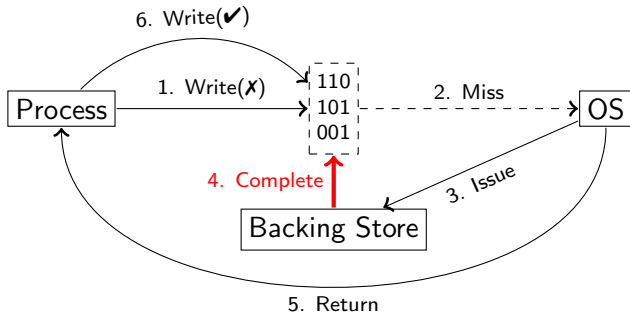
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For writes: why wait for data that the application doesn't need?

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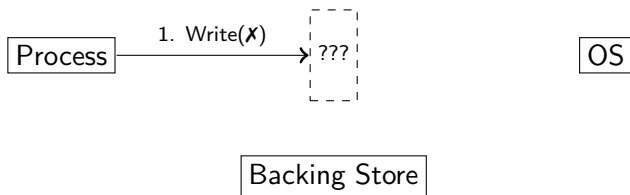
Non-blocking Writes: Basic Approach

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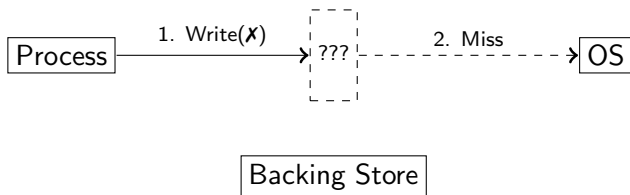
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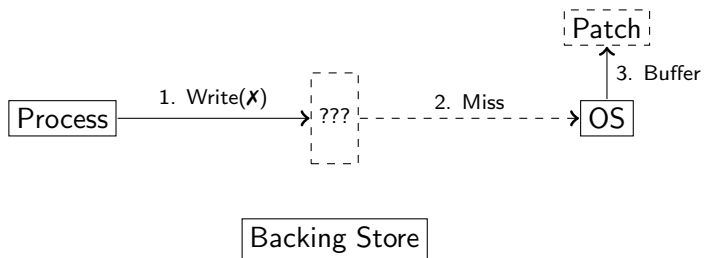
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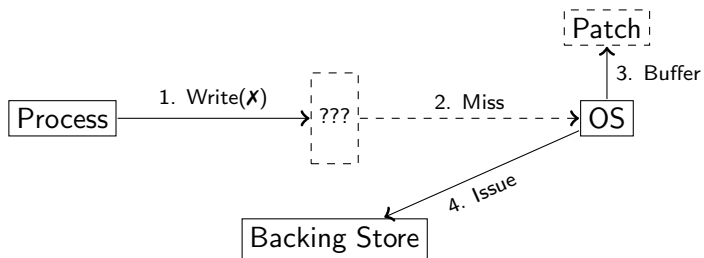
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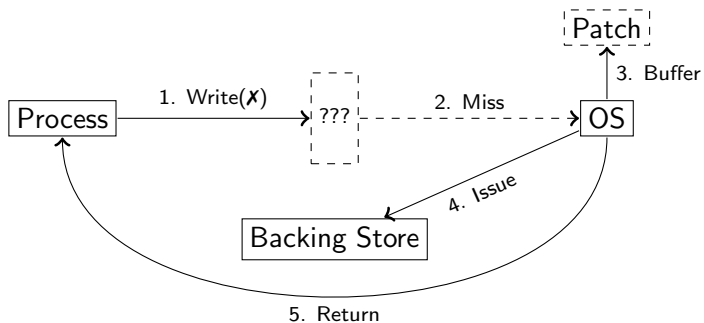
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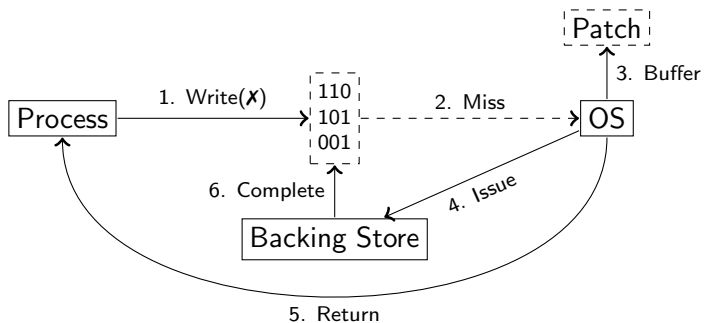
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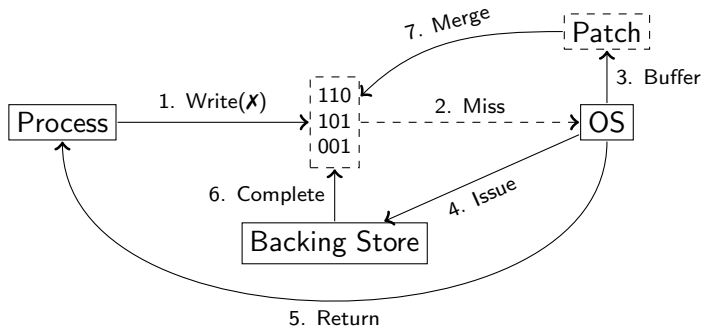
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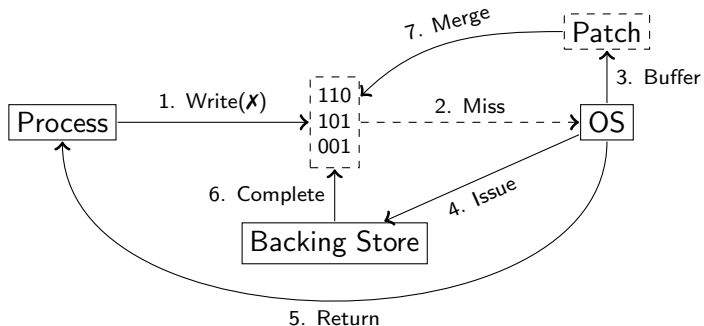
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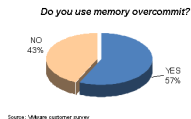
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Benefits

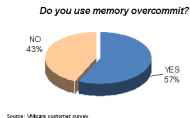
1. Application execution time reduction
2. Increased backing store bandwidth usage

Motivation → Higher Fault Rates



Memory over-committed in virtualized environments

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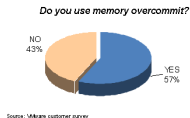


Memory over-committed in virtualized environments



More process running with multi-core and virtualized environments

Motivation → Higher Fault Rates



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Memory hierarchy moving towards a more active and faster backing store

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- ▶ We calculate the % of faults that can benefit in all our workloads

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- Image Processing Rendering of SVG images

- Developer Unit and performance testing

- Server Application, database, and mail server

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- ▶ We calculate the % of faults that can benefit in all our workloads:
 - Image Processing Rendering of SVG images
 - Developer Unit and performance testing
 - Server Application, database, and mail server
- ▶ Simulator with full-system memory traces.
- ▶ RAM set to 50% of app footprint

Motivation → % Non-blocking faults

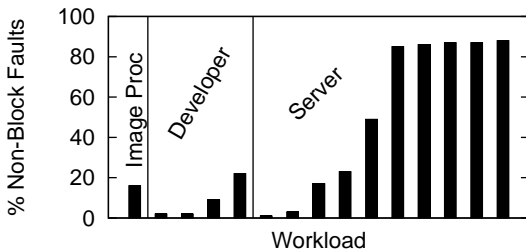
- ▶ We calculate the % of faults that can benefit in all our workloads:

Image Processing Rendering of SVG images

Developer Unit and performance testing

Server Application, database, and mail server

- ▶ Simulator with full-system memory traces.
- ▶ RAM set to 50% of app footprint
- ▶ Up to 80% of page faults benefit



Alternatives to non-blocking writes:

Perfect DRAM Provision

Unpredictable or unbounded.

Prefetching

Can incur false positives and false negatives.

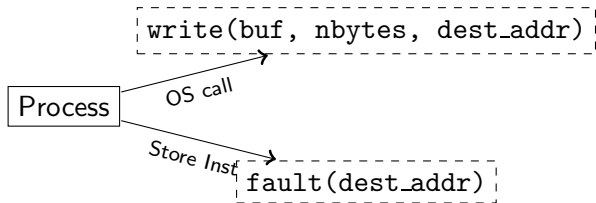
Asynchronous System Calls

1. Do not work with memory mapped pages
2. Written data *not* immediately available for reading

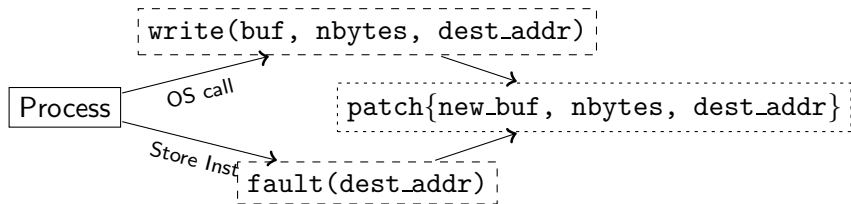
Solution Challenges

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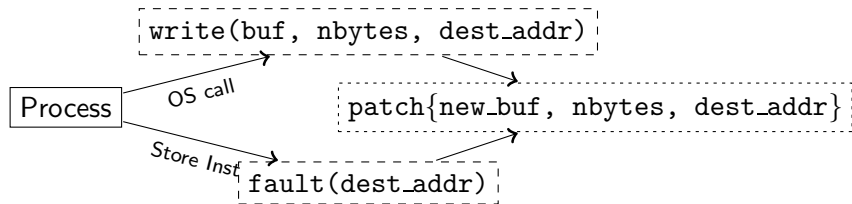
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Information Per Non-blocking Write

Information

Write Offset

Data Written

Size of Data

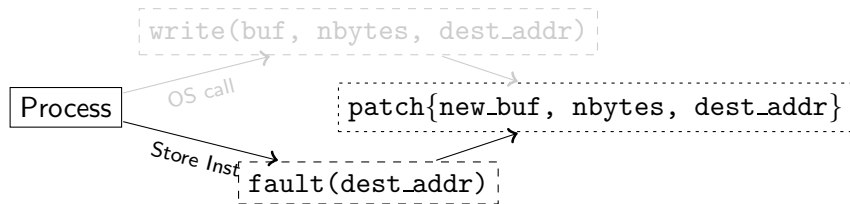
Solution Challenges



Information Per Non-blocking Write

Information	Supervised
	<code>write()</code>
Write Offset	✓
Data Written	✓
Size of Data	✓

Solution Challenges



Information Per Non-blocking Write

Information	Supervised <code>write()</code>	Unsupervised Fault
Write Offset	✓	✓
Data Written	✓	✗
Size of Data	✓	✗

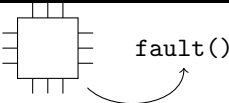
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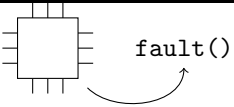
Information Per Non-blocking Write

Information	Supervised <code>write()</code>	Unsupervised Fault
Write Offset	✓	✓
Data Written	✓	✗
Size of Data	✓	✗

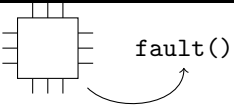
Handling Unsupervised Writes

Approach	Description	Fast	All Arch?	Low Mem?
Full Feature Hardware	 A square component with pins on all four sides. An arrow points from the component to the text "fault()".	✓	✗	✓

Handling Unsupervised Writes

Approach	Description	Fast	All Arch?	Low Mem?
Full Feature Hardware		✓	X	✓
Opcode Disassembly	4 bytes offset <code>sw \$t1, 0xff</code> data	✓	X	✓

Handling Unsupervised Writes

Approach	Description	Fast	All Arch?	Low Mem?
Full Feature Hardware		✓	✗	✓
Opcode Disassembly	4 bytes offset <code>sw \$t1, 0xff</code> data	✓	✗	✓
Page Diff-Merge	Disk Page or 0-buffer and 1-buffer <hr/> Updated Page	✗	✓	✗

Quantifying Benefits

1. Fraction of non-blocking write faults ✓
2. Outstanding write faults (over time)
3. Savings in execution time (new!)

Virtual Memory Simulator

Input RAM size & Full System Memory Traces

Output Performance statistics

- ▶ Memory size set to 50% of workloads footprint
- ▶ Creating patches is not required

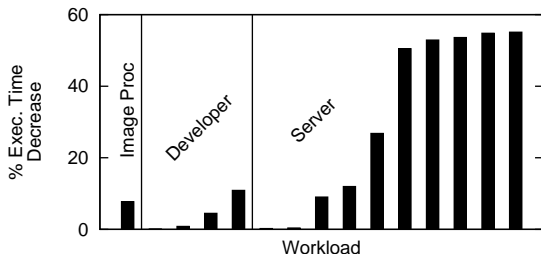
Quantifying Benefits → Metric

- ▶ How to measure the additional parallelism?
- ▶ **Outstanding Write Faults (OWF):** # of parallel write faults at any time
 - ✓ $OWF \leq OIO$
 - ✓ $OWF \leq 1$ for single threaded applications
 - ✓ $OWF \geq 0$ when using non-blocking writes
- ▶ We need the variations over time as well
- ▶ $E[OWF]$: time-weighted average OWF



Quantifying Benefits → Time Reduction

- ▶ These results are not in the paper
- ▶ Execution time = Trace time + Synchronous read time
- ▶ Write time of dirty page on evictions ignored
- ▶ Rough estimate: error proportional to the number of dirty pages evicted



Conclusions and Future Work

- ▶ We presented non-blocking writes: a technique to eliminate read-before-writes
 - ✓ Reduced execution time
 - ✓ Increased device usage
- ▶ We estimate a reduction times of 0.1-54%
- ▶ In the future, we are planning to implement non-blocking writes to better study its implications
 - ✓ What workloads benefit from Non-blocking writes?

Questions?

Virtual Memory Simulator

Input: RAM size & Mem Traces

Output: Per Entry: Timestamp and event (hit, miss, evict);
Global: Performance stats.

- ▶ Writes to out-of-core pages considered non-blocking
- ▶ Non-blocking status revoked when:
 1. The page is read before I/O completion
 2. The page is evicted before I/O completion

Quantifying Benefits → Full System Memory Traces

Modified x86 software-MMU QEMU to log all memory accesses:

- ▶ *Instruction count, CR3, virtual/physical address, access-mode, page privileges.*

Workloads

Type	#	Footprint Avg/Std (MB)
Server	10	294/158
Developer	4	269/183
Image	1	149/0

Solution Approaches → Page Diff-Merge

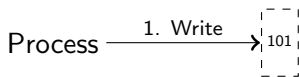
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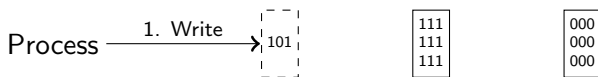
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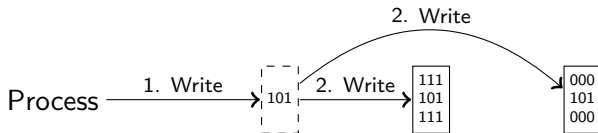
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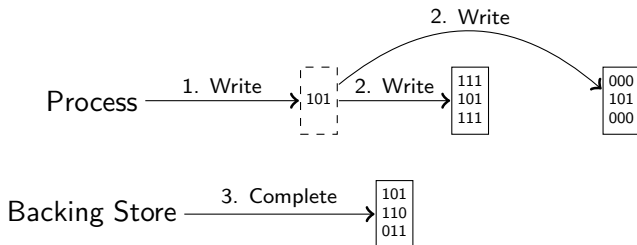
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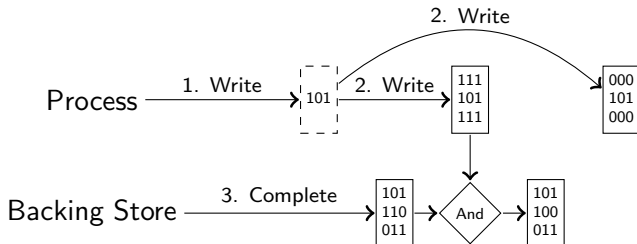
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