

Estimating the Margin of Victory for Instant-Runoff Voting*

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* also known as Ranked-Choice Voting,
preferential voting, and the alternative vote

Overview

- Why estimate?
- What are we talking about?
- Estimates
- Worst-case accuracy
- Real elections
- Conclusions

Why Estimate?

Trustworthy Elections

 Risk-limiting audits

 Margin of Victory

Why Estimate?

IRV Trustworthy Elections



IRV Risk-Limiting Audits



IRV Margin of Victory
(not feasible)

Why Estimate?

IRV Trustworthy Elections ?

 IRV Risk-~~X~~limiting Audits

 IRV ~~X~~Margin of Victory
(not feasible)

Why Estimate?

IRV Trustworthy Elections

IRV Risk-Limiting Audits

IRV Margin of Victory
(not feasible sometimes)

IRV Margin of Victory
Lower Bound

Proposals for IRV Risk-Limiting Audits

Risk-limiting audits for nonplurality elections

Sarwate, A., Checkoway, S., and Shacham, H.
Tech. Rep. CS2011-0967, UC San Diego, June 2011
<https://cseweb.ucsd.edu/~hovav/dist/irv.pdf>

Overview

- Why estimate? **because we can;**
to do risk-limiting audits
- **What are we talking about?**
 - **What is Instant-Runoff Voting?**
 - **What is a margin of victory?**
- Estimates
- Worst-case accuracy
- Real elections
- Conclusions

Model of Instant-Runoff Voting

- Single winner
- Ballot ranks candidates in order of preference.
- Votes are counted and candidates are eliminated in a sequence of rounds.
- In each round, a ballot counts as one vote for the most preferred continuing candidate on the ballot, if one exists.
- In each round, one candidate with the fewest votes is eliminated for subsequent rounds.
- Ties for elimination are resolved by lottery.
- Rounds continue until just one candidate is in the round. That candidate is the winner.

Consistent IRV Features

- Number of candidates ranked on a ballot:
 - require ranking all candidates
 - limit maximum number of ranked candidates
 - can rank any number of candidates
- Multiple eliminations:
 - required*, not allowed, or discretionary*
- Early termination:
 - tabulation stops when a winner is identified*

* may require an extended tabulation for auditing purposes

Defining the Margin of Victory

The *margin of victory* is
the minimum total number* of ballots
that must in some combination
be added and removed
in order for the set of contest winner(s)
to change with some positive probability.

* the number of added ballots, plus
the number of removed ballots

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Estimates for the Margin of Victory

	<u>Time</u>	<u>Space</u>
Last-Two-Candidates upper bound	$O(1)$	$O(1)$
Winner-Survival upper bound	$O(C)$	$O(1)$
Single-Elimination-Path lower bound	$O(C^2)$	$O(1)$
Best-Path lower bound	$O(C^2 \log C)$	$O(C)$

(C = number of candidates)

Example IRV Contest

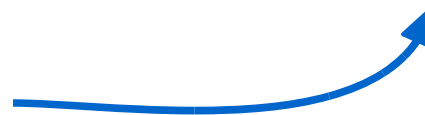
	Round 1	Round 2	Round 3	Round 4	Round 5
Wynda Winslow	107	112	114	186	332
Diana Diaz	130	133	134	146	—
Charlene Colbert	35	46	84	—	—
Barney Biddle	40	41	—	—	—
Adrian Adams	20	—	—	—	—

Candidates are in reverse order of elimination, with the winner first.

Last-Two-Candidates Upper Bound

	Round 1	Round 2	Round 3	Round 4	Round 5
Wynda Winslow	107	112	114	186	332
Diana Diaz	130	133	134	146	—
Charlene Colbert	35	46	84	—	—
Barney Biddle	40	41	—	—	—
Adrian Adams	20	—	—	—	—
				<u>40</u>	

Margin of Survival for Winner in round C – 1, the round with just the last two candidates.



Winner-Survival Upper Bound

	Round 1	Round 2	Round 3	Round 4	Round 5
Wynda Winslow	107	112	114	186	332
Diana Diaz	130	133	134	146	—
Charlene Colbert	35	46	84	—	—
Barney Biddle	40	41	—	—	—
Adrian Adams	20	—	—	—	—
Margin of Survival for Winner	87	71	<u>30</u>	40	

Smallest Margin of Survival for the Winner in the first $C - 1$ rounds.



Vote Totals Not In Sequence By Value

	Round 1	Round 2	Round 3	Round 4	Round 5
Wynda Winslow	107	112	114	186	332
Diana Diaz	130	133	134	146	—
Charlene Colbert	35	46	84	—	—
Barney Biddle	40	41	—	—	—
Adrian Adams	20	—	—	—	—

Vote Totals

In Sequence By Value

Round 1	Round 2	Round 3	Round 4	Round 5
130	133	134	186	332
107	112	114	146	—
40	46	84	—	—
35	41	—	—	—
20	—	—	—	—

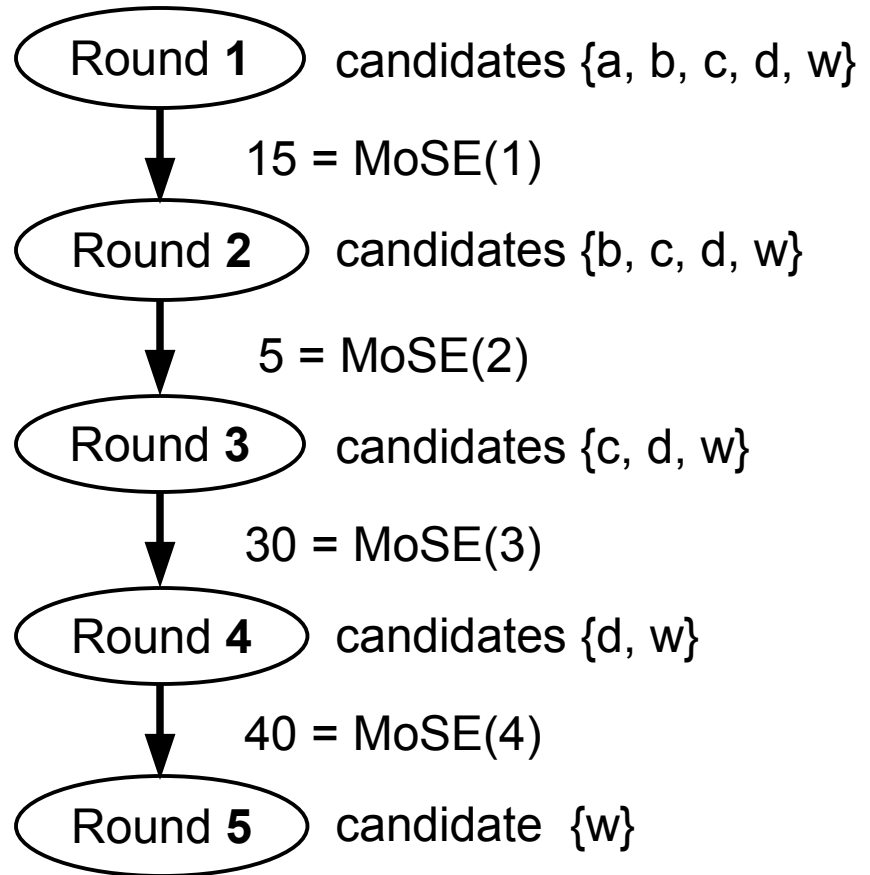
Single-Elimination-Path Lower Bound

	Round 1	Round 2	Round 3	Round 4	Round 5
	130	133	134	186	332
	107	112	114	146	—
	40	46	84	—	—
	35	41	—	—	—
	20	—	—	—	—
Margin of Single Elimination (MoSE)	15	<u>5</u>	30	40	

Smallest Margin of Single Elimination
in the first $C - 1$ rounds.

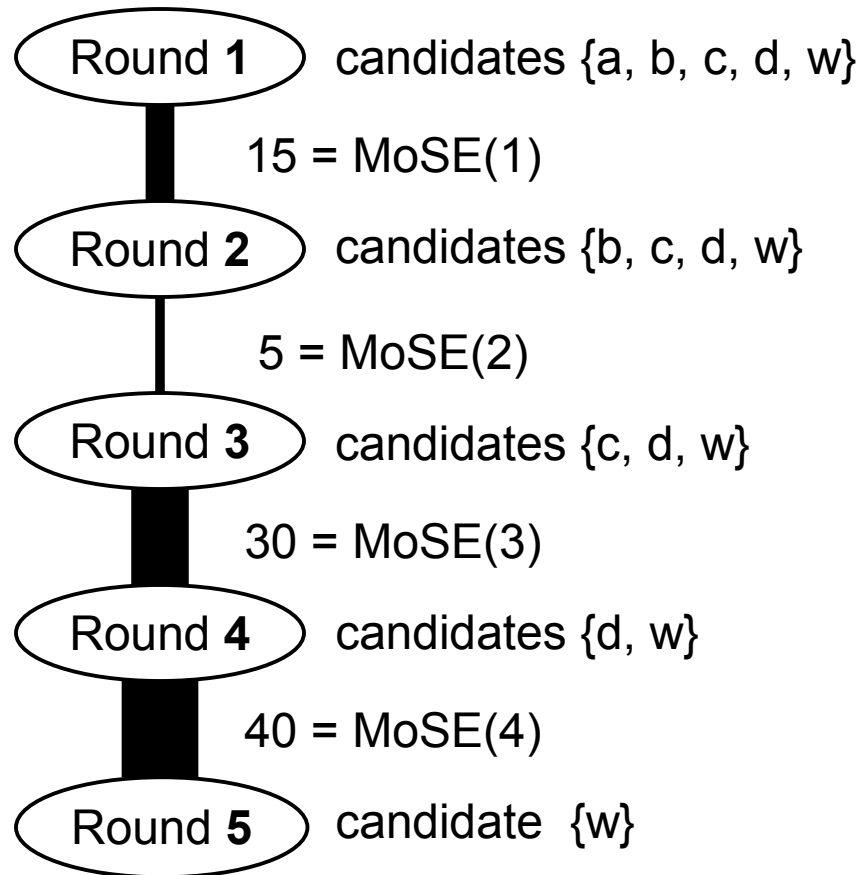


Single Elimination Path



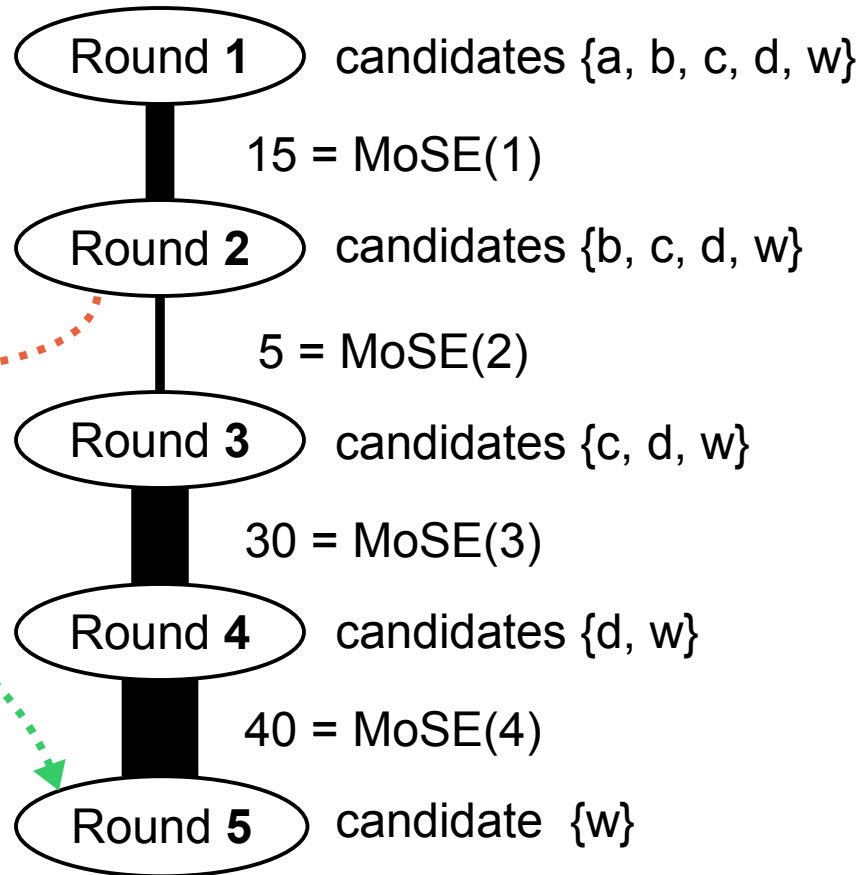
Single-Elimination Path Bottleneck

Edge weight =
a limited capacity
(a bottleneck)
for tolerating
additions and
removals of ballots,
while still staying
on the path.



Exceeding a Bottleneck

Easy guarantee
of same winner:
Stay on the
single-elimination
path



Different
Winner

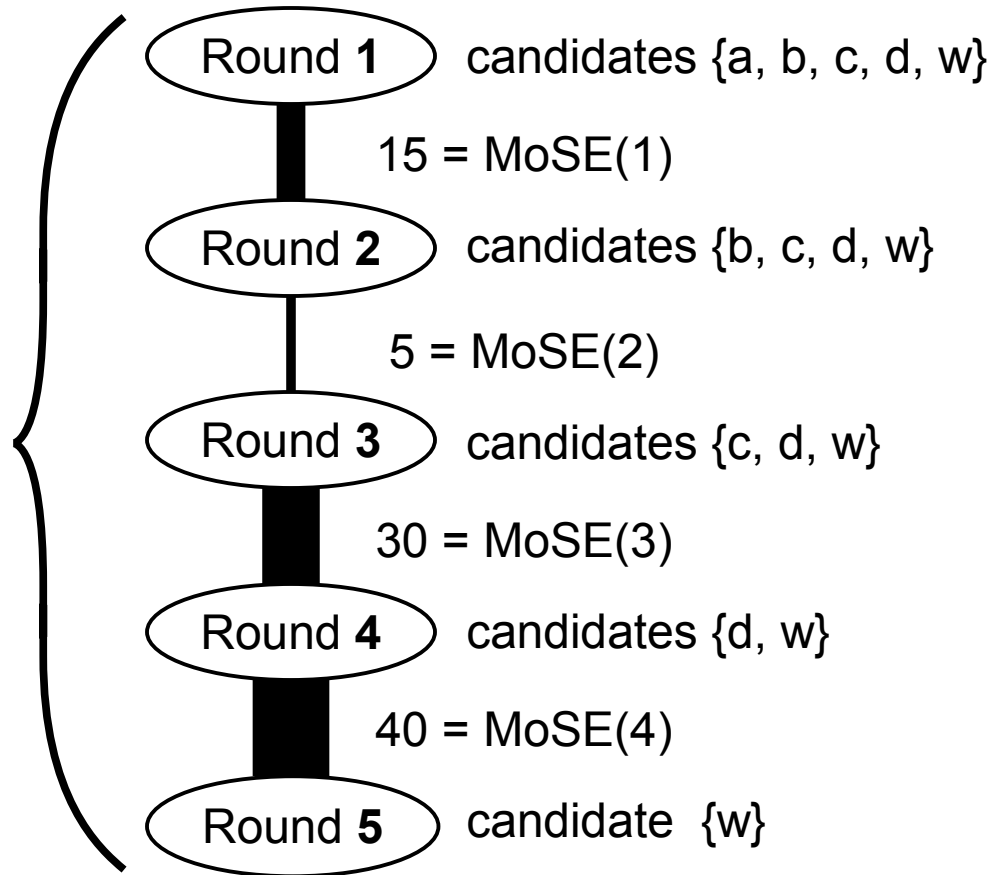
Different
Winner

Path Bottleneck

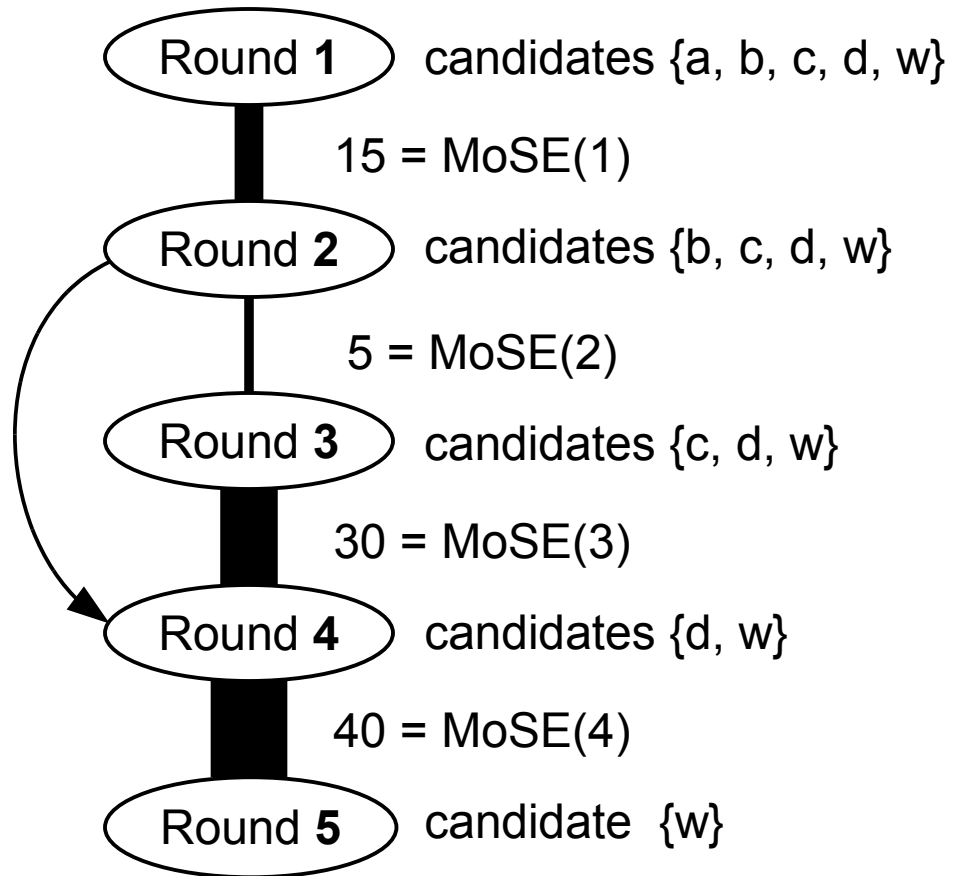
Path Bottleneck
is the smallest
individual
bottleneck
on the path

=

Single-
Elimination-
Path
lower bound



Multiple Elimination as a Detour



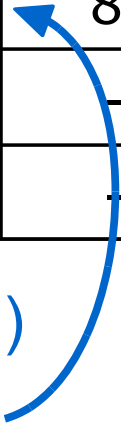
Multiple Elimination of k Candidates

Round 1	Round 2	Round 3	Round 4	Round 5
130	133	134	186	332
107	112	114	146	—
40	87	84	—	—
35		—	—	—
20	—	—	—	—

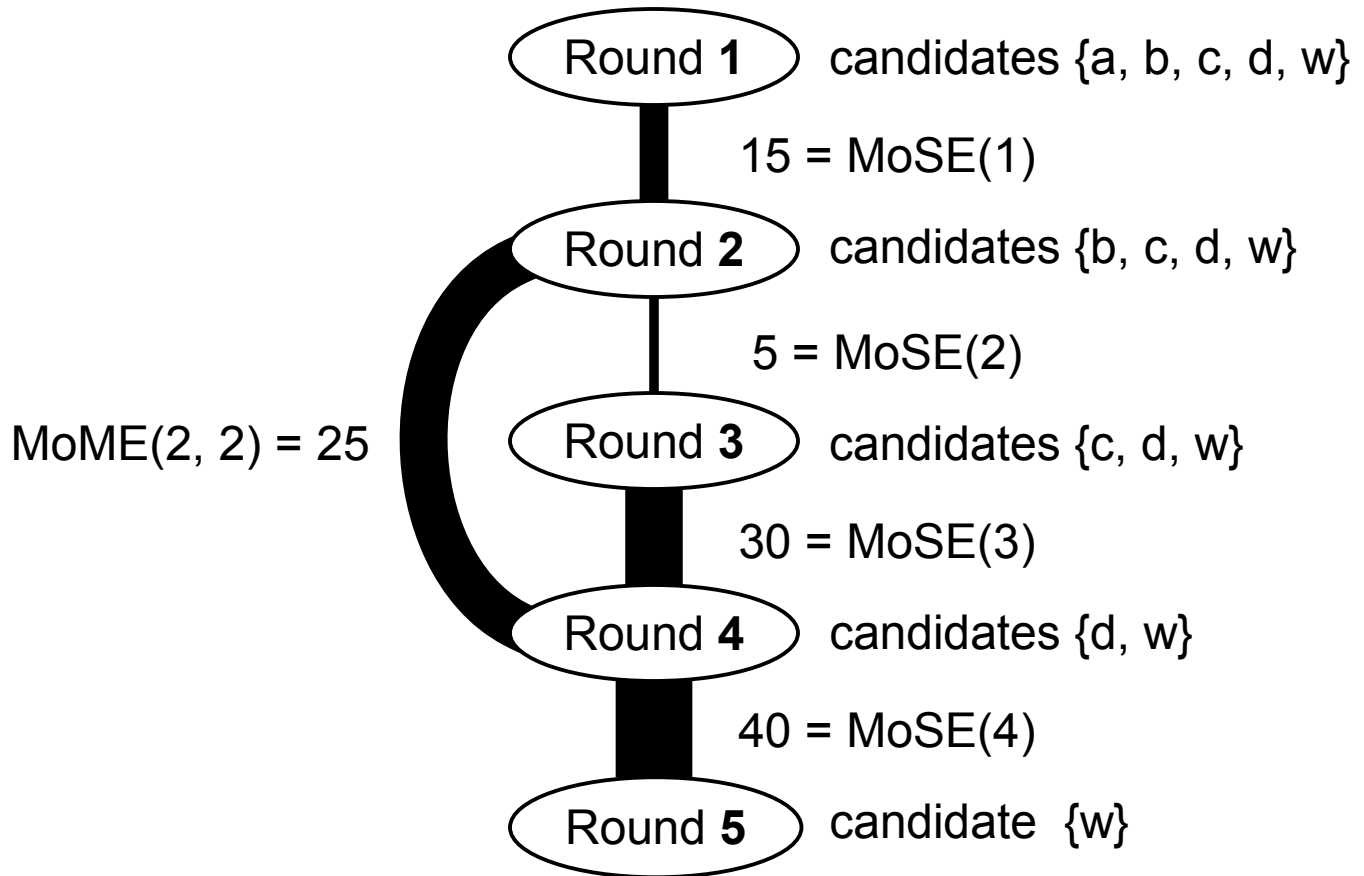
A usable
multiple eliminaton,
if
combined vote total
is still the smallest

Margin of Multiple Elimination

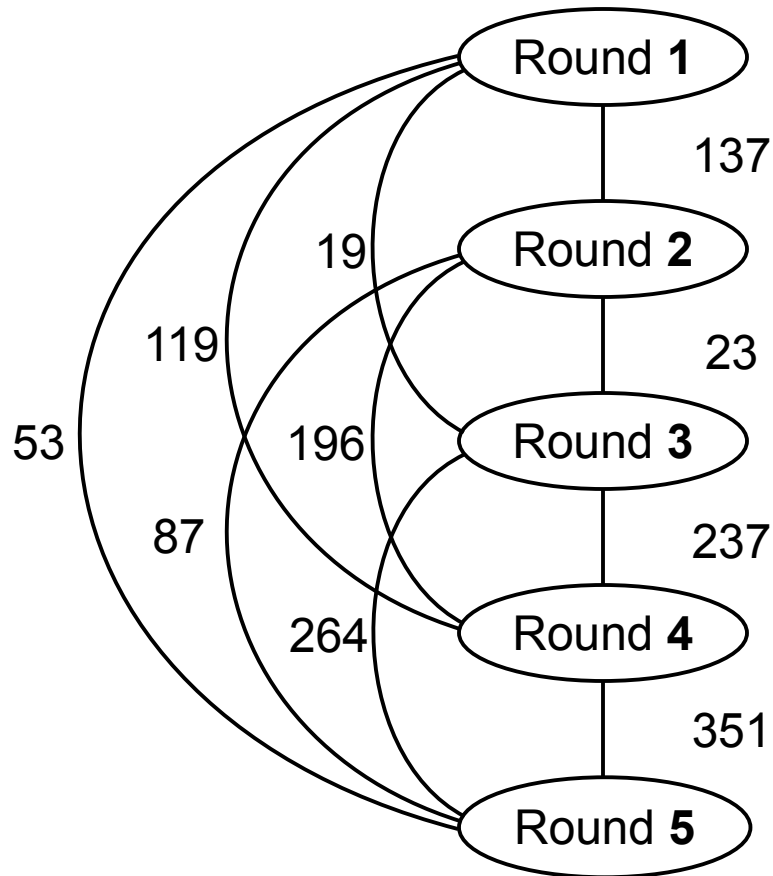
Round 1	Round 2	Round 3	Round 4	Round 5
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35	41	—	—	—
20	—	—	—	—

$$\begin{aligned}\text{MoME}(2, 2) &= 112 - (46 + 41) \\ &= 112 - 87 \\ &= 25\end{aligned}$$


Multiple Elimination as a Detour



Usable Multiple Eliminations



Which path has the largest path bottleneck?

Best-Path Lower Bound

- The largest path bottleneck ...
 - of all paths from round 1 to round C ...
 - that consist of usable multiple eliminations.
- A best path:
 - Guarantees the same winner
 - Maximizes tolerance for additions and removals among usable multiple elimination paths

Best-Path Lower Bound Algorithms

- $O(C^2 \log C)$ time to sort the vote totals within each round.
- The best path can be found in $O(C^2)$ time.
 - Using a bottleneck algorithm, which is ...
 - A longest path algorithm for a weighted directed acyclic graph, but calculating the length as the minimum of its component parts, instead of the sum.

Estimate Relations

Single-Elimination-Path lower bound

≤ Best-Path lower bound

≤ **margin of victory**

≤ Winner-Survival upper bound

≤ Last-Two-Candidates upper bound

Early-Termination Estimates

- For tabulations that stop before C-1 rounds
 - when a candidate has a majority of the continuing votes
 - more than two candidates are in the round
- Accuracy is degraded
 - must allow for possible extreme behavior in the missing rounds of the tabulation.

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- What are we talking about?
- Estimates – **quick: $O(C^2 \log C)$ time**
- **Worst-case accuracy**
- Real elections
- Conclusions

Asymptotic Worst-Case Accuracy

- Ratio with margin of victory is unbounded.

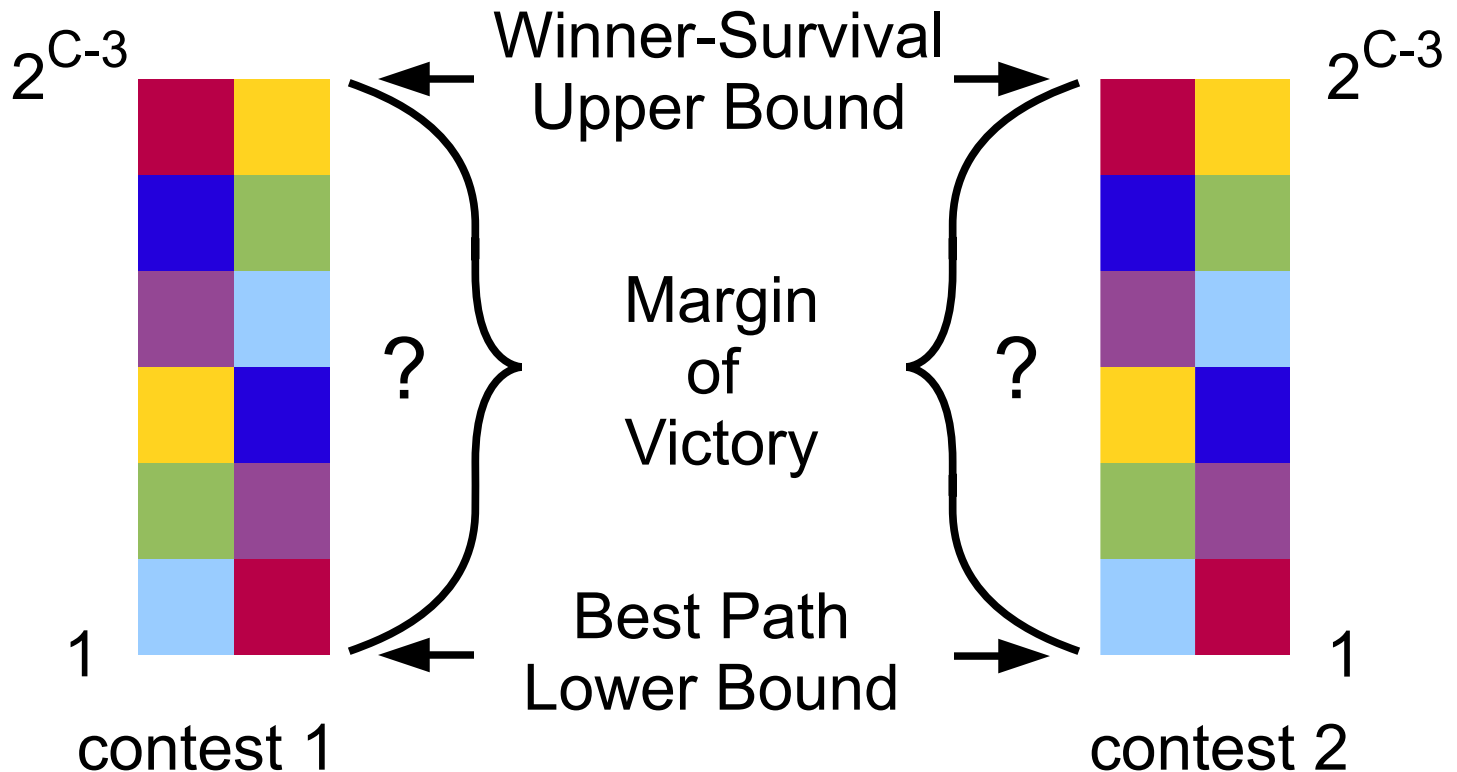
$$\frac{\text{Winner-Survival Upper Bound}}{\text{Margin of Victory}}$$

$$\frac{\text{Margin of Victory}}{\text{Best-Path Lower Bound}}$$

- No estimate can do better if based only on tabulation vote totals.

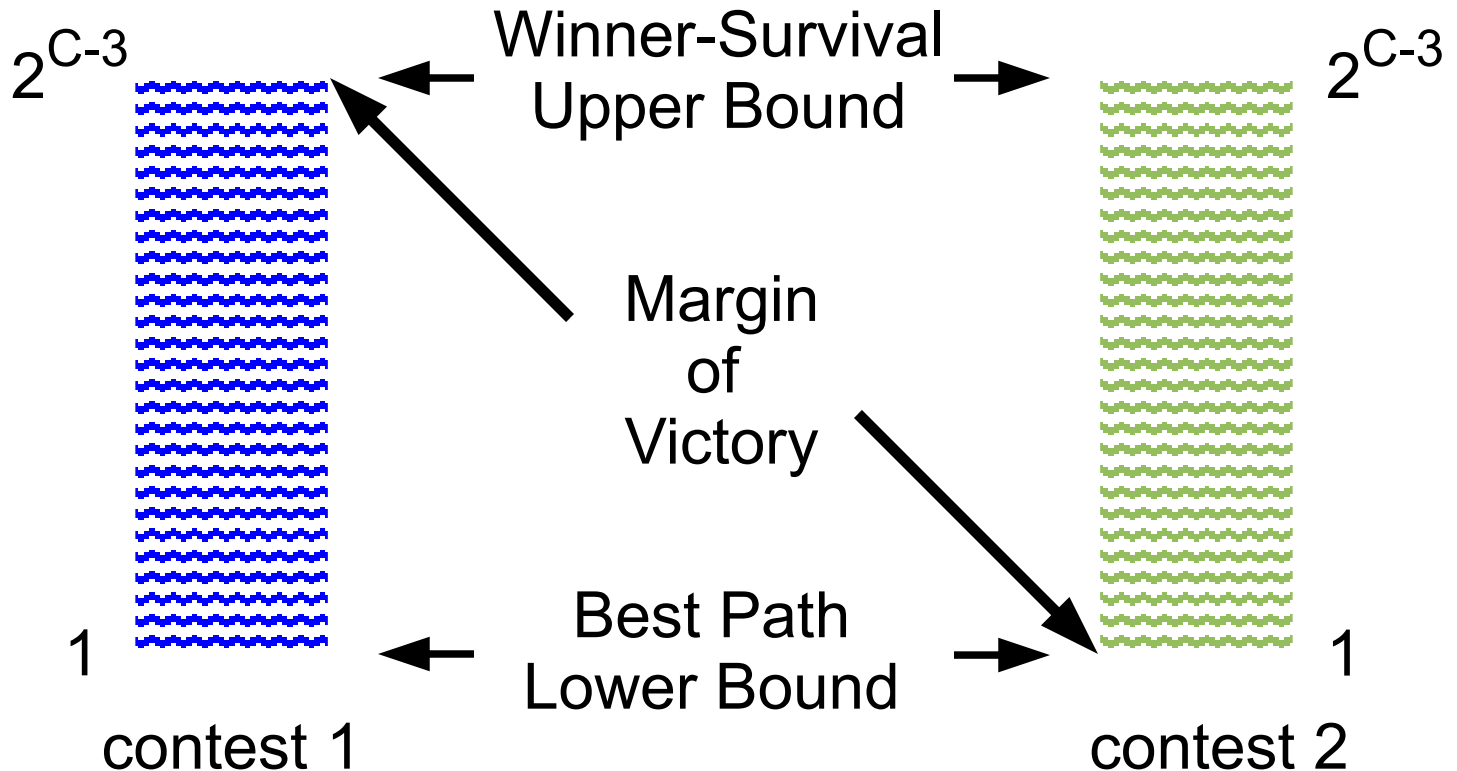
Asymptotic Worst-Case Example

Identical Tabulation Vote Totals



Asymptotic Worst-Case Example

Ballots Show
Different Margins of Victory



Overview

- Why estimate? **to do risk-limiting audits**
- What are we talking about?
- Estimates – **quick: $O(C^2 \log C)$ time**
- Worst-case accuracy – **unbounded ratios**
- **Real elections**
- Conclusions

Estimates for Real Elections

- Australia elections, 2010
 - national House of Representatives
 - 150 contests
- All California IRV contests since 2004
 - local, non-partisan elections
 - 53 contests
 - 36 from San Francisco, 2004-2011
 - 12 using early termination estimates
 - 17 from Alameda county, 2010: Berkeley, Oakland, and San Leandro

Evaluating Estimates

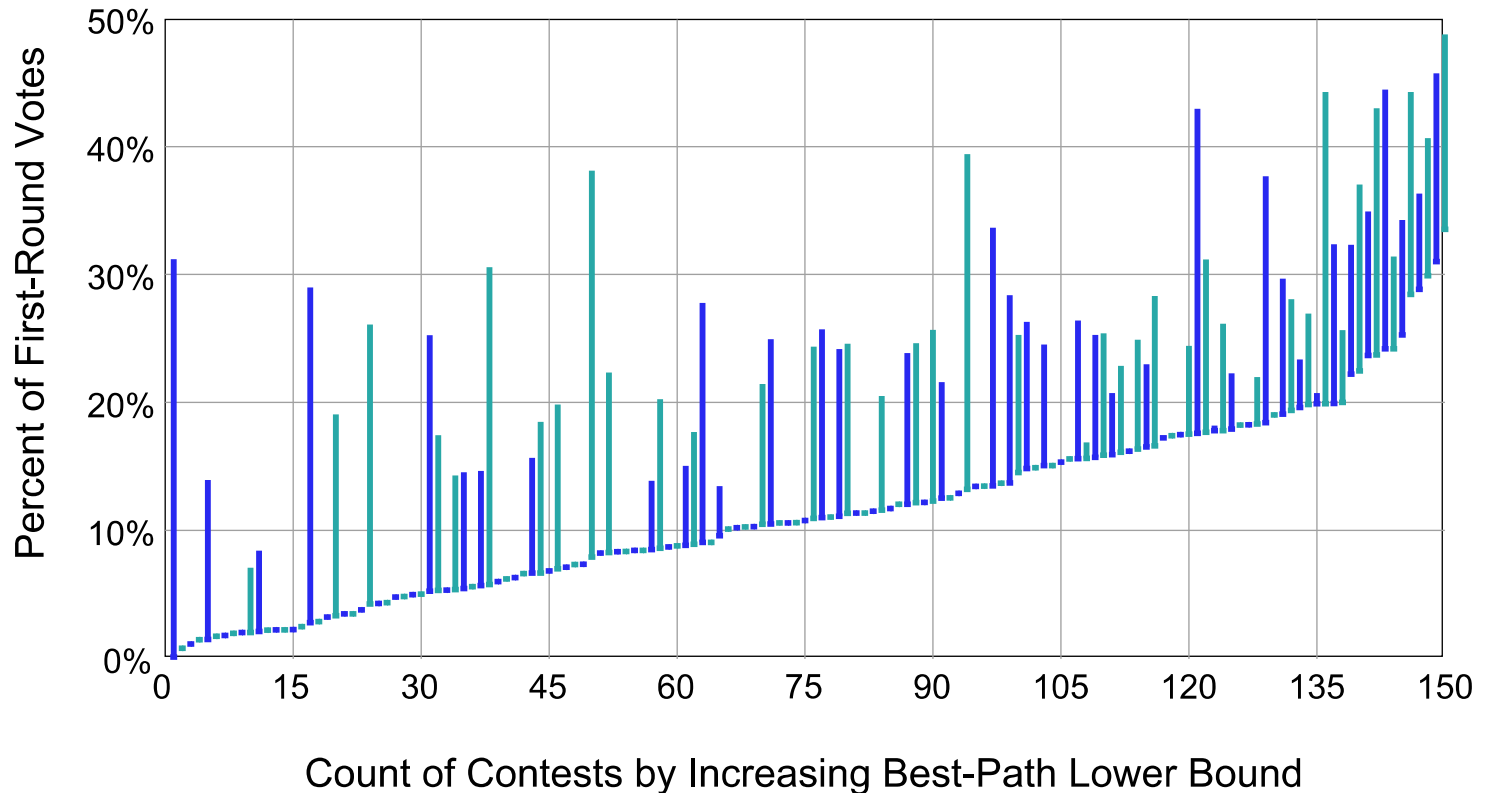
- There are many ways to analyze the data.
- What are relevant metrics?
- A full evaluation requires a context of:
 - specific risk-limiting audit protocols
 - profiles of audit differences.
- Look at:
 - best available lower bound and upper bound,
 - as a percentage of first-round votes.
- What is the distribution of estimates?

Selected Stats

<u>Assessment</u>	<u>Australia</u>		<u>California</u>	
Total Contests	150	100%	53	100%
Contests with LB > 10%	85	57%	35	66%
Contests with LB < 5%	28	19%	14	26%
Contests with LB < 1%	2	1%	7	13%
Contests with LB=MoV=UB	71	47%	16	30%
and LB < 5%	21	14%	4	8%
and LB < 1%	1	1%	0	0%
Contests with UB/LB > 2	34	23%	10	19%
and LB < 5%	7	5%	7	13%
and LB < 1%	1	1%	7	13%

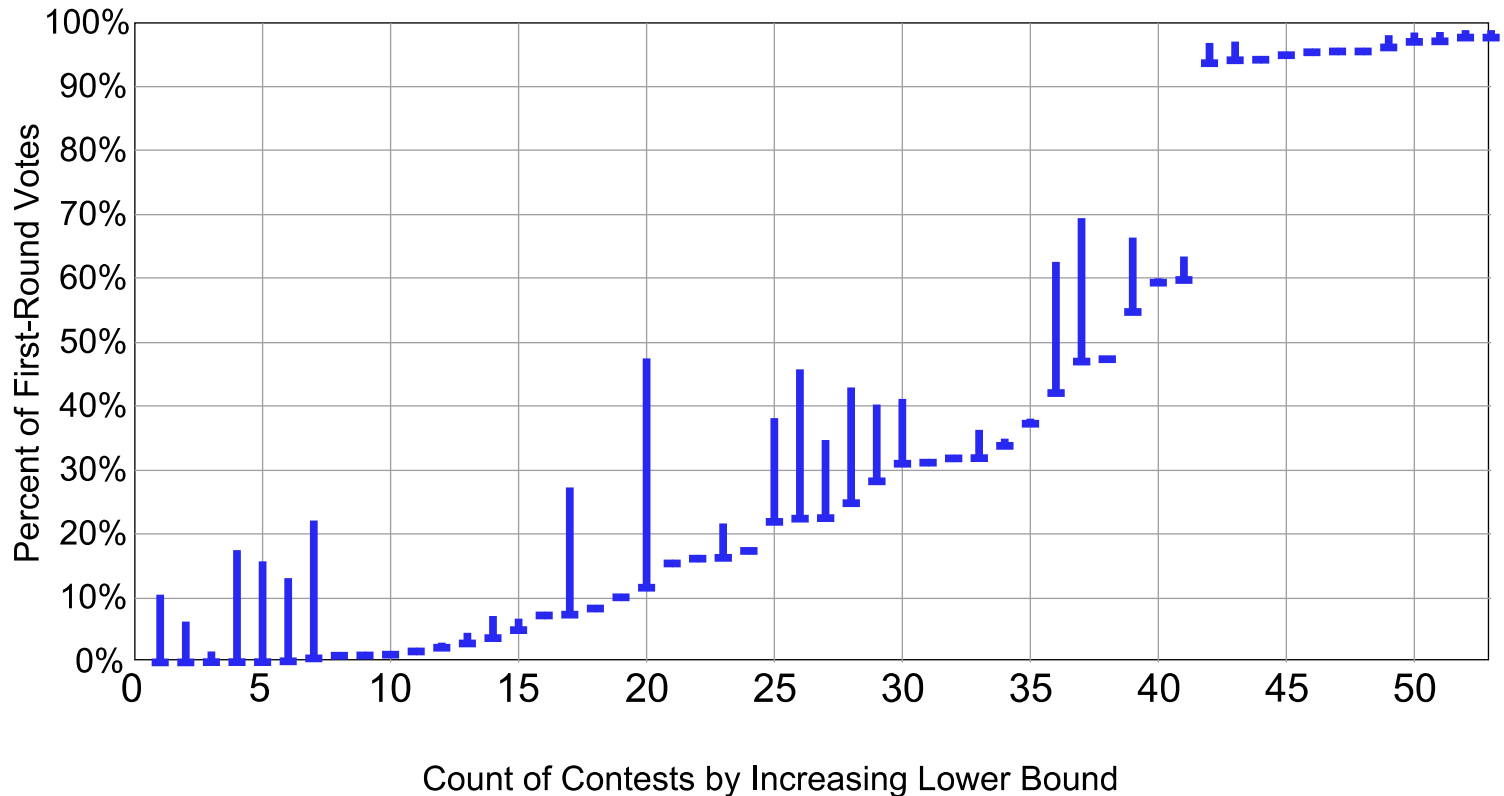
Australia Elections

Margin of Victory, Range of Best Estimates
Australia 2010 House of Representatives



California Elections

Margin of Victory, Range of Best Estimates
California Local IRV Elections, 2004-2010



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- Why estimate? **to do risk-limiting audits**
- What are we talking about?
- Estimates – **quick: $O(C^2 \log C)$ time**
- Worst-case accuracy – **unbounded ratios**
- Real elections – **some estimates useful, some need improvement**
- **Conclusions**

Conclusions

- Risk-limiting audits can use lower bounds for the margin of victory.
- Estimates can be quickly calculated from tabulation vote totals.
- Worst-case ratios with the margin of victory are unbounded.
- The Best-Path lower bound can be used for some risk-limiting audits, but some contests will need better estimates.

Thanks

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