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# Exertion-based billing for cloud storage access

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# Cloud accounting

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- Infrastructure-as-a-Service (IaaS)
  - Provider makes resources available to clients
  - Clients pay for resources used
- Provider wants to recover costs
- Client wants to pay fairly for use
- This talk: Focus on storage

# Provider: Recovering costs

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- More bytes stored → more disks needed
  - Bill for capacity
  - ✓ Providers do this
- More time spent on requests → more disks needed
  - Bill for access
  - ✗ Not all providers do this
  - ✗ Those that do, use the wrong metric

vCloud Express

# Billing for access

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- Providers currently bill for:
  - IOs Amazon EBS, S3, Windows Azure
  - Bytes transferred Amazon S3
  - Performance Amazon EC2

**✗ None matches  $\text{time used} \times \text{cost of resources}$**

# Unsustainable approach: Fixed cost per IO

*Suppose provider charges \$0.01 per IO  
but it takes \$0.01/ms to operate a disk*

Request scenario	Disk time	Real cost	Bill
Cache miss, sequential	1 ms	\$0.01	\$0.01
Cache miss, long seek	20 ms	\$0.20 Provider loses money	\$0.01
Cache hit	0 ms	\$0.00 Client pays too much	\$0.01

# Alternatives

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- Charge per byte or performance
  - ✗ Both also vary as in previous example
- Charge for worst case (\$0.20/IO instead of \$0.01)
  - ✗ Clients with “easy” requests pay too much
- Charge for the average case
  - ✗ Some clients subsidize others
  - ✗ Will the average change over time?

# Solution: Charge for disk time

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*Suppose it takes \$0.01/ms to operate a disk and the provider charges \$0.01/ms*

Request scenario	Disk time	Real cost	Bill
Cache miss, sequential	1 ms	\$0.01	\$0.01
Cache miss, long seek	20 ms	\$0.20	\$0.20
Cache hit	0 ms	\$0.00	\$0.00

✓ Costs recovered fairly

# But, a technical problem remains...

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- Resources may be shared across tenants
- Workload A may be affected by Workload B
  - Workload A sequential, Workload B disrupts locality
  - Workload A cacheable, Workload B evicts A's pages
- ✗ Workload B drives up Workload A's bill



# Interference example

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Workload	Standalone exertion	Exertion w/ interference
Sequential 1 MB/s	1.6%	23%
Random 1 MB/s	67%	74%

✗ Random workload increases Sequential workload's exertion

# How to solve?

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- Bill Workload B for its impact on Workload A?
  - ✗ Not the “fault” of B that it coexists poorly with A
- Provider absorbs the cost?
  - ✗ Provider will either lose money or pass it along as a hidden cost to customers
- ✓ Avoid the interference in the first place

# Performance insulation is needed

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- *Performance insulation*: System property
  - Another workload is not able to affect your...
    - Efficiency
    - Exertion
    - Performance (in your share of time)
  - ... beyond a small (e.g. 10%) factor
- Strictly limits transient influences

# How to insulate storage

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- Preserve locality
  - Disk-head timeslicing
    - Seek between workloads infrequently
- Provide predictable cache allocation
  - Cache partitioning
    - Partition sizes based on access patterns
- Argon storage system [FAST 2007]

# Insulation example

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*Insulation limits impact of other workload to < 10%*

Workload	Standalone exertion	Exertion w/ insulation
Sequential 1 MB/s	1.6%	$\leq 1.8\%$
Random 1 MB/s	67%	$\leq 75\%$

- ✓ Each workload's exertion is close to ideal because it receives dedicated disk time & cache space

# Summary

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- Clients should pay for resources used
- Storage: both capacity and access cost money
- Disk time is what costs, should be the metric
- Bills should be independent and predictable
- Performance insulation is needed to make it fair