Freeware for Cluster Computing

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CyClone

- 144 nodes, 272 Pentiums
    - 1 PII/200
    - 256 P200MMX
    - 3 Cyrix 586/200
    - 12 P90
- “Cluster Fat Tree”:
    - 25 3COM 3C3000 100Bt switches
    - 2 3Com 3C9000 1000Bt switches
- Tiled display
- VIA network for 10 nodes (8 switched, 2 point-point)

Applications

- 3D rendering to replace Sarnoff’s 1024-processor SIMD engine (“Princeton Engine”)
- MPEG prototype encoding development
- Ray Tracing for Optics
- Creating encoded files
- Other device simulation, data transformation tasks
- Research in metacomputing

3D rendering to replace a 1024-processor SIMD engine

- PE is a 1024-processor SIMD supercomputer
- Equivalent to 16K CM-2 except PE has very fast video I/O
- PE: 4 fps, 256^3 cube
- 128-processors on linux: 2 fps
- 32-node paragon at NASA: .25 fps, 128^3 cube

MPEG prototype encoding development

- MPEG motion vector estimation is embarassingly parallel with right software
- Limit before we got to sarnoff was 4 machines scaling
- we have “task bag” software with task dispatch/gather overhead of 2 ms/task
- Task bag supports hierarchy
- Can now easily scale to full CyClone cluster

Ray Tracing for Optics

- We’re using TNT (The Next Taskbag) to support an optical simulation
- Rays are traced via lens configurations
- This step follows analytical analyis and precedes actually building it
- Being used to support a seamless tiled display
- Before: 10K rays overnight on a powermac
- After: 10^9 rays in 1-2 hours on full cluster

Creating encoded files

- Can in the limit perform faster-than-real-time encoding
- Paid for the first cluster in a few months.
- Note: if you have seen parallel MPEG papers, and they’re not from Sarnoff, they’re probably wrong (all the ones we’ve seen to date are wrong ...)

Other device simulation, data transformation tasks

- They come out of the woodwork as soon as the cluster appears ...
- Result: our two clusters (‘94 and ‘97 models) paid for themselves in six months
- BUT: none of the apps were traditional scientific apps such as:

  SOR, LU, Matrix Multiply, SPLASH 2, etc.
- These apps don’t scale on most clusters anyway (unless it’s 4 nodes or so)

Research in metacomputing

- DARPA paid for half the nodes
- Research in private name spaces, better control of TCP connections, and “Network Threads”

Background

- This is fifth in a series, going back to 1991
- Initially we built workstation clusters, but cut over to Pentiums in 1994
- First 16, then 32, then 48, ...
- Comparison:
    - 1994, 16 Indys, $327K
    - 1994, 16 P90s, $36K
    - P90s ran at 90% of Indys for apps of interest
    - Performance/Price was 8:1
    - And we could get source ...
- That ended the need for workstations ...

The case for clusters

- Answer: Clustering became practical in
    - 1997
    - 1994
    - 1991
    - 1984
- And the question is:
- “What is 1984”

1984-1993

- Mary Mock shows how to beat a 7600
- Tom Nash and others at Fermi Labs start construction of “Crates”
- By 1991, clusters were making money for IBM, Sun (in internal use)
- Sun MICA 128-node cluster
- IBM basement full of RS/6000s
- ca. 1993, 1-2% of the world’s supercomputers were retired by clusters
- HP, IBM, Sun, DEC: 100 TFLOPS box will be a cluster

Progress since 1984

- What has happened: hand-built clusters for turnkey applications
- What has happened in some cases:
  - cc myprogram.c
  - a.out
- We had this environment at SRC in 1994
- But people still can not:
  - Watch
  - Debug
  - Easily control
- their program *as an entity*

Are Clusters Multiprocessors?

- Can we pretend:
  - Nodes are processors
  - Network is backplane bus
  - OSes should share name spaces, memory, paging, PID space, etc.?
- In our experience, no
- The cluster::=multiprocessor analogy does not work for us
- The two models have different reliability models, failure modes, latencies, etc.
- Our approach: Process-centric

Process-centric

- We focus on the top-level process, not a single computer or cluster
- Process locates resources and attaches as needed (Via Private Name Space)
- Process creates shared memory segments for export or imports other process's shared memory segments (via Zounds)
- Process efficiently creates groups of processes for remote execution (via vex library)
- Obviously, this is freeware, or I wouldn't be here ...

ZOUNDS

- Quick overview of ZOUNDS
- Rationale
- Programming Interface
- Performance
- Applications
- Conclusions

Zounds is part of a series of DSMs

- Mether (started 1988)
- MNFS (started 1992)
- CMA (started 1994, dropped 1995)
  - User-mode DSM built in C++
  - I really don’t like C++ that much any more ...
- To understand ZOUNDS, need to understand Mether/MNFS

Mether (1988-91)

- Software distributed shared memory
- Use custom protocols on Ethernet
- Supported two page sizes: 32 bytes, 8192 bytes
- Integrated into SunOS VM as a device driver
- Simple consistency model (WORM)
- User-controlled coherency

  Standard cache ops such as purge, update, invalidate from user-mode

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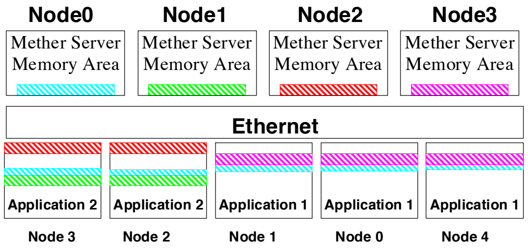
Mether Clients/Servers



- Memory Servers available throughout the network
- Client applications can run on Servers
- Uses UDP, SunOS device driver layers

MNFS (1992-)

- Modified NFS which supports Mether model
- Requires modifications to server and client
- Runs on SunOS, Solaris, Irix, AIX 3.2, and Solaris, FreeBSD 2.0.5R, NetBSD 0.9
- Use to support applications on:
  - 48-node SRC Cluster
  - Aurora Gigabit Network Testbed
  - 100+ nodes at a govt. site
  - U. Koeln in Germany

ZOUNDS Goals

- Simple, low-overhead coherency model
    - A la DEC Memory Channel, SCRAMNet
- Process-centric, not OS-centric, DSM:
    - Clients self-page via SEGV handler (< 10 *u*S)
    - Servers serve from a process image
    - Support multiple memory object types
- Designed to be a good match to the MINI Virtual Interface ATM card
    - I.E. Process-driven, not OS-driven, IO
- Not tied to any particular OS or transport
- Provide kernel/user client and server implementations

ZOUNDS

- Simple library of client and server calls
- Currently supports TCP/IP connections as well as IP multicast
- Experience has shown that non-DSM experts can easily parallelize code
- Not OS-Specific
- Use of IP multicast for updates is (to our knowledge) unique

Zounds library

- All server functions start with the letters *zs*
- All Client functions start with the letters *zc*
- Any user program can be a server-- even an ordinary application
- The program can start the server code and have it run asynchronously
- Client programs are self-paged: allows control of policy such as page size
- No special OS support is required for operation, save on older version of ***BSD
- OS extensions can improve performance

Zounds Server

- Servers are multithreaded or single-threaded (determined by programmer)
- Any application can be a server
- Servers can issue I/O requests for pages to clients -- I/O need not be only client-driven
- Servers track which clients have which pages
- Any data region can be the backing store: mapped files, arrays, SYSV shared memory

Sample Code

ZSINFO *zs;	**Pointer declaration**
zs = zsalloc();	**Allocate the server**
zs->maxclients = 1;	**Set parameters: maxclients**
zs->size = 16384;	**... size of the region**
zs->key = 0;	**... "key"**
zssetup(zs);	**Set up the server.**
zservershow(stdout, zs);	**Show the server status**
if (nodetach)	**Check "nodetach", run:**
zserve(zs);	**Synchronously until done**
else	
zsdetach(zs);	**Asynchronously**

Comments on server code

- There can be more than one server set up for:
  - Same region
  - Different region
  - Overlapping region
- For the asynchronous case, to make the server exit:
  - Set zs->zsexit to 1 (causes that server to exit)
  - zsexit to 1 (causes all servers to exit)
- In the asynchronous case, the server can send updates/invalidates to clients

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ZOUNDS Client

- Clients attach one or more segments from one or more servers
- Clients can attach some or all of the remote segment
- Clients can page from themselves
- Clients can cause pages to be returned to servers, and cause the server to send:
  - invalidates to other clients
  - updates to other clients

Client Code Fragment

ZCINFO zc;	**Declare the client structure**
char *name = "c097/2000";	**This is the name (normally from argv)**
zc.off = 0;	**Set offset into remote segment**
zc.v = 0;	**Allow zounds to pick location in client memory**
zc.size = 16384;	**Set the size**
zcsetup(&zc, name);	**Call setup function.**

Using the client segment

- You can reference it as memory or do “I/O”

```
if (doread)                              doread indicates “I/O” path
  zcread(zc, (off_t) 0, size);           read from server
 else
  {
    int n;
    n = * (int *) zc->v;                 Reference the data
    zcinvalidate(zc, (off_t) 0, size);Throw the reference away
  }
```

- I/O and Memory references can be interspersed and remain consistent

Multicast Setup

- Multicast support is useful for many types of applications
- For servers: set up a server with the normal path, then add a multicast port to it:

  zsmulti(&zs,mcastip);

- For clients: setu up a client and add a multicast port:

  zcmulti(&zc, mcastip);

Additional Server Multicast Ops

- Servers can do multicast sync operations

```
for(; ! zsexit; ) {
        t.tv_sec = 1;
        t.tv_usec = 0;
        zserve_timeout(&zs, &t);
        /* sync all the pages for zs, starting at 0, dirty or not */
        zsmultisync(&zs, (off_t) 0, (size_t) 128, 0);
}
```

- Once a second, update all multicast clients

Multicast Client Ops

- Clients can accept a multicast update:

  Numbytes = zcmultiupdate(zc);

- zcmultiupdate consumes all pending multicast packets for the client

Applications

- Heat transfer solver
- Distributed tiling using simulated annealing (used multicast heavily)
- My favorite: world's most expensive screen saver (video)

Performance

- SEGV handler performance:
    - 7 microseconds, Linux
    - 11 microseconds, FreeBSD
    - Times can be reduced with some careful redo in the kernel fault path:FreeBSD, 5.5 uS
- Page fault on FreeBSD: 1.22 mS/4096 byte page, 2.2 mS/16384 byte page, 512 microsecond/512 byte page
- Page fault on Linux
- Multicast Update: A server can send at least 200 updates/second without loss at the client

Conclusions

- Zounds is simple:

  To program

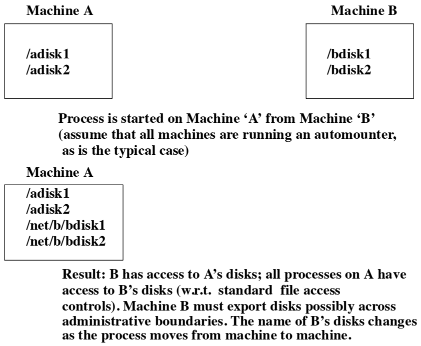
  To use

  To understand performance of

  To understand communications of
- Page fault performance is as good as or better than NFS/MNFS
- Overall performance is quite good
- Policies are easily tuned by the user, should they be so bold

Global Name Spaces

- Global Name Spaces are a security and administration nightmare



Machine A

/adisk1
/adisk2

Machine B

/bdisk1
/bdisk2

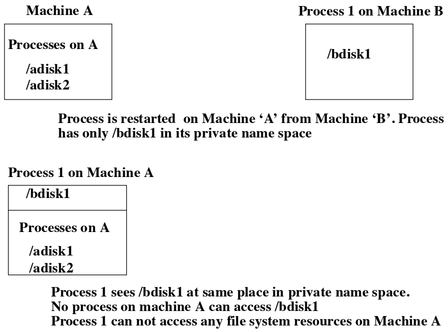
Process is started on Machine 'A' from Machine 'B' (assume that all machines are running an automounter, as is the typical case)

Machine A

/adisk1
/adisk2
/net/b/bdisk1
/net/b/bdisk2

Result: B has access to A's disks; all processes on A have access to B's disks (w.r.t. standard file access controls). Machine B must export disks possibly across administrative boundaries. The name of B's disks changes as the process moves from machine to machine.

Private Name Spaces



Machine A

Processes on A

/adisk1
/adisk2

Process 1 on Machine B

/bdisk1

Process is restarted on Machine ‘A’ from Machine ‘B’. Process has only /bdisk1 in its private name space

Process 1 on Machine A

/bdisk1

Processes on A

/adisk1
/adisk2

Process 1 sees /bdisk1 at same place in private name space.
No process on machine A can access /bdisk1
Process 1 can not access any file system resources on Machine A

Private Name Space Advantages

- User-level mount protocol
- Improved security as a result of reduced unintended sharing and fewer priveleged processes
- Checkpoint/Restart is much simpler
- Processes need less access to system resources (e.g. file system) to access files
- No need to convince sysadmins to export resources across organizations

Current Implementation

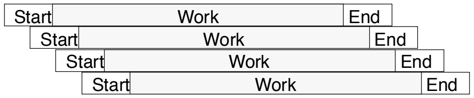
- User-mode on some OSes, kernel mode client on Linux
- One server process per client connection
- Name space inheritance works both locally and on different machines
- Complete transparency (tested on Linux)

How to get it

- www.sarnoff.com:8000
- Go to the metacomputing page and look for the source
- Source is covered under the GNU Programming License

VEX

- Problem: for reasonable numbers of nodes (say 64 or more) it takes too long to start up remote processes



- Ideally, ‘Start’ and ‘End’ are zero

Does this really matter

- Not usually for small numbers (8 or so)
- But: consider a 128-minute job

Time	8	4	2	1
Procs	16	32	64	128
Overhead	7	3	1.5	.6

- Need fast, efficient encapsulation of remote processors for scalable computing
- Need easy ways to specify aggregates

Hostlists

- Efficient way to specify, select, and communicate with aggregates of remote hosts
- Model is to create a hostlist, then apply filters to it
- Filters can be specified by regular expressions on host names or by selecting idle hosts via vector RPC
- Can generate code for static, initialized hostlists
- At end of filter step(s), operate on the hostlist

Basic Types

- A host:

```
struct host
{
  char *hostname;
  struct sockaddr_in hostaddr;
  /* other good stuff */
  char *hosttype; /* quote string of the form "sun4c", "rs/6000", etc.*/
  int mbytes; /* number of mbytes of virtual memory */
  int pmbytes; /* number of mbytes of physical memory */
  int avenrun[3]; /* from avenrun stats */
  int lastmouse, lastkbd;
};
```

- A list of hosts

```
struct hostlist
{
int numhosts;
struct host **hostarr; /* array of pointers to host entries */
};
```

- A hostlist can be sparse

Basic functions

- Many return a hostlist
- Either:

Take a hostlist as an arg, for filtering

Filter to refine the hostlist, to select only certain hosts (via regex or load or ...)

Add information to the list (.e.g given a set of names, fill in the IP addresses)

Netexec

- Used to manage aggregates of remote programs

```
struct netexec
 {
   struct sockaddr_in s;
   char *hostname;
   char *argv, *envp;
   int arglen, argc;
   int envlen, envc;
   int results;
   int fd0, fd1, fd2;
 };
```

Vector RPC

- Used to efficiently call RPC for aggregates of hosts
- hostlists library can generate vector RPC structures and perform vector RPC calls starting with a hostlist

Simple example: alloc a hostlist, filter by RE, run /bin/date on those hosts

```
struct hostlist *hl;
struct netxec *pn;
char *name, *pass, *cmd = “/bin/date”;
char *linuxclusterhosts[] = {“c0[012345].”, “c06[0123]”, 0};
hl = hostlistfromhosts();
selecthostsbynamelist(hl, linuxclusterhosts);
nprocs = hl->numhosts;
res = rexechl(hl, &nx, &nprocs, &name, &pass, cmd, /* numcmds */ 1,
            exec_port, use_priv_port, /* no stderr */ 0);
```

Details

- We only show one command
  can have arrays of commands
  e.g. run odd/even hosts w/different
- nprocs is filled in with successful procs
- nx is filled in with useable netexec struct
- We found with more than 40 or so machines that hostname lookup hurt scaling, so:

Compiled-in host-IP mappings

- Host->IP address lookup was killing scaleup
- IP addresses in most environments are static for years (esp. true for cluster)
- so:

```
hl = hostlistfromhosts(); /* create a hostlist */
selecthostsbynamelist(hl, selections); /* select hosts we want */
gencode(file, "allhosts", hl); /* gen code for inclusion at compile-time*/
```

Gencode output

```
struct host allhostshosts [] = {
{ "p0", {2, 0, {0xb80a2182}}, "NOTYPE", 0, 0,
   {0, 0, 0}, 0, 0},
{ "p1", {2, 0, {0xb90a2182}}, "NOTYPE", 0, 0,
   {0, 0, 0}, 0, 0},
 /* etc. */
```

- Typical use: hl = clonehl(allhostshosts);
- Then filter hl
- Reduces lookup time for 100 hosts from 5-12 seconds to “zero”
- Removes hostname lookup as a factor

VEX: making it easy

- VEX encapsulates hostlist and netexec in one object
- Code looks like this:

```
vex = vexalloc();
vexaddpplist(vex, argv[0], /* regular expression ? 1 */ 1);
if (vexexec(vex) <= 0)
  exit(1);
  while (vex->succ)  /* vex->succ is number of active remote execs */
   {
     vexioloop(vex); /* handles stdin->remote and remote->stdout */
   }
```

Things we don't do

- Signal propagation

  This is for efficient remote exec: we didn't want to pay the overhead

  We can revisit this decision ...

Conclusions

- Successful cluster computing on a large scale requires low-overhead, process-oriented support systems
- Clusters aren't multiprocessors
- We have shown three such systems
- Two (ZOUNDS, hostlist) have been in use for years for real work
- Private name spaces is just coming into use, but solves major problems we have experienced
- This is all open source